## Patterns and Partitions

Yotam Smilansky

Experimental Mathematics Seminar, Rutgers University



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2. Color frequencies:

- Does $\left|r_{m}\right| /\left|x_{m}\right|$ converge?
- Does $\mathcal{L}\left(\left\{x \in \mathcal{I}: x\right.\right.$ is colored red in $\left.\left.\pi_{m}\right\}\right)$ converge?
- In case both limits exist, are they necessarily the same?


## Multiscale substitution schemes

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Prototiles in $\mathbb{R}^{d}$

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## Directed weighted graph model for substitution schemes



Vertices model the prototiles.
Edges originating in a vertex model the tiles appearing in the substitution rule pattern of the prototile.
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## Uniform distribution results

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$\left|r_{m}\right|=\sharp\left\{\right.$ metric paths of length $I_{m}$ terminating on red edge $\}$
$\mathcal{L}\left(\left\{x \in \mathcal{I}: x\right.\right.$ is colored red in $\left.\left.\pi_{m}\right\}\right)$ is the probability that a metric path of length $I_{m}$ terminates on the red edge, if the red edge is assigned probability $\frac{1}{3}$ and the blue edge probability $\frac{2}{3}$.

## Incommensurable and commensurable schemes

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## More examples

A commensurable example - The Rauzy fractal scheme:


Edge lengths: $\log \tau, 2 \log \tau, 3 \log \tau$, where $\tau=$ tribonacci constant.

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For a.e $\theta$ Sadun's generalized pinwheel scheme is incommensurable:



## Multiscale substitution tilings (jointly with Yaar Solomon)

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Tilings of $\mathbb{R}^{d}$ are defined as limits of $\left\{F_{t}(T): t \geq \mathbb{R}\right\}$.
Our study includes:

- Structural, geometrical and statistical properties of tilings: (types and scales, repetitivity, patch frequencies, BD/BL)
- Dynamical properties of the tiling dynamical system. (minimality, invariant measures)










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This Kakutani sequence does not have color frequencies.

The incommensurable case - counting paths on graphs

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M_{i j}(s)=e^{-s \cdot l\left(\varepsilon_{1}\right)}+\cdots+e^{-s \cdot l\left(\varepsilon_{k_{i j}}\right)}
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and $M_{i j}(s)=0$ if there are no such edges in $G$.
Theorem (Kiro, Smilansky $\times 2$ ): Let $G$ be a strongly connected incommensurable graph. There exist $\lambda>0$ and $Q \in M_{n}(\mathbb{R})$ with positive entries, such that if $\varepsilon \in \mathcal{E}$ has initial vertex $h \in \mathcal{V}$, the number of metric paths of length exactly $x$ from vertex $i \in \mathcal{V}$ to a point on the edge $\varepsilon$ grows as

$$
\frac{1-e^{-l(\varepsilon) \lambda}}{\lambda} Q_{i h} e^{\lambda x}+o\left(e^{\lambda x}\right), \quad x \rightarrow \infty
$$

where $\lambda$ is the maximal real value for which $\rho(M(\lambda))=1$,

$$
Q=\frac{\operatorname{adj}(I-M(\lambda))}{-\operatorname{tr}\left(\operatorname{adj}(I-M(\lambda)) \cdot M^{\prime}(\lambda)\right)}
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## Poles of the Laplace transform

The proof follows The Wiener-Ikehara Theorem, originally motivated by the Prime Number Theorem.

This requires the study of the poles of the Laplace transform of a counting function, which in our case is given by

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\mathcal{L}\{f(x)\}(s)=\frac{1-e^{-I(\varepsilon) s}}{s} \cdot \frac{(\operatorname{adj}(I-M(s)))_{i h}}{\operatorname{det}(I-M(s))}
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and so we study the zeroes of the exponential polynomial

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Information on the location of zeroes closest to $\operatorname{Re}(s)=\lambda$ can be used to obtain upper bounds on error terms.

Zeroes of exponential polynomial (jointly with Avner Kiro, Alon Nishry and Aron Wennman)

In the case of graphs modeling an $\boldsymbol{\alpha}$-Kakutani scheme

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\operatorname{det}(I-M(s))=1-e^{-a s}-e^{-b s}
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Incommensurability is equivalent to

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\beta=\frac{\log \alpha}{\log (1-\alpha)} \notin \mathbb{Q},
$$

and by a change of variables $z=s \log \alpha$, we reduce to the study of roots of $e^{z}+e^{\beta z}=1$.

## Zeroes of exponential polynomial (jointly with Avner Kiro, Alon Nishry and Aron Wennman)

In the case of graphs modeling an $\boldsymbol{\alpha}$-Kakutani scheme

$$
\operatorname{det}(I-M(s))=1-e^{-a s}-e^{-b s}
$$

with $a=\log \frac{1}{\alpha}$ and $b=\log \frac{1}{1-\alpha}$.
Incommensurability is equivalent to

$$
\beta=\frac{\log \alpha}{\log (1-\alpha)} \notin \mathbb{Q},
$$

and by a change of variables $z=s \log \alpha$, we reduce to the study of roots of $e^{z}+e^{\beta z}=1$.

The following slides show some approximations of such zeroes in compact strips, for different values of $\beta$. At the moment these experimentations give rise to more questions than answers...

## $\beta=\varphi$ the golden ratio, rightmost roots (up to 10,000 )



## $\beta=\varphi$, all roots (up to 10,000 )



## $\beta=\varphi$, histogram



## $\beta=e$, rightmost roots (up to 10,000 )



## $\beta=e$, all roots (up to 10,000 )



## $\beta=e$, histogram



## $\beta=\pi$, rightmost roots (up to 10,000 )



## $\beta=\pi$, all roots (up to 10,000 )



## $\beta=\pi$, histogram


$\beta=\ell$, a Liouville number, rightmost roots (up to 30,000 )


## $\beta=\ell$, all roots (up to 30,000 )


$\beta=\ell$, histogram


## Extending the model

Next, we now turn to the roots of

$$
e^{z}+e^{\beta z}+e^{\gamma z}=1
$$

which are related to graph with a vertex and three loops, or to schemes in which $\mathcal{I}$ is substituted by three rescaled copy of itself.

## $\beta=1$ and $\gamma=\varphi$, all roots (up to 10,000 )



## $\beta=1$ and $\gamma=\varphi$, histogram



## $\beta=\sqrt{2}$ and $\gamma=\sqrt{3}$, all roots (up to 10,000 )


$\beta=\sqrt{2}$ and $\gamma=\sqrt{3}$, histogram


## $\beta=e$ and $\gamma=\pi$, all roots (up to 10,000 )



## $\beta=e$ and $\gamma=\pi$, histogram




