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Source: *The American Mathematical Monthly*, Jan., 1993, Vol. 100, No. 1 (Jan., 1993), pp. 62-64

Published by: Taylor & Francis, Ltd. on behalf of the Mathematical Association of America

Stable URL: <https://www.jstor.org/stable/2324818>

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The Computer Solves the Three Tower Problem

Arthur Engel

Consider the following probability problem: *We have three piles with a, b, c chips, respectively. Each second a pile X is selected at random, then another pile Y is chosen at random and a chip is moved from X to Y . Find the expected waiting time $f(a, b, c)$ until one pile is empty. (Three Tower Problem, or TTP.)*

This problem is due to Lennart Råde from Gothenburg University, Sweden. During the last 20 years he posed it to numerous people, but nobody could solve it [5]. I heard of it during a Statistics Conference in New Zealand 1990. It became a simulation exercise in a book I was writing at the time [3]. The simulation problem gives numerical answers to specific inputs a, b, c . On January 17, 1992 I loaded the program again and started to experiment. In 15 minutes I guessed the formula

$$f(a, b, c) = \frac{3abc}{a + b + c}. \quad (1)$$

Once you have guessed the formula, the proof is a routine matter. Start in state (a, b, c) . In one step you are in one of the neighboring states $(a, b + 1, c - 1)$, $(a, b - 1, c + 1)$, $(a + 1, b, c - 1)$, $(a - 1, b, c + 1)$, $(a + 1, b - 1, c)$, $(a - 1, b + 1, c)$ with the same probability $1/6$. So we have

$$f(a, b, c) = 1 + \frac{1}{6} \sum f(x, y, z) \quad (2)$$

over all neighbors (x, y, z) of (a, b, c) with boundary conditions

$$f(a, b, 0) = f(a, 0, c) = f(0, b, c) = 0. \quad (3)$$

It looks pretty hopeless to solve the functional equation (2) with the boundary conditions (3). But thanks to the PC we have the guess (1). It obviously satisfies (3) and a short calculation shows that (2) is also satisfied. So we have a solution to our problem. Its uniqueness can be proved by a standard argument, which we reproduce to make the paper self contained. See [1] or [2].

Suppose that $g(a, b, c)$ is another solution. Consider $h(a, b, c) = f(a, b, c) - g(a, b, c)$. Then

$$h(a, b, c) = \frac{1}{6} \sum h(x, y, z) \quad (4)$$

over all neighbors of (a, b, c) .

The function h is defined for finitely many points. At some of these points h assumes its maximum M . Because of (4) $h(x, y, z) = M$ for all six neighbors of (a, b, c) . And their neighbors have also the same h -value M , and so on, until we reach the boundary, at which h has value 0. Thus $h(a, b, c) \leq 0$ everywhere. Similarly we can show that $-h \leq 0$. Thus $h = 0$, and $f(a, b, c) = g(a, b, c)$ everywhere. So f is unique.

Lecturing in Norway, Råde mentioned that the TTP has recently been solved by me. A listener asked about the expected duration $g(a, b, c)$ of the following modification of the TTP: *Start with three towers. As soon as one tower is empty continue playing with two players until just one is left.* At the lecture it was agreed that this would be a harder problem to solve. Råde challenged me to find $g(a, b, c)$. He added that he also would like to know the probability p_a that the a -tower first becomes empty.

With my PC I started to work empirically on $g(a, b, c)$. Instead of 15 minutes it took me several hours of hard work. The trouble was that I was looking for a more complicated formula. At the end I found the much simpler correct formula

$$g(a, b, c) = ab + bc + ca. \quad (5)$$

It is easy to show that (2) is satisfied when $a, b, c > 0$. The new boundary conditions are

$$g(a, b, 0) = ab, \quad g(a, 0, c) = ac, \quad g(0, b, c) = bc. \quad (6)$$

(6) is the expected duration for the Two Tower Problem. This is a classic result, which is equivalent to the Gamblers Ruin Problem. See [4]. Had I looked at (6) first, they would have immediately suggested (5).

By analogy I was able to write down the solution of the modified n -Tower-Problem:

$$g(x_1, \dots, x_n) = \sum_{i < k} x_i x_k. \quad (7)$$

It was also easy to guess the following version of Råde's second problem: The i th tower is the winner (in the game which continues until one pile is left) with probability

$$p_i = \frac{x_i}{x_1 + \dots + x_n}. \quad (8)$$

Both formulas (7) and (8) satisfy the appropriate recurrences and boundary conditions. The proof of (7) involves induction on n . The boundary conditions for a given value of n are determined by the solution of the problem for $n - 1$. Uniqueness is proved as in the case of the Three-Tower-Problem.

A related result could also be found with my PC searching for several hours: *Players 1, 2, 3 start with a, b, c chips, respectively. In one round each player stakes one chip. Then a 3-sided symmetric die labeled 1, 2, 3 is rolled and the winner gets all the chips staked. If a player is broke the game continues with two players until one player has accumulated all the chips.* The expected number of rounds is

$$h(a, b, c) = ab + bc + ca - \frac{2abc}{a + b + c - 2}. \quad (9)$$

If the game stops as soon as one tower is empty the expected duration is

$$h(a, b, c) = \frac{abc}{a + b + c - 2}. \quad (10)$$

This result was communicated to me by a former IMO contestant Michael Stoll. It was found ten years ago during a summer academy for gifted high school students. Despite huge efforts they were unable to handle four players.

The original Four Tower Problem is still unsolved. I experimented extensively for many hours, but all my guesses turned out to be wrong. $f(a, b, c, d)$ seems to be a very complicated function, as can be seen from the exact value $f(3, 2, 2, 2) = 350612/69969$. No simple formula can give such a complicated result for so small values of a, b, c, d . The only thing I could do was to guess a good approximation

$$f(a, b, c, d) \approx \frac{6abcd}{ab + ac + ad + bc + bd + cd}. \quad (11)$$

It is easy to see that $f(a, b, c, d)$ has the form $p(a, b, c, d)/q(a, b, c, d)$ with polynomials which are symmetric in a, b, c, d . In addition q seems to be constant, depending only on $a + b + c + d$. The use of *Mathematica* may bring more success. I worked numerically with Turbo Pascal as in [3].

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A Linear Algebra Approach to Cyclic Extensions in Galois Theory

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A beginning course in Galois theory often includes a discussion of cyclic extensions, that is, Galois extensions whose Galois groups are cyclic. The usual approach (see, e.g., [1] and [2]) is to derive the results on cyclic extensions as