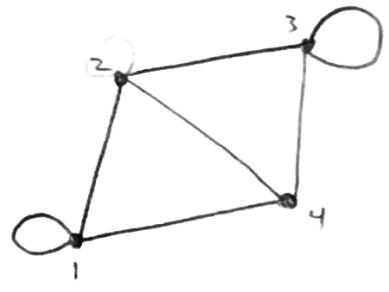


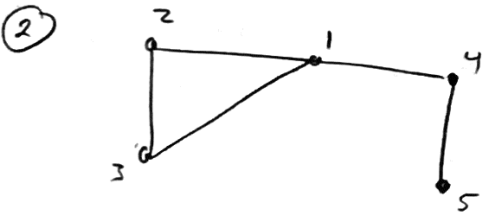
Has none



	1	2	3	4
1	0	1	1	1
2	1	0	1	1
3	1	1	0	1
4	1	1	1	0

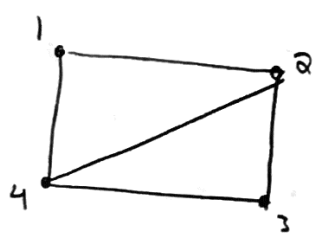
All vertices have odd degrees [sum of row = degree] hence no path or circuit/cycle

	1	2	3	4
1	1	1	0	1
2	1	0	1	1
3	0	1	1	1
4	1	1	1	0



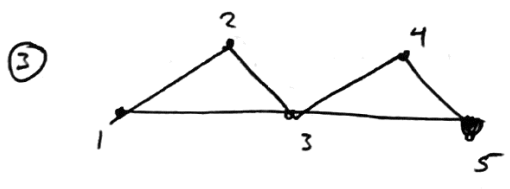
	1	2	3	4	5
1	0	1	1	1	0
2	1	0	1	0	0
3	1	1	0	0	0
4	1	0	0	0	1
5	0	0	0	1	0

5 4 1 2 3 1 is a Eulerian path that starts at 5 and ends at 1.



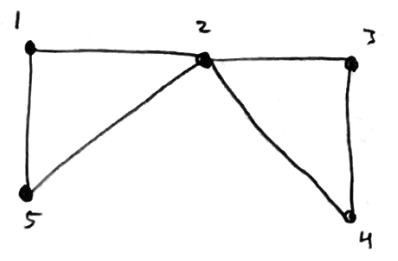
	1	2	3	4
1	0	1	0	1
2	1	0	1	1
3	0	1	0	1
4	1	1	1	0

2 3 4 2 1 4 is a Eulerian path



	1	2	3	4	5
1	0	1	1	0	0
2	1	0	1	0	0
3	1	1	0	1	1
4	0	0	1	0	1
5	0	0	1	1	0

1 2 3 4 5 3 1 is a Eulerian cycle



	1	2	3	4	5
1	0	1	0	0	1
2	0	1	1	1	1
3	0	1	0	1	0
4	0	1	1	0	0
5	1	1	0	0	0

2 1 5 2 3 4 2 is a Eulerian cycle

4) Assume a graph G has a Eulerian circuit/cycle C . Because it is Eulerian, it contains every edge and vertex in the graph. For every vertex $v \in G$ we need to show that v has an even degree. First assume v is not the starting and ending vertex in C . Then v must be in the middle of the cycle. Then everytime v is "touched" in C , there are 2 edges, one that goes into v [v is an endpoint of an edge], and one edge that leaves v [v is a startpoint of an edge]. v can be encountered k times, which means there are $2k$ edges on v , thus has an even degree. The start vertex has an outgoing edge at the beginning and one incoming edge at the end of C , hence all vertices in G are even.

5) Assume graph G has a Eulerian path P that starts at u and ends at v . Then in the path P , it leaves u one more time than an edge enters u , and similarly, edges leave v one less time than it enters v . Therefore using a similar argument as above, it must be that u and v have odd degrees, while all other vertices must be even degrees from an argument similar to 4), hence proved.