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Graph Theory Homework #8

8.1)

Shortest path from A to G : $A \rightarrow B \rightarrow D \rightarrow C \rightarrow F \rightarrow E \rightarrow G$

8.2)

Shortest path from L to A : $L \rightarrow K \rightarrow I \rightarrow F \rightarrow H \rightarrow E \rightarrow B \rightarrow A$

8.3)

Algorithm for constructing the longest path from A to L :

First, run the shortest path algorithm to find the distance of the shortest path from L to each vertex in the graph. Then, we work starting from A . For each neighbor N of A which is not already on the path, calculate $g = \text{weight}(e_{AN}) - (\text{dist}(A) - \text{dist}(N))$, where $\text{weight}(e_{AN})$ is weight of the edge connecting A and N and $\text{dist}(A)$ and $\text{dist}(N)$ are the distances of the shortest paths from L to A and N respectively. The neighbor with the largest value of g is the next vertex in the path. Repeat this procedure with each new vertex in the path, until L is reached.

Longest path from A to L : $A \rightarrow E \rightarrow C \rightarrow F \rightarrow H \rightarrow J \rightarrow L$

8.4)

S to A : $S \rightarrow A$

S to B : $S \rightarrow B$

S to C : $S \rightarrow B \rightarrow C$

S to D : $S \rightarrow A \rightarrow D$

S to E : $S \rightarrow A \rightarrow D \rightarrow E$

S to F : $S \rightarrow B \rightarrow F$

S to T : $S \rightarrow B \rightarrow F \rightarrow T$

8.6)

The solution to the traveling salesman problem is $A \rightarrow B \rightarrow C \rightarrow E \rightarrow D \rightarrow A$ with distance 14.

8.7)

The Hamiltonian cycle with the greatest weight is $A \rightarrow C \rightarrow D \rightarrow B \rightarrow E \rightarrow A$ with a total weight of 32