

Section 6: 6.1, 6.2, 6.3, 6.4, 6.5, 6.6

6.1) Which of the following graphs is Eulerian? Semi-Eulerian?

i) the complete graph  $K_5$ ;Ans: degree =  $5-1 = 4$  (even). Therefore,  $K_5$  is Eulerian.ii) the complete bipartite graph  $K_{2,3}$ ;Ans: degree for  $K_2 = 3$ , degree for  $K_3 = 2$ Therefore,  $K_{2,3}$  is Semi-Eulerian

iii) the graph of the cube

Ans: degree = 3, (odd) Therefore, Cube is Not Eulerian or Semi-Eulerian (neither)

iv) the graph of the octahedron

Ans: degree = 4 (even). Therefore, Octahedron is Eulerian

v) the Petersen graph: 10 vertices, 15 edges

Ans: degree = 3 (odd). Therefore, Petersen Graph is neither

6.2) In the table of Fig 2.9, locate all the Eulerian and semi-Eulerian graphs.

Solution: i) Eulerian: 1, 4, 8, 18, 21, 25, 31

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(i) Semi-Eulerian: 2, 3, 6, 7, 9, 13, 14, 16, 17, 19, 22, 23, 26, 28, 30.

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6.3) (i) For which values of  $n$  is  $K_n$  Eulerian?

Answer:

$K_n$  graphs have  $n-1$  degrees. Therefore,  $K_n$  is Eulerian if  $n-1$  is an even number.

$\therefore K_n$  is Eulerian for values of  $n$ , where  $n-1$  is an even number. e.g.:  $K_3, K_5, K_9, K_{11}, \dots$

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(ii) Which complete bipartite graphs are Eulerian?

Answer: Let's denote the complete bipartite graph as  $K_{r,s}$ . Therefore, the complete bipartite graph  $K_{r,s}$  will be Eulerian if both  $r$  and  $s$  are even numbers.

$\therefore$  Complete bipartite graphs are Eulerian if the two disjoint sets of vertices have even number of vertices.

e.g.:  $K_{2,2}, K_{6,6}, K_{12,10}, \dots$

③

(ii) Which Platonic graphs are Eulerian?

Answer: Platonic graphs are Eulerian if each vertex has an even number of degree. Only the octahedron is Eulerian

This is because the octahedron is the only platonic graph in which all the vertices have even degree

(iv) For which values of  $n$  is the wheel  $W_n$  Eulerian?

Answer: A wheel  $W_n$  has  $2(n-1)$  edges for  $n$  vertices

therefore the degree for each vertex in the cycle is 2, while the degree of the new vertex formed in the central part of the graph will be  $n-1$ . So the wheel will be Eulerian if  $n-1$  is even.

$\therefore W_n$  is Eulerian for values of  $n$  where  $n-1$  is an even number.

E.g:  $W_5, W_7, W_9, \dots$

(v) For which values of  $k$  is the  $k$ -cube  $Q_k$  Eulerian?

Answer: A cubic graph has a degree <sup>regular</sup> of  $k$ . Therefore,  $k$ -cube  $Q_k$  is Eulerian if  $k$  is an even number.

E.g:  $Q_4, Q_6, Q_8, \dots$



4

6.4) Let  $G$  be a connected graph with  $k (> 0)$  vertices of odd degree.

(i) Show that the minimum number of trails, that together include every edge of  $G$  and that have no edges in common is  $k/2$

Solution:  $G$  is a connected graph with  $k$  vertices of odd degree, then  $G$  must have an even number of vertices with odd degrees, in order for  $G$  to meet the Eulerian trail condition for which no edges are in common and the trail includes all edges in  $G$ . We can assume this because the sum of all degrees in a graph is always even, based on the Handshaking Lemma.

Let's assume there are  $2m$  vertices with odd degree in  $G$  where  $m$  is a positive integer i.e.  $m > 0$ . Since each trail will start and end at vertices with odd degrees, then there must be <sup>at least</sup>  $m$  trails <sup>for</sup> all ~~possible~~ vertices with odd degrees which is  $k$  ( $k =$  vertices with all degrees).  
 $\therefore 2m = k, \quad m = k/2$

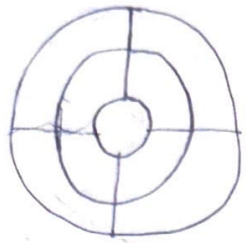
Therefore, the minimum number of trails is  $m$  which is  $\frac{k}{2}$

$\therefore$  This shows that the minimum number of trails, that together include every edge of  $G$  and that have no edges in common is  $k/2$

5

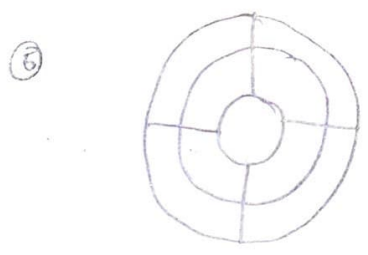
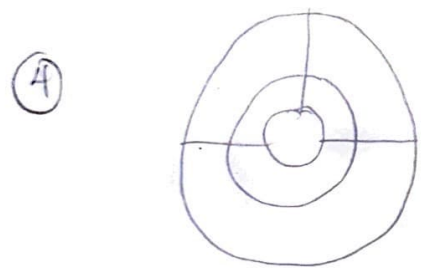
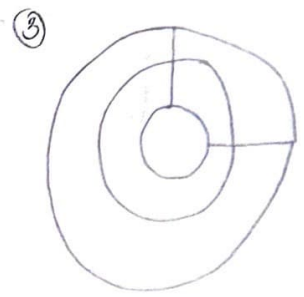
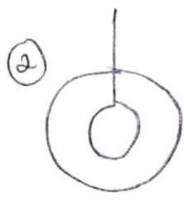
(ii) How many continuous pen-strokes are needed to draw the diagram in Fig 6.8 without repeating any line?

Fig 6.8 :



⇒ a continuous pen stroke involves creating a closed loop without lifting the pen from the paper.

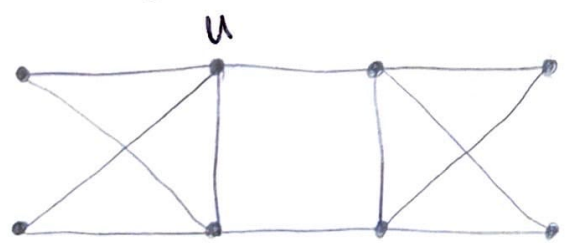
Solution:



∴ There are 5 continuous pen-strokes needed to draw the diagram without repeating any line

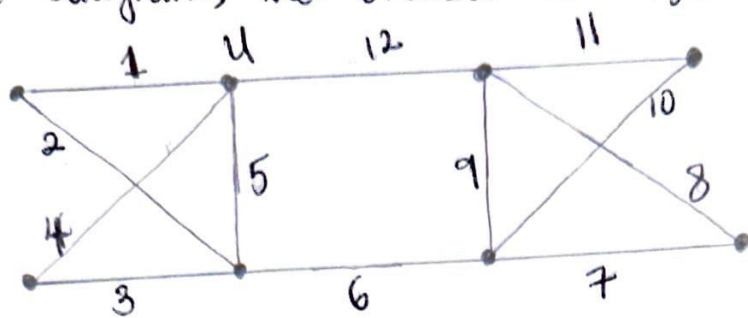
6.5.) Use Fleury's algorithm to produce an Eulerian trail for the graph in Fig 6.9.

Fig 6.9:



②

Solution: start at any vertex  $u$  (this is already given in the diagram) then traverse the edges:



6.6) (i) show that the line graph of a simple Eulerian graph is Eulerian.

Solution: A simple Eulerian graph  $G$  is Eulerian if it has an Eulerian Trail (i.e. every edge is visited exactly once in the <sup>close</sup> walk).

Therefore, to show that the line graph of a simple Eulerian graph also has a closed walk that visits every edge exactly once, we can start a walk from an edge in

$G$ , and move to an adjacent edge on the line graph if the corresponding edges in  $G$  share a common vertex.

Since  $G$  is a Eulerian graph and the line graph is of  $G$ , then we can construct a closed walk <sup>in the line graph</sup> that covers every edge exactly once. This means that the line graph is also Eulerian. Therefore, the line graph of a simple Eulerian graph is Eulerian.



⊕

(ii) If the line graph of a simple graph  $G$  is Eulerian, must  $G$  be Eulerian?

Solution: Let's assume that the line graph of a simple graph  $G$  is Eulerian. This means we can construct a closed walk in the ~~line~~ line graph that covers every edge exactly once. But there might be vertices along  $G$  that we cannot start and end at and will cover all the edges exactly once.

So,  $G$  may not be Eulerian.

Therefore, if the line graph of a simple graph  $G$  is Eulerian, it does not mean  $G$  is Eulerian.