

Math 300–03

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Set 3

Fall 2008

Outline

- 1 Cartesian Products and Relations**
- 2 Equivalence Relations
- 3 Homework #6
- 4 Partitions
- 5 Last Class ... and Today ...
- 6 Ordering Relations
- 7 Homework #7
- 8 Graphs
- 9 Homework #8

Cartesian Products and Relations

Two of the most important mathematical objects are **relations** and **functions**.

They reflect special relationships between elements of a set A or of pairs of elements of two sets A and B .

Ordered Pair

Definition

Let A and B be sets. For $a \in A$ and $b \in B$, the **ordered pair** (a, b) is the set

$$\{\{a\}, \{a, b\}\}.$$

a is called the first coordinate of the pair, and b the second coordinate.

Note that (a, b) may be different from (b, a) :

$$\{\{a\}, \{a, b\}\} \neq \{\{b\}, \{a, b\}\},$$

if $a \neq b$.

Definition

Let A and B be sets. The set of all ordered pairs having first coordinate in A and second coordinate in B is called the **Cartesian product** of A and B and written $A \times B$. Thus

$$A \times B = \{(a, b) : a \in A \text{ and } b \in B.\}$$

Example: Let $A = \{a, b\}$, $B = \{1, 2, 3\}$. Then:

$$A \times B = \{(a, 1), (a, 2), (a, 3), (b, 1), (b, 2), (b, 3)\}.$$

Theorem

If A and B are finite sets, then

$$\overline{\overline{A \times B}} = \overline{\overline{A}} \cdot \overline{\overline{B}}.$$

Ordered triples, quad...

Ordered triples, quadruples, n -tuples—can also be defined:

Definition

Let A, B and C be sets. For $a \in A, b \in B$ and $c \in c$, the **ordered triple** (a, b, c) is the set

$$\{\{a\}, \{a, b\}, \{a, b, c\}\}.$$

a is called the first coordinate of the pair, b the second coordinate, and c the third coordinate.

The set of all these ordered triples is the cartesian product $A \times B \times C$.

For a finite collection A_1, A_2, \dots, A_n , we may define

$$A_1 \times A_2 \times \cdots \times A_n.$$

Some Tools

Theorem

If A, B, C and D are sets, then

- 1 $A \times (B \cup C) = (A \times B) \cup (A \times C).$
- 2 $A \times (B \cap C) = (A \times B) \cap (A \times C).$
- 3 $A \times \emptyset = \emptyset.$
- 4 $(A \times B) \cap (C \times D) = (A \cap C) \times (B \cap D).$
- 5 $(A \times B) \cup (C \times D) \subseteq (A \cup C) \times (B \cup D).$
- 6 $(A \times B) \cap (B \times A) = (A \cap B) \times (B \cap A).$

Proof. To prove (1), $A \times (B \cup C) = (A \times B) \cup (A \times C)$,

- 1 The ordered pair $(x, y) \in A \times (B \cup C)$
- 2 iff $x \in A$ and $y \in B \cup C$
- 3 iff $x \in A$ and ($y \in B$ or $y \in C$)
- 4 iff $(x \in A$ and $y \in B$) or $(x \in A$ and $y \in C)$
- 5 iff $(x, y) \in A \times B$ or $(x, y) \in A \times C$
- 6 iff $(x, y) \in (A \times B) \cup (A \times C)$,
- 7 Therefore, $A \times (B \cup C) = (A \times B) \cup (A \times C)$.

Relations

Definition

Let A and B be sets. R is a **relation from A to B** iff R is a subset of $A \times B$,

$$R \subset A \times B.$$

If $(a, b) \in R$ we write $a R b$ and say that a is R -related to b . If $(a, b) \notin R$, we write $a \not R b$. A relation R from A to A is called a **relation of A** : $R \subset A \times A$.

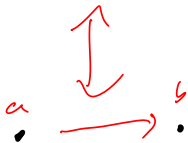
There are many notations for relations: familiar ones are $a \simeq b$, $a \geq b$, $a \mid b$, etc.

Example: I_A (**identity** of A is the relation $a \simeq b$ iff $a = b$. Another: $a R b$ for all $a, b \in A$.)

Graphical Representation

Relation \rightarrow Digraph

$a \ R \ b$



$a \ R \ a$

$a \cdot a$

Example

Solution. Let $A = \{1, 2, 3, 4\}$ and $B = \{-1, 1, 2, 4, 5\}$. We first describe a relation R from A to B with an explicit list of pairs:

$$R = \{(1, 4), (2, 5), (2, -1), (4, 1)\}$$

In table form, we would write R as

1	4
2	5
2	-1
4	1

The same relation R may be written as $R = \{(x, y) \in A \times B : |x - y| = 3\}$. The graph of R is shown in Figure 3.1.

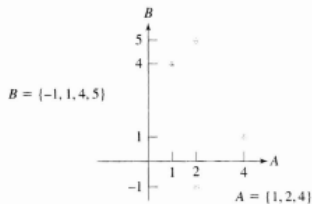


Figure 3.1

Definition

The **domain** of a relation R from A to B is the set

$$\text{Dom}(R) = \{x \in A : \text{there exists } y \in B \text{ such that } x R y\}.$$

The **range** of the relation R is the set

$$\text{Rng}(R) = \{y \in B : \text{there exists } x \in A \text{ such that } x R y\}.$$

Example

Figure 3.3. Let $S = \left\{ (x, y) \in \mathbb{R} \times \mathbb{R} : \frac{x^2}{324} + \frac{y^2}{64} \leq 1 \right\}$. The graph of S is given in:

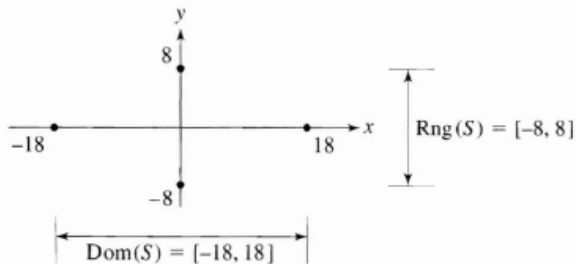
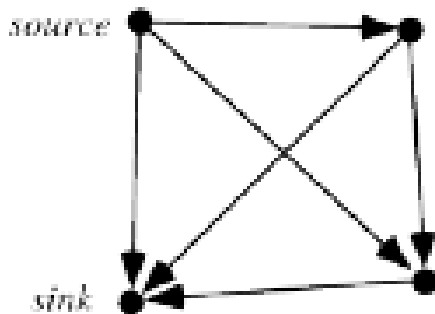


Figure 3.3

From Relations to Graphs



Plenty of Relations

Question: If A is a set with m elements, and B is another set with n elements, how many relations are there from A to B ?

Answer: A lot: Since $A \times B$ has mn elements, and relations are subsets of $A \times B$, each relation is an element of the power set $\mathcal{P}(A \times B)$ of $A \times B$. Thus there are 2^{mn} relations. Thus, if $A = \{a, b\}$, the number of relations of A (that is, from A to A) is $2^{(2)(2)} = 2^4 = 16$.

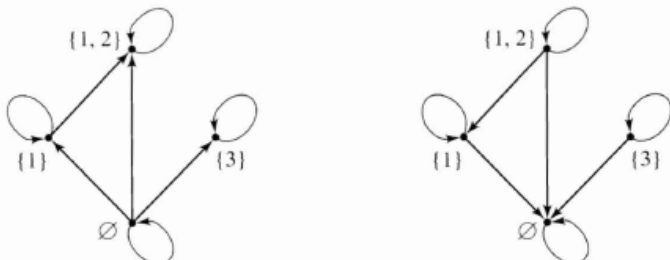
New Relations from Old

Definition

If R is a relation from A to B , the **inverse** of R is

$$R^{-1} = \{(y, x) : (x, y) \in R\}.$$

The digraph of the inverse of a relation on a set differs from the digraph of the relation only in that the directions of the arrows are reversed. Figure 3.8 shows the digraphs of R and R^{-1} , where R is the relation \subseteq on the set $\{\emptyset, \{1\}, \{3\}, \{1, 2\}\}$



Theorem

Let R be a relation from A to B .

- 1 R^{-1} is a relation from B to A .
- 2 $\text{Dom}(R^{-1}) = \text{Rng}(R)$.
- 3 $\text{Rng}(R^{-1}) = \text{Dom}(R)$.

Proof.

- 1 Suppose $(x, y) \in R^{-1}$. Then $(y, x) \in R$. Since R is a relation from A to B , $R \subseteq A \times B$. Thus $y \in A$ and $x \in B$. Therefore $(y, x) \in B \times A$, which proves $R^{-1} \subseteq B \times A$.
- 2 $y \in \text{Dom}(R^{-1})$ iff there exists $x \in A$ such that $(x, y) \in R$ iff $x \in \text{Rng}(R)$.
- 3 Same argument as (2).

New Relations from Old

Definition

Let R be a relation from A to B , and let S be a relation from B to C . the **composite** of R and S is

$$S \circ R = \{(a, c) : \text{there exists } b \in B \text{ such that } (a, b) \in R \text{ and } (b, c) \in S\}.$$

Confusing Diagram...

Let $A = \{1, 2, 3, 4, 5\}$, and $B = \{p, q, r, s, t\}$, and $C = \{x, y, z, w\}$.
Let R be the relation from A to B :

$$R = \{(1, p), (1, q), (2, q), (3, r), (4, s)\}$$

and S the relation from B to C :

$$S = \{(p, x), (q, x), (q, y), (s, z), (t, z)\}$$

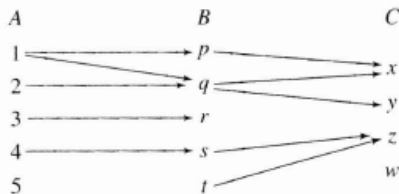


Figure 3.10

Theorem

Suppose A, B, C and D are sets. Let R be a relation from A to B , S a relation from B to C , and T a relation from C to D :

$$A \xrightarrow{R} B \xrightarrow{S} C \xrightarrow{T} D.$$

- 1 $(R^{-1})^{-1} = R.$
- 2 $T \circ (S \circ R) = (T \circ S) \circ R.$
- 3 $I_B \circ R = R$ and $R \circ I_A = R.$
- 4 $(S \circ R)^{-1} = R^{-1} \circ S^{-1}.$

Proof of (2): Note both $T \circ (S \circ R)$ and $(T \circ S) \circ R$ are relations from A to D , that is they are subsets of $A \times D$.

To prove $T \circ (S \circ R) = (T \circ S) \circ R$, let $a T \circ (S \circ R) d$. Note that $S \circ R$ is a relation from A to C .

- 1 Thus there is $c \in C$ such that $a S \circ R c$ and $c T d$.
- 2 Hence there is $b \in B$ such that $b S c$.
- 3 Therefore $b T \circ S d$.
- 4 Since $a R b$ and $b T \circ S d$, it follows that $a (T \circ S) \circ R d$.
- 5 This shows that $T \circ (S \circ R) \subseteq (T \circ S) \circ R$. The reverse inequality has a similar proof.

Last Class ... and Today ...

Relation

A **relation** from a set A to a set B is a subset

$$R \subset A \times B.$$

If $(a, b) \in R$, we also write

$$a R b \quad a \overset{R}{\rightarrow} b \quad a \rightarrow b$$

The two extreme examples are: $R = \{(a, a) : a \in A\}$,
 $R = \{(a, b) : a, b \in A\}$. The first is the **identity** of A .

4 Examples

Let \mathbb{N} be the set of natural numbers. Define

$$a \rightarrow_1 b \quad a \text{ divides } b$$

$$a \rightarrow_2 b \quad a = b + 1$$

$$a \rightarrow_3 b \quad a \text{ does not divide } b$$

$$a \rightarrow_4 b \quad a \text{ and } b \text{ are prime numbers and } a = b + 2$$

Goldbach conjecture

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Equivalence Relations

Definition

Let A be a set and R a relation on A .

- R is **reflexive** iff for all $x \in A$, $x R x$.
- R is **symmetric** iff for all $x \in A$ and $y \in A$, if $x R y$, then $y R x$.
- R is **transitive** iff for all x, y and z in A , if $x R y$ and $y R z$, then $x R z$.

Example

Let R be the set of all $(x, y) \in \mathbb{N} \times \mathbb{N}$ such that $x + y$ is divisible by 3. Is this relation symmetric? reflexive? transitive?

- 1 If $(x, y) \in R$, $x + y = 3m$, for some m . Then $y + x = 3m$ also, so $(y, x) \in R$: **Symmetric**
- 2 $(1, 1) \notin R$: **Not Reflexive**
- 3 $(1, 2), (2, 1) \in R$ but $(1, 1) \notin R$: **Not Transitive**

Example

Let R be the set of all $(x, y) \in \mathbb{N} \times \mathbb{N}$ such that $x + y^2$ is divisible by 2. Is this relation symmetric? reflexive? transitive?

- 1 If $(x, y) \in R$, $x + y^2 = 2m$, for some m , that is $x + y^2$ is even. If x is even then y^2 must be even so y must be even, while if x is odd then both y^2 and y must be odd. Then $y + x^2 = 2n$ also, so $(y, x) \in R$: **Symmetric**
- 2 $(x, x) \in R$ since $x + x^2 = x(x + 1)$ is even: **Reflexive**
- 3 If $(x, y), (y, z) \in R$ we have: If $(x, y) \in R$, if x is even (odd), y is also even (odd), so if $(y, z) \in R$, x is also even (odd). Thus x, z are both even or both odd, so $(x, z) \in R$: **Transitive**

Equivalence Relation

Definition

A relation R on a set A is an **equivalence relation on A** iff R is reflexive, symmetric, and transitive.

Let R be a relation on the set A .

- **R Reflexive:** $a \rightarrow a \forall a \in A$
- **R Symmetric:** $a \rightarrow b \Rightarrow b \rightarrow a$
- **R Transitive:** $a \rightarrow b$ and $b \rightarrow c \Rightarrow a \rightarrow c$

Exercise

Problem: If $A = \{a, b\}$, we saw that there are 16 relations of A . How many of these are equivalence relations? List them all. What if $A = \{a, b, c\}$, a set with $2^9 = 512$ relations, how many of these are equivalence relations?

One Volunteer:

Equivalence Class

Definition

Let R be an equivalence relation on the set A . For $x \in A$, the **equivalence class** of x determined by R is the set

$$x/R = \{y \in A : x R y\}.$$

This is read “the class of x modulo R .” The set of all equivalence classes of R is called A **modulo** R and denoted $A/R = \{x/R : x \in A\}$.

Example: Two integers have the same **parity** if they are both even or both odd. Let

$R = \{(x, y) \in \mathbb{Z} \times \mathbb{Z} : x \text{ and } y \text{ have the same parity.}\}$ R is an equivalence relation with two equivalence classes: the even integers E and the odd integers D . $\mathbb{Z}/R = \{E, D\}$.

Big Example

Let m be a fixed, nonzero integer. Let \equiv_m be the relation on \mathbb{Z} ,

$$x \equiv_m y \text{ iff } m \text{ divides } x - y.$$

This is also written $x \equiv y \pmod{m}$ or even $x = y \pmod{m}$.

It is easy to see that $\mathbb{Z}/\equiv_2 = \{E, D\}$. This set is also denoted by \mathbb{Z}_2 and called the set of integers modulo 2. For $m = 3$, \equiv_3 is also an equivalence relation and there are three distinct equivalence classes.

Theorem

The relation \equiv_m is an equivalence relation on the integers. The set of equivalence relations is called \mathbb{Z}_m and has m distinct elements $\overline{0}, \overline{1}, \overline{2}, \dots, \overline{m-1}$.

Proof

We first prove that \equiv_m is an equivalence relation. Observe that $x \equiv_m y$ means that $x - y = am$, for some integer a .

- 1 **reflexive:** $x \equiv_m x$: $x - x = 0 \cdot m$.
- 2 **symmetric:** $x \equiv_m y$: $x - y = a \cdot m$ for some $a \in \mathbb{Z}$. Thus $y - x = (-a)m$, therefore $y \equiv_m x$.
- 3 **transitive:** If $x \equiv_m y$ and $y \equiv_m z$, $x - y = am$ and $y - z = bm$ for a and b in \mathbb{Z} . Then

$$x - z = (x - y) + (y - z) = am + bm = (a + b)m.$$

Therefore $x \equiv_m z$.

Now we determine the equivalence classes of \equiv_m .

- 1 The numbers $0, 1, 2, \dots, m - 1$ lie in different equivalence classes: For any pair of them, i and j , we cannot have $i \equiv_m j$ since $i - j$ cannot be divisible by m .
- 2 If x is an integer, by the Euclidean algorithm,

$$x = qm + r, \quad 0 \leq r < m.$$

Thus, $x \equiv_m r$. Therefore $r \in \mathbb{Z} / \equiv_m$ (usually denoted by \bar{r}).

- 3 Therefore

$$\mathbb{Z}_m = \{\bar{0}, \bar{1}, \bar{2}, \dots, \overline{m-1}\}.$$

Bonus

The \mathbb{Z}_2 , the set made up by two elements $\{0, 1\}$ (or (even, odd)) with addition defined by the table

+	0	1
0	0	1
1	1	0

 $1 + 1 = 0!$

and multiplication by

×	0	1
0	0	0
1	0	1

More Bonus

The \mathbb{Z}_3 , the set made up by three elements $\{\bar{0}, \bar{1}, \bar{2}\}$ (or more simply, $\{0, 1, 2\}$) has similar properties. For instance, with addition defined by the table

+	0	1	2	1 + 2 = 0!
0	0	1	2	
1	1	2	0	
2	2	0	1	

and multiplication by

×	0	1	2	2 × 2 = 1!
0	0	0	0	
1	0	1	2	
2	0	2	1	

Exercise

Let A be the set of all vectors of the plane. We look at A as the set of all pairs (a, b) of real numbers. The usual notation for A is \mathbb{R}^2 . Define the following relation of A :

$$(a, b) \simeq (c, d) \quad (a, b) - (c, d) = (r, r) \text{ for some } r \in \mathbb{R}$$

This means that the two vectors (a, b) and (c, d) differ by a vector along the diagonal of the first quadrant.

- 1 Prove that \simeq is an **equivalence relation** of A .
- 2 Prove that every **equivalence class** contains exactly a vertical vector (that is, a vector of the form $(0, c)$).

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- 1 3.1: 6(g), 8(d), 9(a), 12, 13(c), 16
- 2 3.2: 2(f), 4(b), 5(b), 8, 10(a), 13
- 3 How many equivalence relations are there on the set $\{a, b, c, d, e\}$? Explain [This is a typical exam question]

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Partitions

Definition

Let A be a nonempty set. A **partition** of A is a set \mathcal{A} of subsets of A such that

- 1 If $X \in \mathcal{A}$, then $X \neq \emptyset$.
- 2 If $X \in \mathcal{A}$ and $Y \in \mathcal{A}$, $X \neq Y$, then $X \cap Y = \emptyset$.
- 3 $\bigcup_{X \in \mathcal{A}} X = A$.

How do partitions arise? It will be a pretty straight answer.

Example

Let $A = \{a, b, c\}$ —the following are partitions of A :

- $\{\{a, b, c\}\}$
- $\{\{a\}, \{b\}, \{c\}\}$,
- $\{\{a\}, \{b, c\}\}$,
- $\{\{b\}, \{a, c\}\}$
- $\{\{c\}, \{a, b\}\}$

Are they all?

Partitions versus Equivalence Classes

Theorem

Let \mathcal{B} be a partition of the nonempty set A . For x and y in A , define $x Q y$ iff there exists $C \in \mathcal{B}$ such that $x \in C$ and $y \in C$. Then

- 1 Q is an equivalence relation on A .
- 2 $A/Q = \mathcal{B}$.

Example: Define the following sets of \mathbb{Z} :

$$A_0 = \{3k : k \in \mathbb{Z}\} = \{\dots, -6, -3, 0, 3, 6, \dots\}$$

$$A_1 = \{3k + 1 : k \in \mathbb{Z}\} = \{\dots, -5, -2, 1, 4, 7, \dots\}$$

$$A_2 = \{3k + 2 : k \in \mathbb{Z}\} = \{\dots, -4, -1, 2, 5, 8, \dots\}$$

The partition defines the relation we denoted \equiv_3 .

How Partitions arise

Theorem

Let R be an equivalence relation on a nonempty set A . Then

- ① For all $x \in A$, $x/R \subseteq A$ and $x \in x/R$. (Thus $x/R \neq \emptyset$.)
- ② $\bigcup_{x \in A} x/R = A$.
- ③ $x R y$ iff $x/R = y/R$.
- ④ $x \not R y$ iff $x/R \cap y/R = \emptyset$.

Thus, the set $\{x/R : x \in A\}$ of equivalence classes is a partition of A .

In words: An equivalence relation on the set A gives rise to a partition of A and vice-versa.

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Last Class ... and Today ...

- Relations
- Equivalence Relation
- Partitions and Equivalent Classes
- Order Relations

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Ordering Relations

Definition

A relation R on a set A is **antisymmetric** if, for all $x, y \in A$, if $x R y$ and $y R x$, then $x = y$:

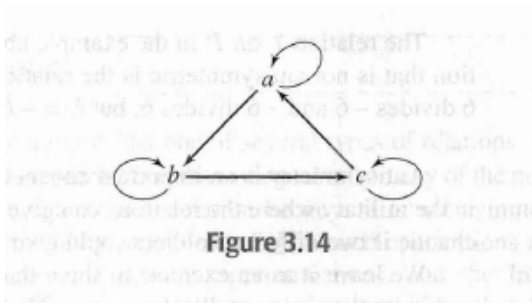
$$x \rightarrow y \rightarrow x \quad \Rightarrow \quad x = y.$$

Definition

A relation R on a set A is a **partial order** (or **partial ordering**) for A if R is reflexive on A , antisymmetric and transitive. A set A with a partial order R is called a **partially ordered set**, or **poset**.

Typical Digraph

Observe the antisymmetry and transitivity:



Top Example

Let X be a set and let $\mathcal{P}(X)$ be its power set. If A and B are in $\mathcal{P}(X)$, i.e. A and B are subsets of X , define the relation

$$A R B \text{ iff } A \subset B.$$

Let us check for the **reflexive**, **antisymmetric** and **transitive** properties of a relation:

- 1 $(A, A) \in R$ since $A \subset A$;
- 2 If $(A, B) \in R$ and $(B, A) \in R$ then $A \subset B$ and $B \subset A$ and therefore $A = B$ by the definition of equality of sets (same elements);
- 3 If $(A, B) \in R$ and $(B, C) \in R$ then $A \subset C$ and therefore $(A, C) \in R$.

Example: Divisors

Let M be a positive integer and let A be the set

$$A = \{n \in \mathbb{N} : n \mid M.\}$$

Consider the relation D 'divides'

For example, for $M = 12$, $A = \{1, 2, 3, 4, 6, 12\}$. Then D consists of the pairs (a, b) where a divides b . For example, $(2, 6) \in D$ but not $(4, 6)$.

Theorem

If R is a partial order for a set A and $x R x_1, x_1 R x_2, x_2 R x_3, \dots, x_n R x$,

$$x_1 \rightarrow x_2 \rightarrow x_3 \rightarrow \dots \rightarrow x_n \rightarrow x_1,$$

then $x = x_1 = x_2 = x_3 = \dots = x_n$.

This means:

$$x \rightarrow x_1 \rightarrow x_2 \rightarrow x \Rightarrow x \rightarrow x_2 \rightarrow x \Rightarrow x = x_2 \Rightarrow x \rightarrow x_1 \rightarrow x$$

and

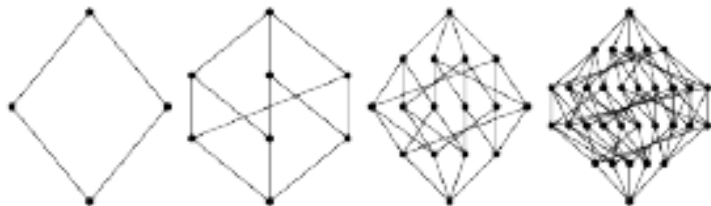
$$x \rightarrow x_1 \rightarrow x \Rightarrow x = x_1$$

Therefore $x = x_1 = x_2$.

Proof. By induction on n .

- 1 For $n = 1$: If $x R x_1$ and $x_1 R x$, by antisymmetry $x = x_1$.
- 2 Suppose that whenever $x R x_1, x_1 R x_2, x_2 R x_3, \dots, x_k R x$, then $x = x_1 = x_2 = x_3 = \dots = x_k$ for some natural number k , and suppose $x R x_1, x_1 R x_2, x_2 R x_3, \dots, x_k R x_{k+1}, x_{k+1} R x$. By transitivity applied to $x_k R x_{k+1}, x_{k+1} R x$, we have $x_k R x$. Now we use the induction hypothesis to deduce $x = x_1 = x_2 = \dots = x_k$. Since $x_k = x$, we have $x R x_{k+1}$ and $x_{k+1} R x$, so $x = x_{k+1}$. Therefore $x = x_1 = x_2 = x_3 = \dots = x_{k+1}$.

Digraphs



Some Terminology

Definition

Let R be a partial ordering on a set A , and let $a, b \in A$ with $a \neq b$. Then a is an **immediate predecessor** of b if $a R b$ and there does not exist $c \in A$ such that $a \neq c$, $b \neq c$, $a R c$ and $c R b$.

Example: For $X = \{1, 2, 3, 4, 5\}$, partially order $\mathcal{P}(X)$ by set inclusion \subseteq . For $b = \{2, 3, 5\}$, there are 3 immediate predecessors: $\{2, 3\}$, $\{3, 5\}$ and $\{2, 5\}$.

Example

Let $M = \{1, 2, 3, 5, 6, 10, 15, 30\}$ be the set of divisors of 30, and let D be the relation “divides”.

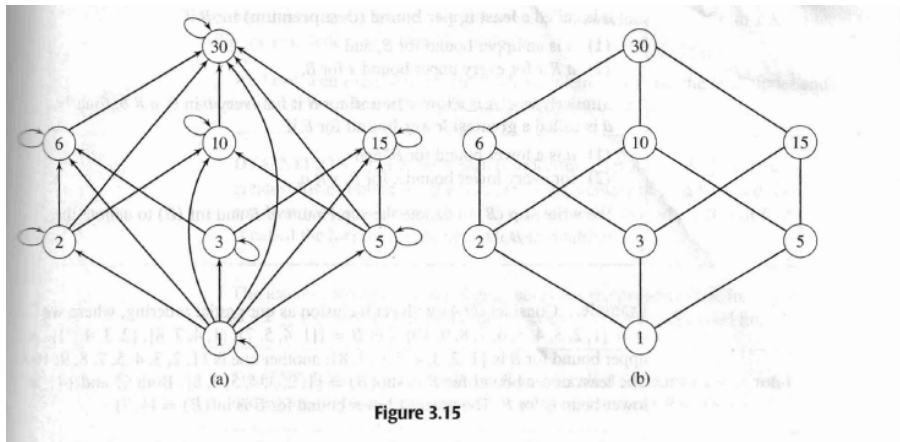


Figure 3.15

Example

Let $A = \mathcal{P}(\{1, 2, 3\})$, ordered by \subseteq :

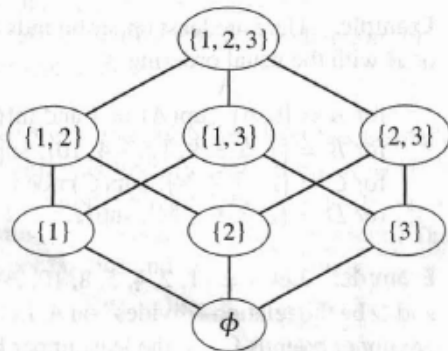


Figure 3.16

More Terminology

Definition

Let R be a partial order for A and let B be a subset of A . Then $a \in A$ is an **upper bound** for B if for every $b \in B$, $b R a$. Also, a is called a **least upper bound** (or **supremum**) for B if

- 1 a is an upper bound for B , and
- 2 $a R x$ for every upper bound x for B .

Similarly, $a \in A$ is an **lower bound** for B if for every $b \in B$, $a R b$. Also, a is called a **greatest lower bound** (or **infimum**) for B if

- 1 a is a lower bound for B , and
- 2 for every lower bound x for B , $x R a$.

In notation: $\sup(B)$ will denote the supremum of B , and $\inf(B)$ will denote the infimum.

Example

Examples: Subsets of \mathbb{R} with the relation \leq :

- $A = [0, 4)$: $\text{Sup}(A) = 4$ and $\text{Inf}(A) = 0$.
- $B = \{2^k : k \in \mathbb{N}\}$: $\text{Sup}(B)$ does not exist and $\text{Inf}(B) = 2$.
- $C = \{x : x \in \mathbb{Q}, x^2 < 2\}$: $\text{Sup}(C) = \sqrt{2}$ and $\text{Inf}(C)$ does not exist.
- $D =$ the set of roots of $x^2 - 3x + 2 = 0$: ??
- $E =$ numbers in $(0, 1)$ which are not fractions (not $p/q, p \neq 0, q \in \mathbb{N}$): ??

More Terminology

Definition

Let R be a partial order for a set A . Let $B \subseteq A$. If the greatest lower bound for B exists and is an element of B , it is called the **smallest** (or least) **element** of B . If the least upper bound for B is in B , it is called the **largest** (or greatest) **element** of B .

Definition

A partial ordering R on A is called a **linear order** (or **total order**) on A if for any two elements x and y of A , either $x R y$ or $y R x$, that is

$$x \rightarrow_R y \quad \text{or} \quad y \rightarrow_R x.$$

Definition

Let L be a linear ordering on a set A . L is a **well ordering** on A if every nonempty subset B of A contains a smallest element.

Examples

- 1 WOP: The Well-Ordering Principle says that \mathbb{N} , with the ordering given by \leq , is a linear ordering with property above.
- 2 Let A be a partial ordered set, called “the alphabet.” Let W be the set of all “words” of length two—that is, combinations of two letters of the alphabet. Define a relation \leq on W as follows: for $x_1x_2 \in W$ and $y_1y_2 \in W$, $x_1x_2 \leq y_1y_2$ iff (i) $x_1 \leq x_2$ or (ii) $x_1 = x_2$ and $y_1 \leq y_2$. We claim that W is a partial ordering of W (called lexicographic ordering, as in a dictionary).

Example

Let $A = \{a, b, c\}$. Given an example of a relation on A that is

- ① antisymmetric and symmetric

$$\{a R a, b R b\}$$

- ② symmetric and not antisymmetric

$$\{a R b, b R a, c R c\}$$

- ③ antisymmetric, reflexive on A and not symmetric

$$\{a R a, b R b, c R c, a R b\}$$

- 1 antisymmetric, not reflexive on A and not symmetric

$$\{a R b, b R c, a R c\}$$

Exercise

Existence: Let $S = \{(x, y) \in \mathbb{R} \times \mathbb{R} : x = 1 - y\}$. Is S antisymmetric?

Need to show: If $(x, y), (y, x) \in S$ then $(x, x) \in S$

$(1, 0) \in S$ and $(0, 1) \in S$ but $(0, 0) \notin S$.

$x = 1 - y \Rightarrow y = 1 - x$ therefore S is symmetric.

Upper and Lower Bounds

- The number b is said to be an **upper bound** of the set $A \subset \mathbb{R}$ if

$$a \leq b \mid \forall a \in A$$

- A number ℓ is said to be a **lower bound** of the set $A \subset \mathbb{R}$ if

$$a \geq \ell \mid \forall a \in A$$

- Consider the set $A = \{q \in \mathbb{Q} \mid q^2 < 2\}$. -2 is a lower bound of A , while $3/2$ is an upper bound. Clearly there are many other bounds.

Least Upper and Greatest Lower Bounds

- A number b is said to be a **least upper bound** of the set $A \subset \mathbb{R}$ if b is an upper bound of A and $b \leq b'$ for any other upper bound b' . Least upper bounds are also known as the **supremum** of A . If $b \in A$, it is called the **maximum** of A .

$$A = \{x_1 = 1, \forall n \quad x_{n+1} = \frac{x_n}{2} + 1\}$$

has 2 for supremum [needs a proof, as we only proved that 2 is an upper bound]

- Similarly we define **greatest lower bound** [and of **infimum/minimum**].

Example

Define the set $\mathbf{A} = \{a_1, a_2, a_3, \dots\}$ by the rule

$$a_1 = \sqrt{2}, \quad a_2 = \sqrt{2\sqrt{2}}, \quad a_3 = \sqrt{2\sqrt{2\sqrt{2}}}, \dots$$

Let us show that $\sup \mathbf{A} = 2$:

$$a_1 = \sqrt{2}, \quad a_2 = a_1 \sqrt[4]{2}, \quad a_3 = a_2 \sqrt[8]{2}, \dots$$

$$a_n = 2^{1/2+1/4+\dots+1/2^n} < 2$$

$$a_n = 2^r, \quad r = \frac{1/2 - 1/2^{n+1}}{1/2} = 1 - 1/2^n$$

Axiom of Completeness

Axiom: Every set A of real numbers with an upper bound has a least upper bound.

This is a defining property of \mathbb{R} . A lot flows out of it. For example consider the set A of all rational numbers x such that $x^2 < 2$. This set has an **upper bound** (in fact many). For instance, $x \leq 3$. The **axiom of completeness** guarantees that there is a **real** number α such that

$$\alpha^2 = 2.$$

Outline

- 1 Cartesian Products and Relations
- 2 Equivalence Relations
- 3 Homework #6
- 4 Partitions
- 5 Last Class ... and Today ...
- 6 Ordering Relations
- 7 Homework #7**
- 8 Graphs
- 9 Homework #8

Homework #7

- 1 3.3: 3(a,b,c), 7(b), 9, 12
- 2 3.4: 2(a,b), 3(all), 9, 12, 13(b), 20
- 3 Let D be a positive integer and let A be the set

$$A = \{n \in \mathbb{N} : n \mid D.\}$$

Prove that A is linearly ordered iff D is the power of a prime.

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Graphs

Definition

A **graph** G is a pair (V, E) , where V is a nonempty set and E is a set of unordered pairs of distinct elements of V .

An element of V is called a **vertex** and an element of E is called an **edge**. An edge between the vertices u and v is written **uv** rather than as the set $\{u, v\}$.

Terminology

Definition

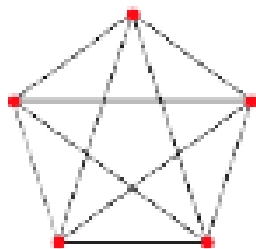
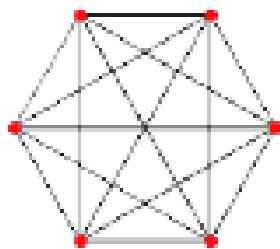
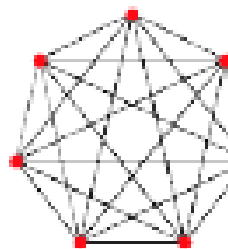
Let $G = (V, E)$ be a graph. The **order** of the graph G is the number of vertices. The **size** of the graph G is the number of edges.

Vertices u and v are **adjacent** if $uv \in E$; the edge uv is said to be **incident** with u and v .

The **degree** of a vertex u is the number of edges incident with u .

A graph G of order n , may have no edges, all the vertices are isolated. Such a graph is called the **null graph**.

A graph such that every pair of vertices are adjacent is called a **complete graph**. If it has n vertices, its size is $\binom{n}{2}$. Its notation; K_n .

K_n  K_5  K_6  K_7

Exercises

If possible, give an example of a graph

- 1 with order 6 and size 6
- 2 with order 4 and size 6
- 3 with order 3 and size 6
- 4 with order 6 and size 3

Properties of Graphs

Theorem

- (The Handshaking Lemma) For each graph G , the sum of the degrees of the vertices of G is even.
- For every graph G the number of vertices of G having odd degree is even.

Proof.

- Each edge is incident with two vertices. Thus

$$\sum_{v \in V} \deg(v) = 2 \times (\text{number of edges}).$$

- Let V_1 be the subset of vertices of odd degree and V_2 the subset of vertices of even degree.

$$\underbrace{\sum_{v \in V} \deg(v)}_{\text{even number}} = \sum_{v \in V_1} \deg(v) + \underbrace{\sum_{v \in V_2} \deg(v)}_{\text{even number}}.$$

Therefore the term

$$\sum_{v \in V_1} \deg(v)$$

being a difference of two even numbers is even. Since each term $\deg(v)$, $v \in V_1$, is odd, we must have an even number of them for the sum to be even. □

Walks

Definition

A **walk** (or **path**) in a graph G is a finite sequence of vertices

$$v_0, v_1, v_2, \dots, v_m$$

where each $v_i v_{i+1}$ is an edge in G . The vertex v_0 is the **initial vertex** and v_m is the **terminal vertex**. The length of the walk is m , the number of edges. If $v_0 = v_m$ the walk is **closed**. A **path** in G is a walk where all the vertices, except for possibly the initial and terminal vertices, are distinct.

The walk $v_0, v_1, v_2, \dots, v_m$ is said to **traverse** the vertices of the sequence.

Theorem

Let G be a graph of order n .

- 1 If there is a walk originating at v and terminating at u , then there is a path from v to u .
- 2 The length of a path in G that is not closed is at most $n - 1$. The length of a closed path is at most n .

Proof. Let v, v_1, v_2, \dots, u be a walk with $u \neq v$. If the walk is not a path, some vertex v_i appears twice in the sequence. Let x be the first such vertex. Then the walk contains at least a closed walk of the form x, v_j, \dots, v_m, x . Delete the vertices v_j, \dots, v_m, x . If the result is a path, we are done. Otherwise another repeated vertex occurs and we repeat the process until all duplications are deleted. The process will result in a path.

Proof Cont'd

In case $v = u$ the process produces a closed path [where v is the only repeated vertex].

The walks that arise at the end are:

$$v \rightarrow v_1 \rightarrow v_2 \rightarrow \cdots \rightarrow v_m \rightarrow u, \quad v \neq u$$

$$v \rightarrow v_1 \rightarrow v_2 \rightarrow \cdots \rightarrow v_m \rightarrow u, \quad v = u$$

Note that the length is the number of arrows. Counting the \rightarrow gives the assertions: at most $n - 1$ in the first case, and at most n in the second.

Reachability

Definition

Let G be a graph and u a vertex of G . The vertex v is **reachable** (or **accessible**) from u if there is a path from u to v . The number of edges in a path of minimum length from u to v is called the **distance from u to v** , written $d(u, v)$. We say that any vertex u is reachable from itself and $d(u, u) = 0$.

Connected Components

Definition

Let G be a graph. If u is a vertex of G , the **component containing u** is $C(u) = \{v \in V : v \text{ is reachable from } u\}$.

Theorem

Let G be a graph with vertex set V and let R be the relation defined by $v R u$ if u is reachable from v . Then R is an equivalence relation on V and the set of equivalence classes for R is $\{C(u) : u \in V\}$.

Definition

A graph G is **connected** iff every vertex is reachable from every other vertex. G is **disconnected** if it is not connected.

Exercise

Give an example of a graph G with 6 vertices having

- 1 One component
- 2 two components
- 3 three components
- 4 six components

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Homework #8

- 1 3.5: 7, 9, 11
- 2 4.1: 2(a,b,c,d), 3(e), 9, 12(a,b), 16(b), 17(a)