

## A minimum of category theory in two pages

*Category:* A class of objects  $obj(\mathcal{C})$ , disjoint sets of morphisms  $Mor_{\mathcal{C}}(A, B)$  between objects  $A, B$  and associative composition  $gf$  of  $f \in Mor_{\mathcal{C}}(A, B), g \in Mor_{\mathcal{C}}(B, C)$  and distinguished morphism  $1_A \in Mor_{\mathcal{C}}(A, A)$  such that  $f1_A = f, 1_Bg = g$ .

*Functor* from category  $\mathcal{C}$  to category  $\mathcal{D}$  : an assignment to each object  $A$  of  $\mathcal{C}$  and object  $F(A)$  and to each morphism  $f \in Mor_{\mathcal{C}}(A, B)$  a morphism  $F(f) \in Mor_{\mathcal{D}}(F(A), F(B))$  such that  $F(gf) = F(g)F(f)$ .

A morphism  $f \in Mor_{\mathcal{C}}(A, B)$  is an *isomorphism* if it has an inverse under composition, an *epimorphism* if  $gf = hf$  implies  $g = h$  and a *monomorphism* if  $fg = fh$  implies  $g = h$  (corresponding to isomorphism, surjection and injection in the categories  $Set, Grp, R - Mod$ )

An *initial object* of a category  $\mathcal{C}$  is an object  $A$  such that  $Mor_{\mathcal{C}}(A, B)$  is a one element set for all objects  $B$ , a *terminal object* is an object  $Z$  such that  $Mor_{\mathcal{C}}(B, Z)$  is a one element set for all objects  $B$ , and a *zero object* is an object which is initial and terminal.

A *natural transformation*  $\eta$  of functors  $F_1, F_2$  from category  $\mathcal{C}$  to  $\mathcal{D}$  is an assignment to each object  $A$  of  $\mathcal{C}$  a morphism  $\eta_A \in Mor_{\mathcal{D}}(F_1(A), F_2(A))$  such that for each morphism  $f \in Mor_{\mathcal{C}}$ ,  $\eta_B \circ F_1(f) = F_2(f) \circ \eta_A$ . A *natural isomorphism* of functors is a natural transformation that has a two sided inverse natural transformation.

Categories  $\mathcal{C}, \mathcal{D}$  are *isomorphic* if there exist functors  $F$  from  $\mathcal{C}$  to  $\mathcal{D}$  and  $G$  from  $\mathcal{D}$  to  $\mathcal{C}$  such that  $FG$  and  $GF$  equal the identity functors on  $\mathcal{D}, \mathcal{C}$  respectively. If we replace the word equal by the terms “are naturally isomorphic to” we say the categories are *equivalent*.

*Diagram* in category  $\mathcal{C}$ : A functor  $D$  from index category  $\mathcal{I}$  to  $\mathcal{C}$ , for example the constant diagram which sends all objects in the index category to the same object  $N$  in  $\mathcal{C}$  and all morphisms to the identity morphism  $1_N$ .

A *cone* over a diagram  $D$  is a natural transformation from a constant diagram to  $D$  (*cocone* : to a constant diagram from  $D$ )

*Limit of a diagram*  $D$  is a terminal object in category of cones over  $D$ , a *colimit* is an initial object in the category of cocones over  $D$ .

Examples: If  $D$  is a category with two objects and no morphisms between them a cocone is a diagram with maps from objects  $A, B$  to  $Y$  and an initial object is an object called the coproduct which has morphisms from  $A, B$  and for any other  $Y$  having morphisms from  $A$  and  $B$  they are obtained via a map from the coproduct to  $Y$ . In the category of modules this is the direct sum  $A \oplus B$ . In other words, to give a pair of maps from  $A, B$  to an object is the same as giving a map of the coproduct to the object. The limit of this diagram in a category is the product  $A \times B$  which has maps to  $A, B$  and any other map of an object  $Y$  to  $A$  and  $B$  factors through a map of  $Y \rightarrow A \times B$ .

*Adjoint functors:* We say  $F : \mathcal{C} \rightarrow \mathcal{D}$  and  $G : \mathcal{D} \rightarrow \mathcal{C}$  are adjoint if  $\text{Mor}_{\mathcal{D}}(FC, D) = \text{Mor}_{\mathcal{C}}(C, GD)$  and there is a natural isomorphism of  $\text{Mor}_{\mathcal{D}}(F-, -)$  and  $\text{Mor}_{\mathcal{C}}(-, G-)$  as functors from  $\mathcal{C}^{op} \times \mathcal{D}$  to SET. We say  $F$  is *left adjoint to*  $G$  and  $G$  is *right adjoint to*  $F$ . Note left adjoints preserve terminal objects, right adjoints preserve initial objects.

Theorem: Left adjoints preserve colimits, right adjoints preserve limits:

Proof: Let  $\text{CoL}$  be a colimit of a diagram  $D$ . Apply  $F$  to the diagram  $D$  and the cocone of it with vertex  $\text{CoL}$ . This is a cocone over the diagram  $FD$  with vertex  $F(\text{CoL})$ . Given any other cocone over  $FD$  with vertex  $Y$ , by the adjoint property we have elements of  $\text{Mor}(FD_i, Y) = \text{Mor}(D_i, GY)$ . Thus the object  $GY$  is the vertex of a cocone for  $D$ , hence there is a unique map of the cone with vertex  $\text{CoL}$  to  $GY$ , and by adjointness a unique morphism of  $F(\text{CoL})$  to  $Y$ . Hence  $F(\text{CoL})$  is a colimit for the diagram  $FD$ . The proof for colimits is similar except for reversing the morphisms appropriately.