

## Homework 6-solutions

MATH 361

(due November 4)

October 28 2022

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**Problem 1.** Prove that for any  $r, s \in \mathbb{R}$ ,  $r \cdot s \in \mathbb{R}$ . (You can use problem 6 from HW5).

**Solution.** The product is defined by cases:

If  $x, y \geq 0$  then  $x \cdot y = 0 \cup \{p \cdot q \mid p \in x, q \in y, p, q \geq 0\}$ . If  $x, y < 0$  then  $x \cdot y = |x| \cdot |y|$ . If  $x < 0 \leq y$  or  $y < 0 \leq x$  then  $x \cdot y = -(|x| \cdot |y|)$

Once we have proven case (1), then since  $|x|, |y|$  are non negative,  $|x| \cdot |y|$  will be in  $\mathbb{R}$  and by HW5 problem 6, also  $-(|x| \cdot |y|)$  and we will be done. So let us prove case (1). Assume that  $x, y \geq 0$  and then  $x \cdot y = 0 \cup \{p \cdot q \mid p \in x, q \in y, p, q \geq 0\}$ . It is clearly non-empty since for example  $-1 \in x \cdot y$ . Also, if  $p^*, q^* > 0$  bound  $x, y$  respectively, then for every  $z \in x \cdot y$ , either  $z < 0 < p^*q^*$ , or  $z = p \cdot q$  for some  $p, q \geq 0$  and  $p \in x$  and  $q \in y$ . It follows that  $p < p^*$  and  $q < q^*$  which in turn implies (since we are dealing with positive rationals) that  $p \cdot q < p^* \cdot q^*$ . Hence  $p^* \cdot q^*$  bounds  $x \cdot y$ . To see it is downward closed, let  $t < z \in x \cdot y$ . if  $t \leq 0$  then it is clearly in  $x \cdot y$ . Otherwise,  $0 < t \leq z$  and therefore  $z = p \cdot q$  for some  $p, q > 0$  rationals. where  $p \in x$  and  $q \in y$ . Let  $q' = \frac{z}{p}$ . Then  $z = p \cdot q'$  and since  $p \cdot q' \leq p \cdot q$ , we have that  $q' \leq q$  (since  $q, q', p$  are all positive). Since  $y$  is a Dedekind cut,  $q' \in y$  and therefore  $z = p \cdot q' \in x \cdot y$ . Finally we need to prove that  $x \cdot y$  has no last element. Let  $q \in x \cdot y$ . If  $q < 0$  then  $q < \frac{q}{2} < 0$  hence  $\frac{q}{2} \in x \cdot y$ . If  $q \geq 0$ , then there are  $0 \leq p_1, p_2, p_1 \in x$  and  $p_2 \in y$  such that  $q = p_1 \cdot p_2$ . Since  $x, y$  are dedekinf cuts, there are  $p_1 < p'_1 \in x$  and  $p_2 < p'_2 \in y$ . Since they are all positive,  $p_1 \cdot p_2 < p'_1 \cdot p'_2 \in x \cdot y$  as wanted.

**Problem 2.** For each of the following statements provide an appropriate function (no need to prove that your functions satisfy the required proper-

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ties):

1.  $\mathbb{R} \approx \mathbb{R} \setminus \{0\}$ .

**Solution.**  $f : \mathbb{R} \rightarrow \mathbb{R} \setminus \{0\}$

$$f(x) = \begin{cases} x - 1 & x \in \mathbb{N}^+ \\ x & o.w. \end{cases}$$

2.  $\mathbb{Z} \approx \mathbb{N}_{\text{even}} \setminus \{0, \dots, 2023\}$ .

**Solution.** Define  $f : \mathbb{Z} \rightarrow \mathbb{N}_{\text{even}} \setminus \{0, \dots, 2023\}$  by

$$f(n) = \begin{cases} 2024 + 4|n| & n \leq 0 \\ 2022 + 4n & n > 0 \end{cases}$$

3.  $\mathbb{N} \times \mathbb{N} \leq^{\mathbb{N}} \{0, 1\}$

**Solution.** We saw in class one example  $f(\langle n, m \rangle) = 00000\dots \underset{n^{\text{th}} \text{ place}}{1} 0\dots 0 \underset{n+m^{\text{th}} \text{ place}}{1} 0\dots$

Formally,  $f(\langle n, m \rangle) : \mathbb{N} \rightarrow \{0, 1\}$  is defined

$$f(\langle n, m \rangle)(k) = \begin{cases} 1 & k \in \{n, n + m\} \\ 0 & o.w. \end{cases}$$

4.  $\{f \in {}^{\mathbb{R}}\mathbb{R} \mid \exists i \in \{0, 1\}, \forall x \in \mathbb{R} \setminus \mathbb{Q}, f(x) = i\} \approx \{0, 1\} \times {}^{\mathbb{Q}}\mathbb{R}$ .

**Solution.** Denote the set by  $A$ .  $F : A \rightarrow \{0, 1\} \times {}^{\mathbb{Q}}\mathbb{R}$  defined by

$$F(f) = \langle f(\sqrt{2}), f \upharpoonright \mathbb{Q} \rangle$$

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**Problem 3.** Prove that

$$\left\{ X \in P(\mathbb{N}) \mid \mathbb{N}_{\text{even}} \subseteq X \right\} \approx P(\mathbb{N})$$

[Hint: First find a function from  $P(\mathbb{N})$  to  $P(\mathbb{N}_{\text{odd}})$ ]

**Solution.** Like in the proof that  $A \approx B \Rightarrow P(A) \approx P(B)$  we fix a bijection  $f : \mathbb{N} \rightarrow \mathbb{N}_{\text{odd}}$ , for example  $f(n) = 2n + 1$ , then  $F(X) = f''X$  is a bijection from  $P(\mathbb{N})$  to  $P(\mathbb{N}_{\text{odd}})$ . Explicitly,  $F(X) = \{2n + 1 \mid n \in X\}$ . Now define  $G : P(\mathbb{N}) \rightarrow \{X \in P(\mathbb{N}) \mid \mathbb{N}_{\text{even}} \subseteq X\}$  defined by  $G(X) = \mathbb{N}_{\text{even}} \cup \{2n + 1 \mid n \in X\}$ . Check that this is a bijection.

**Problem 4.** Let  $C(\mathbb{R})$  be the set of all continuous function  $f : \mathbb{R} \rightarrow \mathbb{R}$ . Prove that

$$C(\mathbb{R}) \leq \mathbb{Q}\mathbb{R}$$

[Hint: use that fact that  $\mathbb{Q}$  is dense in  $\mathbb{R}$  to prove that the restriction function  $G : C(\mathbb{R}) \rightarrow \mathbb{Q}\mathbb{R}$  defined by  $G(f) = f \upharpoonright \mathbb{Q}$  is one-to-one.]

**Solution.** Let  $G : C(\mathbb{R}) \rightarrow \mathbb{Q}\mathbb{R}$  defined by  $G(f) = f \upharpoonright \mathbb{Q}$ , let us prove that it is one-to-one. Suppose that  $f, g$  are two continuous functions, such that  $f \upharpoonright \mathbb{Q} = g \upharpoonright \mathbb{Q}$ . We need to prove  $f = g$ . Let  $x \in \mathbb{R}$ , by density of the rationals we can find a sequence  $(q_n)_{n=0}^{\infty}$  of rationals, such that  $\lim_{n \rightarrow \infty} q_n = x$ , then for each  $n$ ,  $f(q_n) = g(q_n)$  (since  $f \upharpoonright \mathbb{Q} = g \upharpoonright \mathbb{Q}$ ). By continuity,

$$f(x) = \lim_{n \rightarrow \infty} f(q_n) = \lim_{n \rightarrow \infty} g(q_n) = g(x)$$

**Problem 5.** Prove that if  $A \approx B$  and  $C \approx D$  then  $A \times C \approx B \times D$ .

**Solution.** Let  $f : A \rightarrow B$  and  $g : C \rightarrow D$  be bijections. Define  $h : A \times C \rightarrow B \times D$   $h(\langle a, c \rangle) = \langle f(a), g(c) \rangle$ . Prove that  $h$  is one-to-one. Let us prove for

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example that  $h$  is onto. Let  $\langle b, d \rangle \in B \times D$ . Since  $f, g$  are onto, there are  $a \in A$  and  $c \in C$  such that  $f(a) = b$  and  $g(c) = d$ . Then  $\langle a, c \rangle \in A \times C$  and  $h(\langle a, c \rangle) = \langle f(a), g(c) \rangle = \langle b, d \rangle$ .

**Problem 6.** Prove that for every  $\alpha < \beta$  real numbers  $(\alpha, \beta) \approx (0, 1)$ . [Hint: First stretch/shrink  $(0, 1)$  to have length  $\beta - \alpha$ , then shift it by  $+c$  as we did in class.]

**Solution.** Define  $f : (0, 1) \rightarrow (\alpha, \beta)$  by  $f(x) = (\beta - \alpha)x + \alpha$ . It is not hard to check that  $f$  is one-to-one and onto.

### Additional problems

**Problem 7.** Show that  $x \cdot (y + z) = x \cdot y + x \cdot z$  for every  $x, y, z \in \mathbb{R}$ .

**Problem 8.** Show that for every  $n > 0$ ,  $\mathbb{N}^n \approx \mathbb{N}$ . [Hint: Induction. you can  $\mathbb{N} \times \mathbb{N} \approx \mathbb{N}$ .]

**Problem 9.** Show that  ${}^{\mathbb{N}}\{0, 1\} \times {}^{\mathbb{N}}\{0, 1\} \approx {}^{\mathbb{N}}\{0, 1\}$ . [Hint: see HW2 problem 5.]