

Factoring Polynomials over Finite Fields

First Square roots mod p.

Given odd prime p , $a \in \underline{\mathbb{F}_p}$.

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$\{0, 1, \dots, p-1\}$
with operations
mod p .

Goal: Find $x \in \mathbb{F}_p$ s.t. $x^2 = a \pmod{p}$.

Ideal running time: poly (log p) time.

This is open!

$$\mathbb{F}_p = \{0, 1, \dots, p-1\}$$

$$\mathbb{F}_p^* = \{1, 2, \dots, p-1\}$$

a group under multiplication.

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$$\frac{2}{p-1}$$

Equivalently, $\exists g \in \mathbb{F}_p^*$ s.t.

$$\{1, g, g^2, \dots, g^{p-2}\} = \mathbb{F}_p^*$$

$$\{\text{Nonzero Perfect Squares}\} = \{g^{2i} : 0 \leq i \leq p-2\}$$

$$\downarrow = \{g^{2i} : 0 \leq i \leq \lfloor \frac{p-2}{2} \rfloor\}$$

$\frac{p-1}{2}$ of them.

Test for whether $a \in \mathbb{F}_p^*$ is a perfect square

Non-squares \rightarrow Consider $b = a^{\frac{p-1}{2}}$.
If $a = g^{\text{odd}}$, $b = g^{\frac{p-1}{2}} = -1$

Squares \rightarrow If $a = g^{\text{even}}$, $b = a^0 = 1$

How do we find $a^{\frac{p-1}{2}}$?

Repeated Squaring

First find $a, a^2, a^4, a^8, \dots, a^{2^{\lceil \log p \rceil}}$.
by $a^{2^i} = (a^{2^{i-1}})^2$

Then find $a^{\frac{p-1}{2}} = \prod_{\substack{j \\ \text{certain } j\text{'s}}} a^{2^j}$.

Whole thing takes $\text{poly}(\log p)$ time.

If $p \equiv 3 \pmod{4}$,

$$\left(a^{\frac{p+1}{4}}\right)^2 = a^{\frac{p+1}{2}}$$

$$= a^{\frac{p-1}{2}} \cdot a$$

$$= (+1) \cdot a \rightarrow \text{If } a \text{ is a square}$$

$$\text{So } a^{\frac{p+1}{4}} \equiv \sqrt{a} \pmod{p}.$$

If $p \equiv 5 \pmod{8}$, algorithm to find square roots in $\text{poly}(\log(p))$ time

There is an algorithm that runs in time $\text{poly}(a, \log p)$ that finds $\sqrt{a} \pmod{p}$

Schoof, via Elliptic curves.

Today Randomized $\text{poly}(\log p)$ time algorithm for square roots \pmod{p} .

An algorithm $F(a, p, y)$, where

p prime, $a \in \mathbb{F}_p$, $y \in \{0, 1\}^{\text{poly} \log p}$.

s.t.

$\forall p, \forall a \in \mathbb{F}_p$ a perfect square,

P_n
 $y \in \{0, 1\}^{\text{poly} \log p} \left[F(a, p, y) \neq \begin{array}{l} \text{a square} \\ \text{root} \\ \text{of } a \end{array} \right]$
in \mathbb{F}_p

≤ 0.1 .

Consider $Q(x) = x^{\frac{p-1}{2}} - 1$

Dist-
recting $\left[\begin{array}{l} 0. \text{ Roots of } Q(x) \text{ are the } \frac{(p-1)}{2} \text{ perfect} \\ \text{squares in } \mathbb{F}_p. \end{array} \right.$

Input-
cent $\left[\begin{array}{l} 1. Q(x) \text{ has } \Omega(1) \text{ and } 1 - \Omega(1) \\ \text{fraction roots in } \mathbb{F}_p. \\ \\ 2. Q(x) \text{ is sparse.} \end{array} \right.$

Consider $f(x) = x^2 - a$.

$$= (x-d)(x+d)$$

where $d = \sqrt{a}$

Idea 1: Compute $\text{GCD}(Q(x), f(x))$

If it is degree 1, then
we found a nontrivial factor
of $f(x)$, hooray!

Idea 2 Randomize the roots of
 $f(x)$ by looking at

$\tilde{f}(x) = f(cx+d)$ for random
 $c, d \in \mathbb{F}_p$.

If $\alpha, \beta \in \mathbb{F}_p$ $\alpha \neq \beta$.

then for uniformly random $c, d \in \mathbb{F}_p$,
 $(c\alpha + d, c\beta + d)$ is uniformly
 distributed in \mathbb{F}_p^2 .

$$\begin{bmatrix} \alpha & 1 \\ \beta & 1 \end{bmatrix} \begin{pmatrix} c \\ d \end{pmatrix}$$

η
invertible

So the chance that exactly one of
 $\{c\alpha + d, c\beta + d\}$ ~~is~~ is a
 root of $Q(x)$
 is $= \frac{1}{2} + O\left(\frac{1}{p}\right)$.

To find $\text{GCD}(Q(x), \tilde{f}(x))$

need to find $x^{\frac{p-1}{2}} - 1 \pmod{\tilde{f}(x)}$
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deg 2

First find $x, x^2, x^4, \dots, x^{2^f} \pmod{\tilde{f}(x)}$.

as deg 1 polys mod $\tilde{f}(x)$.

Find $x^{2^f} \pmod{\tilde{f}(x)}$ as a product of
 $\leq \log p$ of these.
mod $\tilde{f}(x)$.

Find roots of $f(x)$ reduces
to finding roots of $\tilde{f}(x)$ provided
 $C \neq 0$. (this holds w. prob $1 - \frac{1}{p}$).

Berlekamp's algorithm 1970

Factoring polys over finite
fields.

Above alg. easily generalizes to finding roots in \mathbb{F}_p of any polynomial. (odd prime p).

Also to \mathbb{F}_q for odd prime powers q .

Also to \mathbb{F}_2 for even prime powers 2 .

(FACT = In \mathbb{F}_{2^m} , the polynomial

$$\text{Tr}_{\mathbb{F}_{2^m}/\mathbb{F}_2}(x) = x + x^2 + x^4 + x^8 + \dots + x^{2^{m-1}}$$

has 2^{m-1} roots in \mathbb{F}_{2^m}

and has only $m = \log_2(2^m)$ monomials).

Next time: Factoring.

