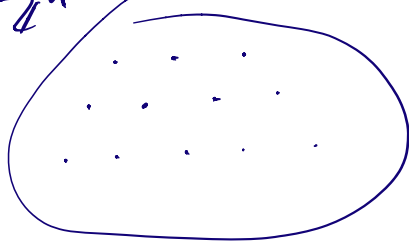


Algorithmic Number Theory

Some problems

1. Algorithms for arithmetic
+, \times , other stuff
representations
2. Primality testing
3. Factoring integers
4. Discrete logs
5. Polynomial analogues of all these
6. Lattices eg \mathbb{Z}^n

7. Recognizing algebraic numbers.
8. Elliptic Curves

9 Cryptography and motivations
From there

10. Hilbert's 10th problem.

Undecidability of solvability of
Diophantine equations.

11. Quantum Algorithms

Problem sets

Project at the end 10-15 page write up
on a recent research paper.

Generalities

Represent integers in base 2.

1. Can add, multiply, divide, subtract,
convert bases in time
poly(n) where the inputs
are n-bits long.



0.

2. Rational numbers, represented as
 a/b $a, b \in \mathbb{Z}$.

"Complexity" of $a/b \stackrel{p}{=} \log_2(a) + \log_2(b)$

2 problems

1. GCDs: computing GCD of 2 numbers
 2. Recognizing approximate rational numbers.
-

Naive algorithm for $\text{GCD}(a, b)$

- ① factor a, b
- ② use prime factorization.

Factoring is difficult!

- Trial division requires time $\approx \left(\frac{n}{2}\right)$

Euclid's GCD algorithm

Observation: $\text{GCD}(a, b) = \text{GCD}(a-b, b)$
 $= \text{GCD}(a \bmod b, b)$

Alg (a, b):

Make $a > b$ by swapping a, b if necessary.

If $b = 0$

return a

else

return Alg (a-b, b)

Alg (a mod b, b)

$a > b$

$\log a + \log b$ keeps reducing by

vs

~~log a~~

$\log(a \bmod b) + \log b$

Claim: $(a \bmod b) \cdot b < \frac{a \cdot b}{2}$

If $a > 2b$, then yes.

$$\text{If } a = (1+d) \cdot b$$

$$\text{then } a \bmod b = 2b < \left(\frac{1+d}{2}\right)b.$$

So $\text{Alg}(a, b)$, if a, b are n bit long integers, runs in $\leq 2n$ recursions.

Runtime of $\text{Alg} \leq O(n^3)$.

Bezout Identity

$$\text{If } \text{GCD}(a, b) = d,$$

then $\exists x, y \in \mathbb{Z}$ s.t.

$$ax + by = d.$$

Can also find x, y like this using Euclid's algorithm.

$$\dots d \dots \dots \dots 0 \dots$$

$$a = q_1 b + \underline{r_1}$$

$$b = q_2 r_1 + \underline{r_2}$$

$$r_1 = q_3 r_2 + \underline{r_3}$$

$$r_{s-1} = q_{s+1} r_s + \underline{d} (= r_{s+1}) \leftarrow$$

$$r_s = q_{s+2} \cdot \underset{r_{s+1}}{d} + 0$$

This gives us d as integer combination of a, b .



Given relatively prime a, b ,

can find x st. $ax \equiv 1 \pmod{b}$.

[Find x, y s.t. $ax + by = 1$.
Then $ax \equiv 1 \pmod{b}$]

continued FRACTIONS

$$1.73 = a_0 + \frac{1}{a_1 + \frac{1}{a_2 + \frac{1}{a_3 + \dots}}}$$

where all $a_i \in \mathbb{Z}$

$$a_i > 0.$$

$$a_0 = 1$$

$$\frac{1}{0.73} = a_1 + \frac{1}{\dots}$$

$$1.3 = 1 + \frac{1}{\dots}$$

$$\frac{1}{0.3} = a_2 + \frac{1}{\dots}$$

$$3.1 = 3 + \frac{1}{\dots}$$

$$1.73 = 1 + \frac{1}{1 + \frac{1}{3 + \frac{1}{\dots}}}$$

Problem Given rational number $\frac{a}{b}$,
and a denominator bound Q ,
where a, b, Q are all $\leq n$ bits long.
Find p, q s.t.

$\left| \frac{a}{b} - \frac{p}{q} \right|$ is minimized over all
rationals $\frac{p}{q}$ with $q \leq Q$.