

Intro to Mathematical Reasoning (Math 300)–Honors  
Assignment 9 <sup>1</sup>

1. (30 points) Say that a relation  $R$  on  $A$  is a TR relation if it is both transitive and reflexive. (TR relation is not a standard term.) We know that a TR relation that is symmetric is an equivalence relation, and a TR relation that is anti-symmetric is a partial order. The purpose of this problem is to give a general characterization of TR relations. We'll prove: A relation  $R$  on  $A$  is a TR relation if and only if there is a set  $B$  with a partial order relation  $Q$ , and a function  $f : A \rightarrow B$  such that for all  $x, y \in A$ ,  $xRy$  if and only if  $f(x)Qf(y)$ .

This is an “if and only if” theorem so there are two directions of the implication to prove:  $\Rightarrow$  and  $\Leftarrow$ . The  $\Leftarrow$  direction is easy and is the first part of the problem. The rest of the problem is to prove the  $\Rightarrow$  direction.

(Warm up problem, not to be handed in: Prove the theorem in the case that (a)  $R$  is an equivalence relation, and (b)  $R$  is a partial order relation.)

- (a) Suppose that  $Q$  is a partial order on  $B$  and that  $f : A \rightarrow B$ . Define the relation  $R$  on  $A$  by  $xRy$  if  $f(x)Qf(y)$ . Prove that  $R$  is a TR relation.
  - (b) For the rest of the problem, suppose  $R$  is an arbitrary TR relation on  $A$ . Define the relation  $W$  on  $A$  by  $xWy$  if and only if  $xRy$  and  $yRx$ . Prove that  $W$  is an equivalence relation.
  - (c) Let  $\mathcal{C}$  denote the set of equivalence classes of  $W$ . Define a relation  $P$  on the set  $\mathcal{C}$  where  $CPD$  if there exists an  $x \in C$  and a  $y \in D$  such that  $xRy$ . Prove that in fact, for all  $C, D \in \mathcal{C}$  if  $CPD$  then for all  $x \in C$  and  $y \in D$ ,  $xRy$ .
  - (d) Prove that  $P$  is a partial order on  $\mathcal{C}$ .
  - (e) Finish the proof of the  $\Rightarrow$  direction.
2. Recall that if  $R$  is a relation on  $A$ , an  $R$ -chain is a list  $(a_1, \dots, a_k)$  such that for each  $i \in \{2, \dots, k\}$ ,  $a_{i-1}Ra_i$ . An  $R$ -chain is said to be *non-repeating* if all of the entries of the chain are different. An  $R$ -chain is said to be a *Hamilton path* (named after the mathematician William Rowan Hamilton) provided that it is non-repeating and every element of  $A$  appears on this list. (Notice that this requires that  $A$  is finite.) Recall that a relation  $R$  on  $A$  is *full* if for all  $x, y \in A$  if  $x \neq y$  then  $xRy$  or  $yRx$ . Prove the following neat theorem: For any finite set  $A$  and full relation  $R$  on  $A$ , there is a Hamilton path in  $R$ . (A hint is given below.<sup>2</sup>)
3. Define a relation  $R$  on  $\mathbb{Z}_{>0}$  as follows: For  $m, n \in \mathbb{Z}_{>0}$ , we have  $mRn$  provided that there are odd numbers  $a$  and  $b$  so that  $ma = nb$ .
- (a) Prove that this relation is transitive, symmetric and reflexive, and is therefore an equivalence relation.
  - (b) Let  $S$  be the set consisting of the smallest member from each  $R$ -equivalence class. Determine explicitly what the set  $S$  is (and prove that your answer is correct.)
4. (20 points) For a relation  $R$  on  $A$ , a subset  $X$  of  $A$  is  $R$ -independent provided that for all  $x, y \in X$ , if  $x \neq y$  then  $x \not R y$ . Also, if  $A$  is finite, define  $c(R)$  to be the length of the largest  $R$ -chain.

The purpose of this problem is to prove the following theorem: Suppose  $A$  is a finite partially ordered set with partial order  $R$ . Then there exists a partition of  $A$  into exactly  $c(R)$  sets, each of which is  $R$ -independent. (As a warm-up, test out the theorem on a few small partially ordered sets.)

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<sup>2</sup>Let  $(a_1, \dots, a_t)$  be an  $R$ -chain of longest length that is non-repeating. Assume for contradiction that not all members of  $A$  appear in this  $R$ -chain, and derive a contradiction.

- (a) Define the function  $h : A \rightarrow \mathbb{Z}_{>0}$  as follows:  $h(x)$  is the length of the longest chain that ends with  $x$ . Prove that for all  $x \in A$ ,  $1 \leq h(x) \leq c(R)$ .
- (b) Let  $(a_1, \dots, a_{c(R)})$  be an  $R$ -chain of length  $c(R)$ . Prove that for  $i \in \{1, \dots, c(R)\}$ ,  $h(a_i) = i$ . (Hint: Prove separately that  $h(a_i) \geq i$  and  $h(a_i) \leq i$ .)
- (c) Prove that for each  $i \in \{1, \dots, c(R)\}$ ,  $h^{-1}(i)$  is an  $R$ -independent set. (Hint: Use proof by contradiction.)
- (d) Finish the proof of the theorem.