## Quiz 2 Solutions

## **Question 1** Consider the following graph *G*:



1. Is G bipartite? If so, provide a partition of its vertices into appropriate sets; if not, rigorously explain why not.

**Answer** No. A graph is bipartite if and only if it contains no odd cycles, and G contains an odd cycle.

2. Define the *diameter* of a graph and find the diameter of G.

**Answer** The *diameter* of G is the maximum value of d(x, y) for x, y in G; for G this value is 3.

3. Find the quantities  $\delta(G)$  and  $\Delta(G)$ .

**Answer** These values indicate the minimum and maximum degree, respectively, so  $\delta(G) = 1$  and  $\Delta(G) = 4$ .

4. Give the order and size of G.

**Answer** The order is 7 and the size is 8.

**Directions** Answer AT LEAST ONE of questions 2 and 3. If you attempt to answer both, clearly indicate which one should be graded first: if you earn at least 80% credit on this one then points for the other will count for extra credit in the Quizzes category.

\*\*\*

**Question 2** Prove or disprove that each of the following relations is an equivalence relation:

- On the set  $\mathbb{Z}$  of integers,  $m \sim n$  if and only if m and n are either both odd or both even (*i.e.*, iff they have the same parity)
- On the set V(G) of vertices in an arbitrary graph G,  $u \sim v$  if and only if u and v are NOT adjacent.

## Answer

- For any integer  $m, m \sim m$  trivially. We have  $m \sim n \rightarrow n \sim m$ ; it doesn't matter what order the integers are given in. Finally, to prove transitivity, consider that two integers m and n have the same parity if and only if m n is even. So if  $m \sim n$  and  $n \sim p$  then m n is even and n p is even, so m p is even and thus  $m \sim p$ . Because it satisfies all three prongs,  $\sim$  is an equivalence relation.
- No vertex can be adjacent to itself in a graph, and if  $u \not\sim v$  then  $v \not\sim u$ . But the relation is not transitive: for example, consider the following graph:



We have  $u \not\sim v$  and  $v \not\sim w$  in the graph, yet  $u \sim w$ .

**Question 3** Recall that the graph G[S] is the induced subgraph of G on the subset S of V(G). Prove that the vertices of every connected graph of order n can be labelled in an order

 $\{v_1, ..., v_n\}$ 

such that the induced graph

 $G[v_1, \ldots, v_k]$ 

is connected for every k with  $1 \le k \le n$ .

**Answer** We simultaneously prove two things by induction: (1) that we can "greedily" choose vertices at each step and (2) that this choice results in a connected induced graph for all k.

(BASE CASE.) Let  $v_1$  be arbitrary in G. Then  $G[v_1]$  is trivial and thus connected.

(INDUCTION STEP.) For  $k \geq 1$  such that k < n, assume the theorem to be true. Then there exists some  $v_{k+1} \in G \setminus \{v_1, ..., v_k\}$  that is adjacent to a vertex in  $G[v_1, ..., v_k]$ . Once we prove this claim, we will have proven that we can find a  $v_{k+1}$  such that  $G[v_1, ..., v_k]$  is connected and our proof will be complete.

The proof of the claim uses the connectedness of G. Indeed, take a point x in the induced graph and a point y outside it and draw an x - y path: because y is not in the induced graph, this path contains an edge between a vertex in the induced graph and a vertex outside.