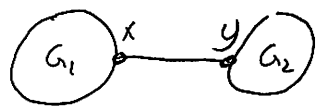


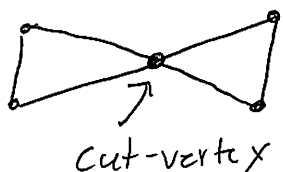
Hw #6: Solutions

- 1) a- Assume $e=xy$ is a bridge in G . Then $G-e$ has two comp G_1 & G_2 as follows:

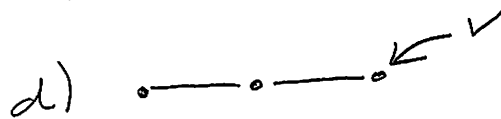


In G_1 , x now has odd degree (as e has been deleted). Moreover all the other vertices in G_1 are even. This is impossible b/c no graph can have an odd # of odd degrees.

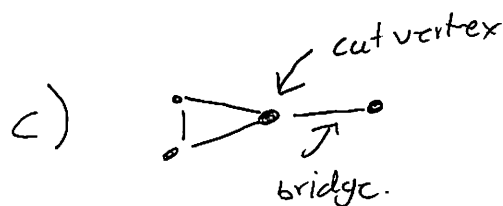
b -



- 2) a) same ex as #1b



- 3) a) same ex as #1b

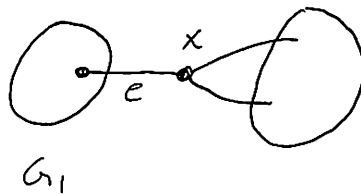


4) ($\Rightarrow$) As G is 3-regular, its order is ≥ 2 .

Therefore thm 7.1 tells us that if $e=xy$ is a bridge then either x or y is a cut-vertex.

($\Leftarrow$) Assume G has a cut-vertex, x . Then $G-x$ has at least two components G_1, G_2 .

As x is adj to at least one vertex in each comp and $\deg x = 3$ we see that x must be adj to exactly one vertex in G_1 (wlog).



It is now clear that e is a bridge.

5) Observe that if T is any tree and we delete a leaf, u , the result is still cntd. If T happens to be a spanning tree for some graph G , then since any $T-u \subseteq G-u$, and $T-u$ is cntd it follows that $G-u$ is cntd as well.

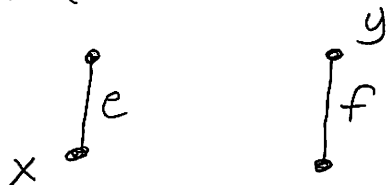
$\therefore u$ is not a cut-vertex.

If G is nontrivial, then any spanning tree for G will also be nontrivial, and hence have at least two leaves. By what we just proved these leaves cannot be cut-vertices.

$\therefore G$ has at least two non-cut-vertices.

6) To prove the reverse direction, we may assume our graph has at least two edges so $n \geq 3$. (Note the only cntd graph w/ 1 edge is K_2 which is nonseparable.)

Let $x, y \in V(G)$, and choose two edges e, f w/ $e \neq f$ so that e is incident to x and f is incident to y : (this is possible as $n \geq 3$).



By assumption, $e + f$ lie on a common cycle. It now follows that x, y lie on a common cycle. By thm 7.2 we can conclude that G is nonseparable.