

Solutions: HW 4

1) Consider the graph



order = ~~4~~, size = 3.

2) Let F be a forest w/ components

$T_1, \dots, T_k$

where T_i is a tree w/ order n_i + size $n_i - 1$.

so

$$\begin{aligned} |E(F)| &= \sum_{i=1}^k |E(T_i)| = \sum_{i=1}^k n_i - 1 = \left(\sum_{i=1}^k n_i \right) - k \\ &= |V(F)| - k. \end{aligned}$$

3) Let G have order n . As $G-u$ is a tree:

$$\begin{aligned} |E(G)| - \deg u &= |E(G-u)| = |V(G-u)| - 1 \\ &= n - 1 - 1 = n - 2. \end{aligned}$$

we also know $G-v$ is a tree:

$$|E(G)| - \deg v = |E(G-v)| = |V(G-v)| - 1 = n - 2.$$

$$\therefore \deg u = \deg v$$

4) Let P & Q be longest paths in G .

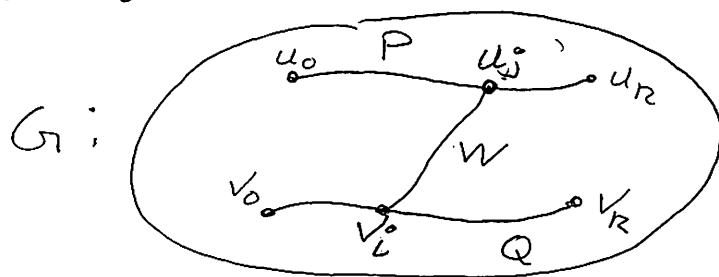
$$P = (u_0 \dots u_k)$$

$$Q = (v_0 \dots v_k)$$

Assume for a contradiction that P & Q have NO vertices in common. As G is cnt'd we must have a $v_i - u_j$ path:

$$R = (w_0 \dots w_m)$$

where $w_0 = v_i$ & $w_m = u_j$. SO



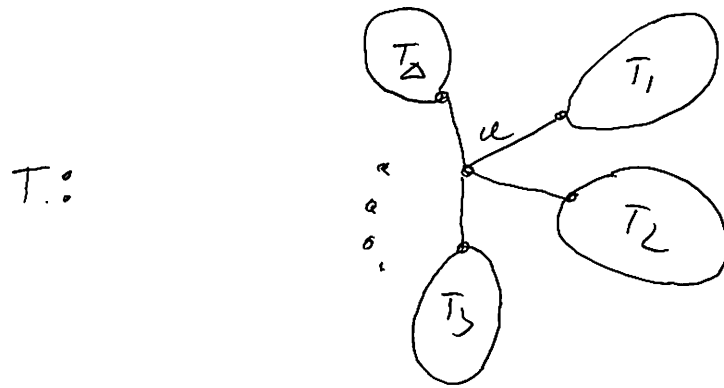
Moreover we can choose W so only its end pts have vertices in common w/ P & Q . If $j \geq k/2$ & $i \geq k/2$, then the path

$$\underbrace{u_0 \dots u_j}_{\geq k/2} \cdot \underbrace{w_1 \dots w_{m-1}}_{\geq 1} \cdot \underbrace{v_{i-1} \dots v_0}_{\geq k/2}$$

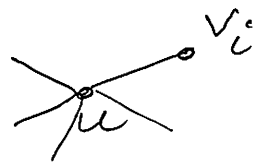
has length $\geq k+1 \Rightarrow \text{contradiction}$. The other cases are similar, and lead to paths $\geq k+1$.

$\therefore P$ & Q must have some vertex in common.

5) Let u be a vertex in T w/ $\deg u = \Delta$.
 so T looks like:



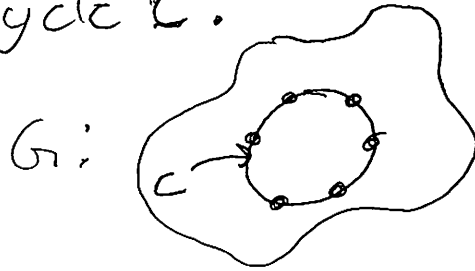
So $T_1, \dots, T_{\Delta}$ are the cntd components of $T - u$. They must be trees. Either T_i is trivial, in which case T_i is just a single vertex v_i and then



v_i is a leaf in T . Otherwise T_i is not trivial and by thm T_i has at least 2 leaves: l_1 & l_2 . It follows that in T either l_1 or l_2 is still a leaf. (Possibly one of them is adj. to u in T , in which case it is no longer a leaf in the big tree.)
 We conclude T has at least Δ leaves.

6) a) Assume G does not have a cycle. Then it is a forest, in which case all comp. are trees. If one of its trees is trivial then it's a vertex of $\deg 0 \Rightarrow \Leftarrow$. Otherwise, if one of its trees is non-trivial then it must have a leaf (in fact two of them!) $\Rightarrow \Leftarrow$.
 $\therefore G$ has a cycle.

b) If G has order n , then $G = C_n$.
 To see why, we know from a) that G has a cycle C :



No vertex on C can be adj to a vertex not on C . As G is cont'd this implies the only vertices are those on C . Moreover if we had edges other than those on C then some vertex would have $\deg \geq 3 \Rightarrow \Leftarrow$.

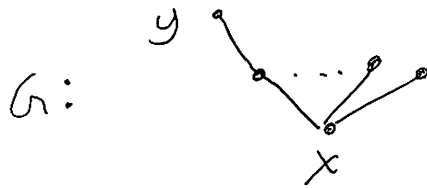
$\therefore G = C_n$.

7) Let P be a path of max length in G . As $\text{diam}(G) = 2$, then

$$P = (u, x, v).$$

Now let $N_G(x)$ be the set of all n'bers of x in G . We claim $|N_G(x)| = n-1$, where $|V(G)| = n$. If $N(x) \neq V(G) \setminus \{x\}$ then as G is cndd $\exists$ some vertex y $\therefore$ y is adj to a n'bor of x but $x \not\sim y$.

Then



which results in a path of length 3 $\Rightarrow \nabla$.

Thus $N(x)$ consists of all the vertices in G except x . Thus x is isolated in $\overline{G}$.

8) Let T be a tree of order n . We know $m := |E(T)| = n-1$.

As $\bar{T}$ is the complement of T we must have

$$|E(\bar{T})| = \binom{n}{2} - |E(T)| \\ = \frac{n(n-1)}{2} - (n-1)$$

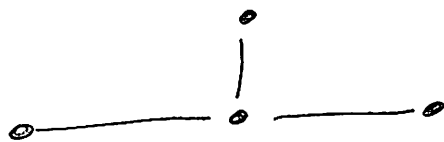
If $\bar{T}$ is a tree, then, as $n = |V(T)| = |V(\bar{T})|$,

$$\frac{n(n-1)}{2} - (n-1) = |E(\bar{T})| = n-1$$

$$\text{so } n(n-1) = 2(n-1) \Rightarrow n^2 - 5n + 4 = 0 \\ (n-4)(n+1) = 0. \\ \therefore n = 4 \quad (n \geq 1).$$

By checking all trees of order 4. we see P_4 is the only one $\rightarrow \bar{P}_4$ is also a tree (In fact $\bar{P}_4 = P_4$!).

The ^{only} other order 4 tree is:



whose comp. is NOT a tree.

8) As T is a tree, then

$$|E(T)| = |V(T)| - 1$$