

Back to S-T


Lists: same as Abn, but \bar{w} bigger t (i.e. better anal)

$$\Gamma t \sim \text{Co}_{\mathbb{Z}_2} d \sim \text{IPA}, \quad s = t^2, \quad \Gamma = [s],$$

t 's unif (indep) $\in \binom{\Gamma}{t} \rightarrow$ whp \neq L-coloring \downarrow

Persp: $X := |\{L\text{-colorings}\}| \rightarrow$ whp $X = 0$?

how ab. $\mathbb{E}X$ small?

NO:  K_{dd}
 fix $\Gamma = \Gamma_X \cup \Gamma_Y$ equipartite

$$\rightarrow \left(\frac{s}{2}\right)^n \text{ potential cols} \rightarrow \mathbb{E}X > \left(\frac{s}{2}\right)^n \left(\frac{t}{s}\right)^n = \left(\frac{t}{2}\right)^n$$

$\rightarrow \text{??}$

ws. (for too small t) $\exists t_\lambda \subseteq \Gamma_Y$ as n.s. $\rightarrow X = 0$

but if not then $X \approx (t/2)^n$ is typ.

[sim (e.g.) threshold for p.m.s.]

Apply container thm \bar{w} arb (fixed, pos) $\approx \rightarrow \mathcal{C}$

\mathcal{C} L-coloring is (I_1, \dots, I_s) \bar{w} I_j indept \dagger
 $v \in I_j \Rightarrow j \in L_v$ ~~\otimes~~

\mathcal{C} L-coloring $\mathcal{C} = (C_1, \dots, C_s)$ \bar{w} $C_j \in \mathcal{C}$ \dagger
 $\forall v \exists j \in L_v \ v \in C_j$ ~~\otimes~~

ETS whp. \nexists L-coloring $(\exists$ L-coloring $\Rightarrow \exists$ L-coloring)

loss: ~~\otimes~~ is easier than ~~\otimes~~

gain: # potential L-colorings $< \exp[\mathbb{R} \frac{n}{d} \log d \cdot s]$
 (vs. possibly $\exp[\Omega(n \cdot s)]$ pot. L-colorings)

Ex. Kdd again, Here can take $\mathcal{C} = \{X, Y\}$ \rightarrow

potential coll's = 2^s (not using thm)

$C = (\underbrace{X, X, \dots, X}_{\Gamma_X}, \underbrace{Y, \dots, Y}_{\Gamma_Y})$ actual coll'g \Leftrightarrow

$x \in X \Rightarrow L_x \not\subseteq \Gamma_Y \ \frac{1}{2} \dots \rightarrow \text{prob} \approx (1 - 2^{-t})^{2d}$

$\sum \mathbb{P}(C \text{ coll'g}) \approx \underbrace{2^{t^2}}_{\text{negl.}} \exp[-2^{-t} 2d]$

$= o(1)$ if $t = \log_2 d - 2 \log_2 \log d$

\rightarrow smaller if $|\Gamma_X| \neq |\Gamma_Y|$
 (as in Erdős-Rubin-Taylor)

(Back to thm)

$$P(X \neq 0) \leq \sum_C P(\underbrace{C \neq \text{good}}_{\text{good}})$$

given $C: F_v = \{j: v \notin C_j\}$

$$\alpha_v = |F_v|/s$$

C good $\iff t_v \notin F_v \quad \forall v$

$$\alpha = n^{-1} \sum \alpha_v = n^{-1} s^{-1} \sum |F_v| = n^{-1} s^{-1} \sum_j (n - |C_j|) > \frac{1}{2} - \varepsilon$$

$$P(C \text{ good}) \approx \prod (1 - \alpha_v^t) < \exp[-\sum \alpha_v^t]$$

Sensen \leftarrow \leq $\exp[-n \alpha^t] < \exp[-n (\frac{1}{2} - \varepsilon)^t]$

$$\sum_C P(C \text{ good}) \leq \exp\left[B \frac{n}{d} \log_2 d \cdot s - n (\frac{1}{2} - \varepsilon)^t\right]$$

check $o(1)$ for some $t = (1 - o(\varepsilon)) \log_2 d$

* larger is α_v small

\Rightarrow

Back to AKS [\leadsto int'g conditioning + more LL, list-col'g, neg. corr.]

recall (from lec 2)

\triangleright AKS '81: G Δ -free, $\bar{d} \leq d \Rightarrow \alpha(G) > c \frac{n \ln d}{d}$

\ominus mot: Sidon sets, $R(\xi, k) = O(k^2 / \log k)$

\ominus pf in 1 line: random greedy

\ominus Shearer '83: $c \sim 1$ ("easy" ind. pf)

can't beat $c \approx 2$ (typ. for $G_{n,p}$)

$\frac{1}{2}$ random greedy does work

\ominus nat: what ab K_r -free ($r > 3$ fixed)?

AKS '81: $\alpha > c_r \frac{n \log \log d}{d}$

(WE DO) \triangleright Shearer '95: $c_r \frac{n \log d}{d \log \log d} \rightarrow$ get rid of this?

(w amazingly simple pf of AKS - will see)

Actually show:

cf. AS p.336

$$\Delta_G \leq d \Rightarrow \bar{\alpha}(G) > \begin{cases} c \frac{n \log d}{d} & \Delta\text{-free} \\ c_r \frac{n \log d}{d \log \log d} & K_r\text{-free } (r \geq 4) \end{cases}$$

\triangleright this is enough (why?)

$$[\text{eq. } \bar{d} \leq d \Rightarrow |\{v: d_v \leq 2d\}| \geq n/2]$$

\triangleright EX: statement for $\alpha \Rightarrow$ statement for $\bar{\alpha}$ (w diff c)

Pf \mathbb{I} unif $\in \mathcal{D}(G)$

$$X_v := d \mathbb{1}_{\{\sum_{w \in \mathbb{I}}\}} + |\mathbb{I} \cap \mathcal{N}_v|$$

MAIN $\forall v$

$$\mathbb{E} X_v = \begin{cases} \Omega(\log d) & r=3 \\ \Omega_r\left(\frac{\log d}{\log \log d}\right) & \text{in gen'l} \end{cases}$$

then done (why?):

$$\sum X_v = d |\mathbb{I}| + \sum_{w \in \mathbb{I}} d_w \leq 2d |\mathbb{I}|$$

$$\rightarrow \mathbb{E} |\mathbb{I}| \geq (2d)^{-1} \sum X_v \quad \text{etc.}$$

□