

GOSTS Descriptive Set Theory

Prewellorderings, Scales, and Uniformization

James Holland

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- Last time we showed one direction of the following equivalences:

Theorem

The following are equivalent:

- ① $\omega_1^{L[x]} < \omega_1$.
- ② $\text{PSP}(\Sigma_2^1(x))$.
- ③ $\text{PSP}(\Pi_1^1(x))$.

- Namely, (1) implies (2), and (2) implies (3).
- To get (3) implies (1), we need concepts related to *prewellorders*,
- Basically, a way to talk about the complexity of functions $f : X \rightarrow \text{Ord}$ for $X \subseteq \mathcal{N}$.
- This motivates the notion of a *scale* as well as bounds on the lengths of prewellorders.

Because we'll be talking a bit in the abstract, it helps to talk about properties common to the bold and lightface pointclasses pointclasses:

Definition

A pointclass $\Gamma \subseteq \mathcal{P}(\mathcal{N})$ is *adequate* iff

- Γ contains all computable relations;
- Γ is closed under computable substitution/preimages;
- Γ is closed under finite unions and intersections; and
- Γ is closed under bounded quantification (over ω).

- So clearly all borel, arithmetical, projective, and analytical pointclasses are adequate.
- Basically, simple operations don't change complexity.

Definition

A *prewellorder* is a relation $\leq$ that is transitive, total, and well-founded.

- This is a *prewellorder* in the following sense.

Result

Let $\leq$ be a prewellorder on X . Define the equivalence relation $x \approx y$ iff $x \leq y \leq x$. Therefore $\leq/\approx$ is a well-order on $\{[x]_{\approx} : x \in X\}$.

- So what do prewellorders look like?
- Basically just well-orders with clusters of loops, but which don't fundamentally change the order.
- So how do we get examples of prewellorders?

Definition

A *prewellorder* is a relation $\leq$ that is transitive, total, and well-founded.

Definition

A *norm* on a set X is a function $\varphi : X \rightarrow \text{Ord}$.

Result (19 C • 5)

Every prewellorder has a norm (its rank function). Moreover, every norm $\varphi : X \rightarrow \text{Ord}$ gives a prewellorder $x \leq y$ iff $\varphi(x) \leq \varphi(y)$.

- Note that the norm associated with a prewellorder isn't unique.
- For example, $\varphi = \{\langle 0, 0 \rangle, \langle 1, 1 \rangle\}$ and $\varphi' = \{\langle 0, 0 \rangle, \langle 1, 4 \rangle\}$ give the same prewellorder.

Result (19 C • 5)

Every prewellorder has a norm (its rank function). Moreover, every norm $\varphi : X \rightarrow \text{Ord}$ gives a prewellorder $x \leqslant y$ iff $\varphi(x) \leqslant \varphi(y)$.

- Also note that this gives lots of trivial prewellorders: any constant function $\varphi : X \rightarrow \{\alpha\}$.
- On the other hand, assuming φ is injective, we get a well-order.
- Any set has lots of prewellorders, even in ZF.
- But we are interested in prewellorders with definability restrictions.

Definition

Let $\Gamma \subseteq \mathcal{P}(\mathcal{N})$ be a pointclass. $X \subseteq \mathcal{N}$ has a Γ -norm iff there's a norm $\varphi : X \rightarrow \text{Ord}$ such that $\leq_\varphi$ and $<_\varphi$ are both in Γ , defined by

$$x \leq_\varphi y \quad \text{iff} \quad x \in X \wedge (y \in X \rightarrow \varphi(x) \leq \varphi(y))$$

$$x <_\varphi y \quad \text{iff} \quad x \in X \wedge (y \in Y \rightarrow \varphi(x) < \varphi(y))$$

- Let's consider the complexity here: $x \leq_\varphi x$ iff $x \in X$ so having a Γ -norm implies $X \in \Gamma$ (if Γ is adequate).
- We know the constant 0 function $\varphi = \text{const}_0$ is a norm on every set, but it may not have the best complexity:
- $\text{const}_0(x) < \text{const}_0(y)$ is always false so that $x <_\varphi y$ iff $x \in X$ and $y \notin X$.
- This requires complexity $\Gamma \wedge \neg\Gamma$.

Corollary

If $\Delta \subseteq \mathcal{P}(\mathcal{N})$ is adequate with $\neg\Delta = \Delta$, then every $X \in \Delta$ has a Δ -norm. E.g. each Δ_n^1 has this.

Definition

$\Gamma \subseteq \mathcal{P}(\mathcal{N})$ has the *prewellordering property*, $\text{PWO}(\Gamma)$, iff every $X \in \Gamma$ has a Γ -norm.

Corollary

- $\text{PWO}(\Delta_\alpha^0)$ and $\text{PWO}(\Delta_n^1)$ for every $\alpha < \omega_1$, $n < \omega$ (and their boldface counterparts).
- $\text{PWO}(\Sigma_0^1)$ (and $\text{PWO}(\Sigma_1^0)$)

Proof.

- $A \in \Sigma_0^1 = \Sigma_1^0$ is $\bigcup_{n \in \omega} \mathcal{N}_{f(n)}$ for some computable f .
- Define $\varphi(x)$ as the least $n \in \omega$ with $x \in \mathcal{N}_{f(n)}$.
- This is a Σ_0^1 -norm: each cone is Δ_1^0

$$x \leq_\varphi y \quad \text{iff} \quad x \in A \wedge \exists n \in \omega (x \in \mathcal{N}_{f(n)} \wedge \forall m < n (y \notin \mathcal{N}_{f(m)}))$$

$$x \leq_\varphi y \quad \text{iff} \quad x \in A \wedge \exists n \in \omega (x \in \mathcal{N}_{f(n)} \wedge \forall m \leq n (y \notin \mathcal{N}_{f(m)})).$$

⊥

Definition

$\Gamma \subseteq \mathcal{P}(\mathcal{N})$ has the *prewellordering property*, $\text{PWO}(\Gamma)$, iff every $X \in \Gamma$ has a Γ -norm.

A much harder pointclass to show PWO for is Π_1^1 .

Theorem

$\text{PWO}(\Pi_1^1)$

To prove this, recall the following nice properties for Π_1^1 -sets.

Lemma

Every $X \in \Pi_1^1$ is the computable preimage of WO.

This is quite nice for us, because we already showed that WO has a norm on it: $x \mapsto \|x\|$ where $\|x\|$ is the length of the well-order coded by $x \in \text{WO}$.

Lemma (19 B • 8 and 19 C • 8)

$\psi(x) = \|x\|$ is a Π_1^1 -norm on WO.

Lemma

Every $X \in \Pi_1^1$ is the computable preimage of WO, and $\psi(x) = \|x\|$ is a Π_1^1 -norm on WO.

Theorem

$\text{PWO}(\Pi_1^1)$

Proof.

As a result, we can define a Π_1^1 -norm for $X \in \Pi_1^1$ as follows:

- Let $X = f^{-1}''\text{WO}$ for $f : \mathcal{N} \rightarrow \mathcal{N}$ computable.
- Define $\varphi = \|f \upharpoonright X\|$. We get that

$$\begin{aligned} x \leq_{\varphi} y &\leftrightarrow f(x) \leq_{\psi} f(y) \\ &\leftrightarrow f(x) \in \text{WO} \wedge (f(y) \in \text{WO} \rightarrow \|f(x)\| \leq \|f(y)\|) \\ &\leftrightarrow x \in X \wedge (y \in X \rightarrow \varphi(x) \leq \varphi(y)). \end{aligned} \quad \dashv$$

This also lifts to Σ_2^1 .

Corollary

$\text{PWO}(\Sigma_2^1)$ (and $\text{PWO}(\Sigma_2^1)$)

The idea is actually quite easy and nice, having the following form.

Theorem

*If Γ is adequate and X has a Γ -norm, then $\exists^{\mathcal{N}} X$ has a $\exists^{\mathcal{N}} \forall^{\mathcal{N}} \Gamma$ -norm.
So $\text{PWO}(\Gamma)$ implies $\text{PWO}(\exists^{\mathcal{N}} \forall^{\mathcal{N}} \Gamma)$.*

Proof.

- Let φ be a Γ -norm for X .
- Define $\psi : \exists^{\mathcal{N}} X \rightarrow \text{Ord}$ just by

$$\psi(x) = \min\{\varphi(x, y) : \langle x, y \rangle \in X\}.$$
- This is a $\exists^{\mathcal{N}} \forall^{\mathcal{N}} \Gamma$ -norm:

$$x \leq_{\psi} y \quad \text{iff} \quad \exists z \forall t (\langle x, z \rangle \leq_{\varphi} \langle y, t \rangle)$$

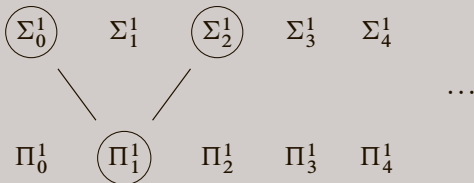
$$x <_{\psi} y \quad \text{iff} \quad \exists z \forall t (z = t \vee \langle x, z \rangle <_{\varphi} \langle y, t \rangle).$$

⊢

Question

We know $\text{PWO}(\Sigma_2^1)$. Does this imply $\text{PWO}(\Sigma_1^1)$ just because $\Sigma_1^1 \subseteq \Sigma_2^1$?

- The answer here is no: every Σ_1^1 -set has a Σ_2^1 -norm, but not necessarily a Σ_1^1 -norm.
- In fact, these three pointclasses (Σ_0^1 , Π_1^1 , and Σ_2^1) are the only three analytical pointclasses we can show PWO for under ZFC:



Prewellorders

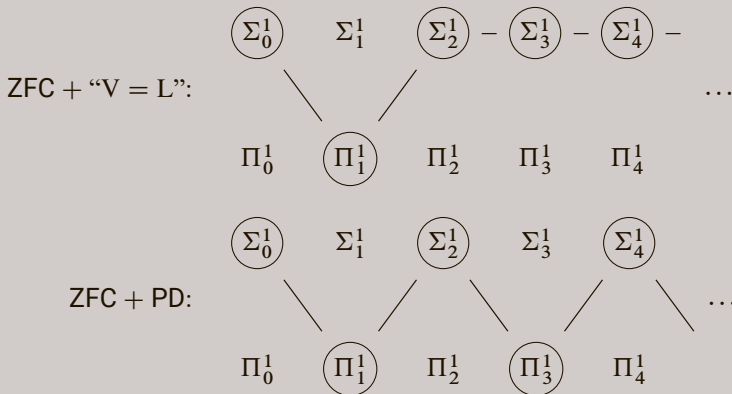
Γ^1 -norms

PWO

Scales

Uniformization from
Scales

The Main Goal



- In particular (assuming $\text{Con}(\text{ZFC} + \text{PD})$), $\text{PWO}(\Sigma_n^1)$ is independent for odd $n > 2$.
- For even $n > 2$, Harrington has some (unpublished) work which gives a model where $\text{PWO}(\Sigma_n^1)$ and $\text{PWO}(\Pi_n^1)$ both fail.

- One question that arises from this picture is can we have both $\text{PWO}(\Sigma_n^1)$ and $\text{PWO}(\Pi_n^1)$?
- The answer is “no”, and this is a result of properties more topological in nature.

Theorem (19 D • 2, 4, 6, and 7)

Let Γ be an adequate pointclass. Therefore

- *$\text{PWO}(\Gamma)$ implies Γ has the reduction property.*
- *Γ has the reduction property implies $\neg\Gamma$ has the separation property.*
- *If Γ has a universal set, Γ cannot have both the reduction and separation properties.*
- *Every σ -algebra (e.g. Δ_n^1) has both the reduction and separation properties.*

- What these properties *are* exactly isn't too important for us.
- They at least establish $\neg(\text{PWO}(\Gamma) \wedge \text{PWO}(\neg\Gamma))$ for $\Gamma = \Sigma_n^1$.
- They also tell us we shouldn't be considering Δ_n^1 , since they trivially have these properties.

Theorem

The following are equivalent:

- ❶ $\omega_1^{L[x]} < \omega_1$.
- ❷ $\text{PSP}(\Sigma_2^1(x))$.
- ❸ $\text{PSP}(\Pi_1^1(x))$.

- Let's get back on track.
- The main way we'll show (3) implies (1) above is to use Π_1^1 -uniformization (aka *Kondô's theorem*).
- In general, the way to show Γ -uniformization is with the scale property on Γ .
- As usual, we can prove the scale property on Π_1^1 and Σ_2^1 , but we can't go beyond this in ZFC alone.
- Determinacy axioms can push this further with the same alternating pattern as with PWO.

Definition

Let $X \subseteq \mathcal{N}$. A *scale* on X is a sequence $\vec{\varphi} = \langle \varphi_n : n < \omega \rangle$ such that

- each φ_n is a norm on X ;
- for all convergent $x \in {}^\omega X$, if each $\varphi_n \circ x$ is eventually constant then
 - $\lim x \in X$ and
 - (lower semi-continuity) $\varphi_n(\lim x) \leq \lim(\varphi_n \circ x)$ for each $n < \omega$.

Without lower semi-continuity, $\vec{\varphi}$ is called a *semi-scale*.

- As before, we are interested in scales with definability restrictions.
- We can easily get scales on any X we want by using AC.
- $f : X \rightarrow |X|$ a bijection with $\varphi_n = f$ for every n works.
- If $x \in {}^\omega X$ is convergent and $f \circ x$ is eventually constant then x is eventually constant. So $\lim x \in \text{im } x \subseteq X$.
- Lower semi-continuity is also easy in this case:
 $f(\lim x) = \lim(f \circ x)$.
- As before, we are interested in scales with definability restrictions.

- So what are some less trivial ways to get scales?
- How could scales possibly be useful?

Result (19 E • 3)

For $X \subseteq \mathcal{N}$ and $\kappa \geq \aleph_0$, the following are equivalent.

- *X has a scale $\vec{\varphi}$ where $\varphi_n(x) < \kappa$ for all $n < \omega$, $x \in X$.*
- *X has a semi-scale where $\vec{\varphi}$ where $\varphi_n(x) < \kappa$ for all $n < \omega$, $x \in X$.*
- *X is κ -suslin.*

- Suslin representations of sets are very important because they allow us to ask questions about trees instead of whatever weird reals we have.
- We will only prove the downward direction. The upward isn't very interesting to me.

Result (19 E • 3)

If X has a semi-scale bounded by κ then X is κ -suslin.

Proof.

- A scale is a semi-scale so one direction is obvious.
- If $\vec{\varphi}$ is a semi-scale, consider the tree building up sequences of elements and their norms: $\langle \tau, \rho \rangle \in {}^{<\omega}\omega \times {}^{<\omega}\kappa$ is in T iff there's an $x \in X$ with
 - $\tau \triangleleft x$;
 - $\rho = \langle \varphi_n(x) : n < \text{lh}(\tau) \rangle$.
- This is a tree over $\omega \times \kappa$. So $\mathfrak{p}[T]$ is κ -suslin.
- $\langle x, \langle \varphi_n(x) : n < \omega \rangle \rangle \in [T]$ for any $x \in X$ so $X \subseteq \mathfrak{p}[T]$.

Result (19 E • 3)

If X has a semi-scale bounded by κ then X is κ -suslin.

Proof.

- $\langle \tau, \rho \rangle \in T$ iff there's an $x \in X$ with
 - $\tau \triangleleft x$;
 - $\rho = \langle \varphi_n(x) : n < \text{lh}(\tau) \rangle$.
- To show $\mathfrak{p}[T] \subseteq X$, if $x \in \mathfrak{p}[T]$, then we get a sequence $\langle x_n \in X : n < \omega \rangle$ witnessing $x \upharpoonright n \in \mathfrak{p}T$.
- $x \upharpoonright n = x_n \upharpoonright n$ so $\langle x_n \in X : n < \omega \rangle$ converges to x with $\langle \varphi_n(x_m) : m < \omega \rangle$ eventually determined and thus constant.
- Hence $\lim x_n = x \in X$. ⊢

- The motivating result for PWO was that WO had a Π_1^1 -norm given by $x \mapsto \|x\|$.
- We get a similar result for scales.

Result (19 E • 4)

There is a scale on WO. In fact, the relations on triples $\langle x, y, n \rangle \in \mathcal{N}^2 \times \omega$ are both Π_1^1 :

$$x \leq_{\varphi_n} y \quad \text{iff} \quad x \in \text{WO} \wedge (y \in \text{WO} \rightarrow \varphi_n(x) \leq \varphi_n(y))$$

$$x <_{\varphi_n} y \quad \text{iff} \quad x \in \text{WO} \wedge (y \in \text{WO} \rightarrow \varphi_n(x) < \varphi_n(y)).$$

- What are φ_n ? They code $\|x\|$ and its initial segment “up to n ”.
- More precisely, we can consider

$$(E_x)_{<n} = \{\langle a, b \rangle : a E_x b \wedge a E_x n \neq b\}.$$

- These $(E_x)_{<n}$ tell us how E_x is built up.
- So we define $\varphi_n(x) = \text{code}(\|E_x\|, \|(E_x)_{<n}\|)$.
- Showing $\langle \varphi_n : n < \omega \rangle$ is in fact a scale is a boring, technical process.

- This motivates the concept of a Γ -scale.

Definition

Let Γ be a pointclass. A Γ -scale is a scale $\vec{\varphi}$ on a set X such that the relations on $\langle x, y, n \rangle \in \mathcal{N}^2 \times \omega$ are in Γ :

$$x \leq_{\varphi_n} y \quad \text{iff} \quad x \in X \wedge (y \in X \rightarrow \varphi_n(x) \leq \varphi_n(y))$$

$$x <_{\varphi_n} y \quad \text{iff} \quad x \in X \wedge (y \in X \rightarrow \varphi_n(x) < \varphi_n(y)).$$

- So for adequate pointclasses, X having a Γ -scale implies $X \in \Gamma$ by $x \in X$ iff $x \leq_{\varphi_0} x$.
- We can also show that every Π_1^1 -set has a Π_1^1 -scale in the same way as with PWO.

Theorem

Every Π_1^1 -set has a Π_1^1 -scale, i.e. the scale property holds for Π_1^1 .

Proof.

- Let $A \in \Pi_1^1$, $A = f^{-1}\text{"WO"}$ for some computable $f : \mathcal{N} \rightarrow \mathcal{N}$.
 - Let $\vec{\psi}$ a Π_1^1 -scale on WO.
 - We can form $\langle \psi_n \circ f : n < \omega \rangle$ and get a Π_1^1 -scale on A . $\dashv$
-
- We can also show the scale property for Σ_2^1 .
 - Indeed, we get the same restrictions as with PWO: after Σ_2^1 ,
 - the scale property is independent of ZFC;
 - PD gives a full zig-zag pattern; and
 - $V = L$ gives an initial zig followed by a line.
 - Again, we mostly care about scales for now to get uniformization.

Definition

Let $X \subseteq A \times B$. A *uniformization* is a function $f \subseteq X$ such that $\text{dom}(f) = \text{dom}(X)$.

For Γ a pointclass, Γ -uniformization is the statement that every $X \in \Gamma$ has a uniformization $f \in \Gamma$.

Let's prove Π_1^1 -uniformization now that we have the scale property.

Theorem

Π_1^1 -uniformization holds (and $\mathfrak{N}_1^1$ -uniformization).

- Let $A \in \Pi_1^1$ be arbitrary and let $A''x = \{y \in \mathcal{N} : \langle x, y \rangle \in A\}$.
- To get a uniformization, we need to find a unique $y \in A''x$ defined from x .
- Uniqueness is the easy part. The hard part is showing existence and ensuring the resulting f is Π_1^1 with $\text{dom}(f) = \text{dom}(A)$.
- That is where scales come in.

Theorem

Π_1^1 -uniformization holds (and $\tilde{\Pi}_1^1$ -uniformization).

- Let $x \in A$ be arbitrary and $\vec{\varphi}$ a Π_1^1 -scale on $A \in \Pi_1^1$.
- Set $f(x) = y$ iff $\langle x, y \rangle \in A$ and for all $z \in \mathcal{N}$ and all $n < \omega$,
 If $(*) \langle x, z \rangle \in A, z \upharpoonright n = y \upharpoonright n$, and $\varphi_m(x, z) = \varphi_m(x, y)$ for all $m < n$,
 Then $(**) y(n) < z(n)$, or else $y(n) = z(n) \wedge \varphi_n(x, y) \leq \varphi_n(x, z)$.
- Basically, y is lexicographically least among those whose first norms agree with y 's initial first segments.
- Let's show that this indeed uniquely defines a y (if there exists one).
- Suppose this holds for two $y_1, y_2 \in \mathcal{N}$ and let N be their first disagreement.
- $(*)$ therefore holds inductively for $n \leq N$ (consider $(*)$ with $n = 0$ and then $(**)$ with $n = 0$ to get $\varphi_n(x, y_1) = \varphi_n(x, y_2)$)
- $(**)$ thus holds. But applying $(*)$ with $z = y_1$ and $y = y_2$ gives $y_2(N) < y_1(N)$ and vice versa yields $y_1(N) < y_2(N)$, a contradiction.

Theorem

Π_1^1 -uniformization holds (and $\tilde{\Pi}_1^1$ -uniformization).

- Set $f(x) = y$ iff $\langle x, y \rangle \in A$ and for all $z \in \mathcal{N}$ and all $n < \omega$,
 If $(*) \langle x, z \rangle \in A, z \upharpoonright n = y \upharpoonright n$, and $\varphi_m(x, z) = \varphi_m(x, y)$ for all $m < n$,
 Then $(**) y(n) < z(n)$, or else $y(n) = z(n) \wedge \varphi_n(x, y) \leq \varphi_n(x, z)$.
- It's not hard to see that $(*)$ is Σ_1^1 and $(**)$ is Π_1^1 .
- Thus f (if well-defined) will be $\forall^{\mathcal{N}} (\neg \Sigma_1^1 \vee \Pi_1^1) = \Pi_1^1$.
- So we need to show there *is* such a y , which is much harder.
- We use the scale to do the heavy lifting in the construction and ensure the limit (i.e. end result) is in A .
- Start out with $A_0 = A \restriction x$.
- At stage $n + 1$, we first consider the set of all $y \in A_n$ with minimal $\varphi_n(x, y)$ (among other elements in A_n).
- Then we thin out that set to only consider the y s with minimal $y(n)$. The result is A_{n+1} .
- We now want a $y \in \bigcap_{n < \omega} A_n$.

Theorem

Π_1^1 -uniformization holds (and $\tilde{\Pi}_1^1$ -uniformization).

- Start out with $A_0 = A''x$.
- At stage $n + 1$, we first consider the set of all $y \in A_n$ with minimal $\varphi_n(x, y)$ (among other elements in A_n).
- Then we thin out that set to only consider the y s with minimal $y(n)$. The result is A_{n+1} .
- Any sequence $\langle y_n \in A_n : n < \omega \rangle$ is necessarily convergent since we're deciding more and more when moving from A_n to A_{n+1} .
- Such a sequence also has $\langle \varphi_n(y_k) : k < \omega \rangle$ as eventually constant because again we're deciding more of the norm moving from A_n to A_{n+1} .
- As a scale, it follows that $y = \lim_{n \rightarrow \infty} y_n \in A$ and an easy induction shows it satisfies the definition of $f(x) = y$.
- Hence f is well-defined and a Π_1^1 -uniformization of A . $\dashv$

- Again, we have the same picture with uniformization as with PWO, scales, and other properties: Σ_2^1 is the limit of what can be shown in ZFC.
- The zig-zag (or lackthereof in L) of PD also holds with uniformization.
- This is quite useful with determinacy because uniformization is a choice-like principle telling us not only that we can pick out an element, but we can do so in a simply definable way.
- Let's return to the main goal here:

Theorem

The following are equivalent:

- 1 For every $x \in \mathcal{N}$, $\mathbb{L}[x] \models \text{"}\omega_1^V \text{ is inaccessible"}$.
- 2 $\text{PSP}(\prod_1^1)$.

- Restated, the main goal is the following.

Theorem

$\text{PSP}(\Pi_1^1(x))$ implies $\omega_1^{\mathbb{L}[x]} < \omega_1$ for every $x \in \mathcal{N}$.

Proof.

- Let $x \in \mathcal{N}$ and suppose $\omega_1^{\mathbb{L}[x]} = \omega_1$, aiming to show $\neg \text{PSP}(\Pi_1^1(x))$.
- For each $\alpha < \omega_1^{\mathbb{L}[x]}$, let $f(\alpha)$ be the $<_{\mathbb{L}[x]}$ -least real in $\text{WO} \cap \mathbb{L}[x]$ coding $\langle \alpha, \in \rangle$.
- Set $X = f''\omega_1^{\mathbb{L}[x]}$ so that every ordinal is coded by an element of X .
- We actually considered this set already to show that $\mathbb{L} \models \neg \text{PSP}(\Sigma_2^1)$.
- The relevant theorems generalize to show
 - $X \in \Sigma_2^1(x)$ is uncountable (since $\aleph_1^{\mathbb{L}[x]} = \aleph_1$), and
 - X has no uncountable Σ_1^1 -subset (by the Boundedness lemma).

Theorem

$\text{PSP}(\Pi_1^1(x))$ implies $\omega_1^{L[x]} < \omega_1$ for every $x \in \mathcal{N}$.

Proof.

- $X = f''\omega_1^{L[x]}$ where $f(\alpha) \in \text{WO} \cap L[x]$ codes α .
 - ① $X \in \Sigma_2^1(x)$ is uncountable (since $\aleph_1^{L[x]} = \aleph_1$), and
 - ② X has no uncountable Σ_1^1 -subset (by the Boundedness lemma).
- Let $X = pY$ for $Y \in \Pi_1^1(x)$.
- By $\Pi_1^1(x)$ -uniformization, we get a Π_1^1 -function $f \subseteq Y$ with $\text{dom}(f) = X$.
- Hence $|f| = |X| = \aleph_1$.
- f also can't have any perfect subset. To see this, firstly any perfect subset must be closed and uncountable.
- But any closed $g \subseteq f$ has $\text{dom}(g) \in \Sigma_1^1$ and is thus countable since $\text{dom}(g) \subseteq X$ and (2) holds.
- Thus $\neg \text{PSP}(\Pi_1^1(x))$. ⊥

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