

Lecture 11.

Review for MT1.

Propositional Algebra and Proofs (Ch.1).

1. Propositions and connectives.

Connectives are functions of some propositional variables $f(P_1, \dots, P_n)$ taken proposition depending of $P_1, \dots, P_n$ (their True values).

They define by their True tables.

2. Basis connectives (operations)

conjunction	$\wedge$
disjunction	$\vee$
negation	$\sim$
conditional	$\Rightarrow$
biconditional	$\Leftrightarrow$

Supplementary operations

Sheffer $\square$

Alternative disjunction

3. Propositional formula -

composition (superposition)
of formulas. operations.

Don't forget ~~an~~ brackets which we
need for reading of formulas.

4. 2 functions ^{are} equivalent

$$f(P_1, \dots, P_n) \equiv g(P_1, \dots, P_n)$$

⚡ (same variables!)

iff they have identical True tables.

$f \equiv g$ iff $f \Leftrightarrow g$ is Tautology.

2 ways to prove equivalency:

- by True tables

- using Laws of Propositional

operations (Tautologies)

Remarks. Instead of complete tables
we can use **Truncated** ones:

~~which can~~ compare only values of
 $P_1, \dots, P_n$ on which f, g are True
(or False).

5. We have several groups of Laws:

- basic laws for $(\vee, \wedge, \neg)$: commutative, associative, distributive, de Morgan.
- absorption laws (like $P \vee (P \wedge Q) = P$, $P \vee (\neg P \wedge Q) = P \vee Q$).
- 2 groups for conditionals: direct and for contraposition.

Always give clear references when making transformations.

6. 2 important types problems:

- find if $f(P_1, \dots, P_n)$ is Tautology, Contradiction or Neither. At first 2 cases you need verify that f takes the same values for all values of variables; or 3rd one - find 2 combinations on which it takes T and F.
- transformation formulas to convenient form: when negations stay only for variables.

7. Complete systems of operations.

We can any connective $f(P_1, \dots, P_n)$ represented given by True table to represent by an equivalent formula with $\wedge, \vee, \neg$.

So these 3 operations give a complete system of operations. Then we can transform any formula with only $\vee, \neg$, either $\wedge, \neg$.

8. Correspondence between Propositional and Boolean functions.

9. Conditional, Biconditional, Theorems.

Properties of $A \Rightarrow B, A \Leftrightarrow B$.

Theorem $P \Rightarrow Q$; Inverse Theorem $Q \Rightarrow P$
($P \Rightarrow Q \Leftrightarrow Q \Rightarrow P$).

Proofs direct, by contradiction ($\neg Q \Rightarrow \neg P \Leftrightarrow P \Rightarrow Q$)

To prove $P \Leftrightarrow Q$ we need prove both:

direct $P \Rightarrow Q$ and inverse Theorems $Q \Rightarrow P$.

10. Proofs of 3 Theorems: irrationality $\sqrt{p}$; infinity of prime numbers; numbers from only 5.

$$9b) b = aq + r, 0 \leq r < |a|$$

$$a = 5, b = 36$$

$$36 = 5q + r \Rightarrow q = 7, r = 1.$$

~~Sect~~

$$b) a = -8, b = -52$$

$$-52 = -8 \cdot q + r, 0 \leq r \leq 8$$

$$q = 7, r = 4$$

Sect. 1.3.

1. (c) $P(x)$ - x is isosceles

$Q(x)$ - x is ~~not~~ right

$$\exists x (P(x) \wedge Q(x))$$

$$(d) \forall x (Q(x) \Rightarrow \neg P(x)).$$

(f) $P(x)$ x is honest

$$(\exists x P(x)) \wedge (\exists y \neg P(y))$$

$$(g) \forall x \quad x \neq 0 \Rightarrow (x > 0 \vee x < 0)$$

(iv) $P(x, y)$ x loves y

$$\forall x \exists y P(x, y)$$

$$(10) \quad \forall x \quad x > 0 \Rightarrow (\exists y \quad x = 2^y) \wedge \forall z \quad (x = 2^z \Rightarrow y = z)$$

$$(\forall x) (x > 0 \Rightarrow (\exists y!) \quad x = 2^y)$$

5. $P(x, y)$ Person x dislikes taxes y

$$\forall x \forall y P(x, y), \quad \forall x \exists y P(x, y), \quad \exists x \forall y P(x, y)$$

$$\exists x \exists y P(x, y).$$

6. b) T, c) T, d) T

8. a) ~~F~~ g) F

9. a) For all ~~$x \in \mathbb{N}$~~ natural x we have
 $x \geq 1$.

c) ~~Each~~ Each prime number $x \neq 2$ ~~is~~
 is odd.