SPANNING TREES AND CONTINUED FRACTIONS

SWEE HONG CHAN, ALEX KONTOROVICH, AND IGOR PAK

ABSTRACT. We prove the exponential growth of the cardinality of the set of numbers of spanning trees in simple planar graphs on n vertices, answering a question from 1969. The proof uses a connection with continued fractions and recent advances towards Zaremba's conjecture.

1. INTRODUCTION

1.1. Foreword. Often lacking structure and extensive background, many combinatorics results are so simple that they can be proved "in a few lines by anybody obtuse enough to feel the need for verification", cf. [Dys44]. In the opposite extreme, rather unsurprisingly, hard problems require an inordinary amount of work and technology, either homegrown or borrowed from other areas.

In this paper, we analyze the number of spanning trees of planar graphs, by employing Bourgain– Kontorovich's technology towards Zaremba's conjecture on continued fractions. Put simply, we use recent results in number theory and dynamical systems to settle an old problem in graph theory. These results are technical and require some delicacy in applying. We choose a combination of both research and survey styles, in an effort to be both accessible and concise.

1.2. Main results. Let G = (V, E) be a simple graph, i.e. a graph without loops and multiple edges. Denote by $\tau(G)$ the number of spanning trees in G. Almost 60 years ago, Sedláček initiated in [Sed66] the study of the values taken by $\tau(G)$ on planar graphs, graded by the number of vertices of G. To this end, for $n \geq 3$, let $\mathcal{T}(n)$ denote the set of numbers of spanning trees of connected planar simple graphs on n vertices:

(1.1)
$$\mathcal{T}(n) := \{ \tau(G) : G = (V, E) \text{ is connected, simple and planar, } |V| = n \}.$$

(See also §5.3 for a related set of interest.) For example, $\mathcal{T}(3) = \{1,3\}$, $\mathcal{T}(4) = \{1,3,4,8,16\}$ and $\mathcal{T}(5) = \{1,3,4,5,8,9,11,12,16,20,21,24,40,45,75\}$. It is easy to see that the sets $\mathcal{T}(n)$ are nested, i.e. $\mathcal{T}(n) \subseteq \mathcal{T}(n+1)$, that $n \in \mathcal{T}(n)$, and that $2 \notin \mathcal{T}(n)$ for all n.

Let $|\mathcal{T}(n)|$ be the cardinality of $\mathcal{T}(n)$, that is, the number of distinct values of $\tau(G)$ in $\mathcal{T}(n)$; so $|\mathcal{T}(4)| = 5$ and $|\mathcal{T}(5)| = 15$. A simple argument using Euler characteristic (see §5.1) shows that (1.2) $\max \mathcal{T}(n) < C^n$,

for some C > 1, and hence $|\mathcal{T}(n)|$ grows at most exponentially:

$$|\mathcal{T}(n)| < C^n,$$

for all *n* sufficiently large. The study of lower bounds on $|\mathcal{T}(n)|$ was initiated by Sedláček in [Sed69], who proved that $|\mathcal{T}(n)| = \Omega(n^2)$. In [Aza14], Azarija showed that $|\mathcal{T}(n)| = e^{\Omega(\sqrt{n/\log n})}$. It follows from Stong's theorem (Theorem 1.3) that $|\mathcal{T}(n)| = e^{\Omega(n^{2/3})}$. Our main result shows that the sequence $|\mathcal{T}(n)|$ indeed has exponential growth:

Theorem 1.1 (Main Theorem). There exists a constant c > 1 such that

(1.3) $|\mathcal{T}(n)| > c^n$

holds for all sufficiently large n.

Date: November 11, 2024.

See §5.1 for an explicit estimate for the value $c \approx 1.0087$ derived from our proof method. To understand why the lower bound remained out of reach until now, note that there are only exponentially many nonisomorphic connected simple planar graphs to begin with,¹ and exponentially many of them may have the same number of spanning trees.²

Theorem 1.1 is a consequence of the following stronger result, that the set $\mathcal{T}(n)$ contains a positive proportion of an exponentially long set of integers.

Theorem 1.2. There exists a constant c > 1, such that

(1.4)
$$\liminf_{n \to \infty} \frac{1}{c^n} \left| \mathcal{T}(n) \cap \{1, \dots, c^n\} \right| > 0.$$

In fact, though we do not need it for the Main Theorem 1.1, it is possible to improve Theorem 1.2 from positive proportion to density one, that is, that the expression in (1.4) tends to 1; see Remark 1.16.

1.3. Ingredients. As we describe in more detail below, there are four main steps in the proof of Theorem 1.2.

- (i) The first is to dualize (cf. §1.4), replacing the study of $\mathcal{T}(n)$ with the closely related quantity $\alpha(t)$ defined in (1.5).
- (ii) The second step is to employ a mechanism that constructs simple planar graphs G = (V, E) with a prescribed number of spanning trees, $\tau(G)$, in such a way as to simultaneously allow some control on the number |V| of vertices. It turns out (cf. §1.5) that this control relies on certain Diophantine-theoretic properties of continued fraction expansions of rational numbers. This reduces the graph theory problem to one in Diophantine geometry.
- (iii) The Diophantine problem turns out to be amenable to techniques similar to those attacking Zaremba's conjecture (cf. §1.6); one is then able to produce a large enough collection of the desired fractions, as long as a certain Cantor-like fractal's Hausdorff dimension is sufficiently large.
- (iv) And the final ingredient is to verify numerically and rigorously (cf. §1.7) that this fractal dimension is indeed large enough.

It is in this last step that we turn out to be extraordinarily lucky: the currently best available sufficiency condition is indeed satisfied by the numerics, but only just barely, in the hundredths place! (See Remark 1.18.)

1.4. **Dualizing.** To better understand Theorem 1.2, consider the sets dual to $\mathcal{T}(n)$, namely, for $t \geq 3$, we define $\widehat{\mathcal{T}}(t) := \{n : t \in \mathcal{T}(n)\}$. In light of the nesting property $\mathcal{T}(n) \subseteq \mathcal{T}(n+1)$, the only quantity of interest for $\widehat{\mathcal{T}}(t)$ is its minimal element, which we denote by

(1.5)
$$\alpha(t) := \min \widehat{\mathcal{T}}(t) = \min\{n : t \in \mathcal{T}(n)\}.$$

That is, $\alpha(t)$ is the smallest number of vertices of a planar simple graph with exactly t spanning trees. Then $\widehat{\mathcal{T}}(t)$ consists of $\alpha(t)$, followed by every subsequent integer. By (1.2), we have that

(1.6)
$$\alpha(t) = \Omega(\log t).$$

The study of $\alpha(t)$ was also initiated by Sedláček [Sed70] and continued over the years, see [Aza14, AŠ13, CP24a] and §5.3. Until Stong's recent breakthrough, even $\alpha(t) = o(t)$ remained open, see [AŠ13, Question 1].

Theorem 1.3 (Stong [Sto22, Cor. 7.3.1]). $\alpha(t) = O((\log t)^{3/2}/(\log \log t))$.

¹While the exact asymptotics remain open, the number of unlabeled simple planar graphs on n vertices is $O(30.061^n)$, see [B+06], [Noy15, §6.9.2] and [OEIS, A003094].

²For example, there are $\Omega(2.955^n)$ unlabeled trees on *n* vertices, see e.g. [OEIS, A000055] and [OEIS, A051491].



FIGURE 1.1. A typical graph constructed in the Main Graph Theorem 1.7

It is natural to conjecture (see also Proposition 1.9) that the true upper bound for $\alpha(t)$ matches the lower bound in (1.6):

Conjecture 1.4. $\alpha(t) = O(\log t)$ for all $t \ge 3$.

Note that $\alpha(pq) \leq \alpha(p) + \alpha(q)$ (by taking disjoint union of graphs), so it suffices to prove the conjecture for prime t. Theorem 1.2 follows from the resolution of the following weaker version of this conjecture.

Theorem 1.5 (= Theorem 1.2). Conjecture 1.4 holds for a positive proportion of t. That is, there is a set of positive proportion within the natural numbers for which the estimate $\alpha(t) = O(\log t)$ holds as $t \to \infty$ within this set.

Conjecture 1.4 remains out of reach. In fact, even the following weaker problem is open. Denote by $\beta(t)$ the smallest number of edges of a planar graph with exactly t spanning trees. This function was introduced by Nebeský in [Neb73]. Note that multiple edges are allowed in this case, and that we have $\beta(t) < 3\alpha(t)$. It was proved in [CP24c], that $\beta(t) = O(\log t \log \log t)$, see §5.5. This is very close but still shy of the natural upper bound that would follow from Conjecture 1.4:

Conjecture 1.6. $\beta(t) = O(\log t)$.

1.5. Continued Fractions and Graphs. Given $a_0 \ge 0$, $a_1, \ldots, a_\ell \ge 1$, where $\ell \ge 0$, the corresponding *continued fraction* is defined as follows:

$$[a_0; a_1, \dots, a_\ell] := a_0 + \frac{1}{a_1 + \frac{1}{\ddots + \frac{1}{a_\ell}}}.$$

Integers a_i are called *partial quotients*, see e.g. [HW08, §10.1]. We use the notation $[a_1, \ldots, a_\ell]$ when $a_0 = 0$. The following elementary result, proved in §2, gives a connection between spanning trees and continued fractions:

Theorem 1.7 (Main Graph Theorem). Let $t, u \ge 1$ be positive integers with t < u and gcd(t, u) = 1. Suppose that

(1.7)
$$\frac{t}{u} = [b_1, 1, b_2, 1, \dots, b_m, 1],$$

for some $b_1, \ldots, b_m \geq 1$. Then there exists a simple planar graph G such that

$$\tau(G) = t$$
, and $|V| = b_1 + \ldots + b_m + 2$.

The construction, explained in §2, is completely explicit; see Figure 1.1 for the typical structure of such a graph. The proof is inductive, and is a variation on the proof of Theorem 1.5 in [CP24c] and a construction of Bier in [Bier10]. Motivated by this theorem, we formulate the following.

Conjecture 1.8 (Diophantine Conjecture). There is a universal constant A > 0 so that, for every integer $t \ge 3$, there is a coprime integer u > t, such that the quotient t/u has continued fraction expansion (1.7) with $b_1, \ldots, b_m \le A$.

We discuss in §5.6 why this conjecture may be plausible. Regardless, together with the Main Graph Theorem 1.7, it would settle Conjecture 1.4.

Proposition 1.9. The Diophantine Conjecture 1.8 implies Conjecture 1.4.

Proof. For any integer $t \ge 3$, Conjecture 1.8 produces a fraction t/u having continued fraction expansion as in (1.7). By Theorem 1.7, there is a simple planar graph G such that $t = \tau(G)$. This graph also has $|V| \le \frac{(A+1)}{2}m + 2 = O(\log t)$ vertices; hence $\alpha(t) = O(\log t)$, as required.

While we are not able to establish Conjecture 1.8, we prove the following approximation.

Theorem 1.10. Conjecture 1.8 holds for a set of positive proportion. That is, there is some A > 0, so that the set of t for which the conclusion of Conjecture 1.8 holds, has positive proportion in the natural numbers.

We can thus reduce all our efforts to proving Theorems 1.7 and 1.10.

Proof of Theorem 1.5. The proof follows verbatim the proof of Proposition 1.9, but applied to a positive proportion subset of t.

1.6. Zaremba's conjecture. The Diophantine Conjecture 1.8 is itself closely related to the following celebrated open problem (see, e.g., [Kon13] and §3 for more background):

Conjecture 1.11 (Zaremba's Conjecture [Zar72, p. 76]). There is a universal constant A > 0 so that, for every integer $u \ge 1$, there is a coprime integer $1 \le t < u$, such that $t/u = [a_1, \ldots, a_\ell]$ and $a_1, \ldots, a_\ell \le A$.

Here the change to denominators from numerators in the Diophantine Conjecture 1.8 is completely anodyne, see Remark 3.1. Zaremba himself conjectured that one can take A = 5 in Conjecture 1.11, see also §5.7. In [BK14], Bourgain–Kontorovich proved the following result, on which the proof of Theorem 1.10 is based, and which shows that Zaremba's conjecture holds for a positive proportion of denominators (in fact, they prove a density one version, but again, we will only quote the weaker statement):

Theorem 1.12 ([BK14]). There is a universal constant A > 0 so that the conclusion of Zarembra's Conjecture 1.11 holds for a positive proportion of u.

A key role is played here by the Cantor-like sets comprising the limit points of these Diophantine fractions. In the setting of Zaremba's conjecture, the limit set, for a given A > 0, is the following:

$$\mathfrak{F}_A := \left\{ [a_1, a_2, \ldots] : 1 \le a_i \le A, \forall i \right\}.$$

Let $\delta_A = \text{Hdim}(\mathfrak{F}_A)$ denote the Hausdorff dimension of \mathfrak{F}_A ; it is classical that $\delta_A \nearrow 1$ as $A \to \infty$. Then the more precise version of Theorem 1.12 is the following.

Theorem 1.13 ([BK14]). There is a constant $\delta_0 < 1$, such that, for any A with $\delta_A > \delta_0$, Zaremba's Conjecture 1.11 holds for a positive proportion of u.

We explain in \$3 how the same methods can be adapted to prove an analogous theorem in the context of §1.5. For A > 0, define the limit set:

$$\mathfrak{C}_A := \{ [b_1, 1, b_2, 1, \ldots] : 1 \le b_i \le A, \forall i \},\$$

and let

$$\vartheta_A := \operatorname{Hdim}(\mathfrak{C}_A)$$

be its Hausdorff dimension.

Theorem 1.14 (Main Diophantine Theorem). There is a constant $\delta_0 < 1$, such that, for any A > 0 with $\vartheta_A > \delta_0$, the Diophantine Conjecture 1.8 holds for a positive proportion of t.

Let

(1.8)
$$\mathfrak{C}_{\infty} = \bigcup_{A=1}^{\infty} \mathfrak{C}_A$$

and $\vartheta_{\infty} := \operatorname{Hdim} \mathfrak{C}_{\infty}$, so that $\vartheta_A \nearrow \vartheta_{\infty}$ as $A \to \infty$.

In order to conclude Theorem 1.10, we thus need to verify that the dimension ϑ_{∞} indeed exceeds the numerical value of δ_0 coming from the proof of Theorem 1.14. The original work [BK14] gave the value of $\delta_0 = 0.984$, and there have been a number of subsequent improvements bringing down this critical parameter. As explained in $\S3$, the proof of (a density-one version of) Theorem 1.13 follows the "orbital circle method" which relies on an analysis of a "minor arcs" estimate and "major arcs" estimate.

In [FK14], Frolenkov–Kan gave an improved treatment of the minor arcs only, allowing them to reduce δ_0 to $5/6 = 0.8\overline{3}$, but at the cost of concluding that Zaremba's conjecture holds for only a positive proportion of t, rather than a density one set. In [Hua15], Huang was able to combine the major arcs analysis from [BK14] with the minor arcs treatment in [FK14] to recover $\delta_0 = 5/6$ for a density one set. Since then, there have been a number of improvements on the minor arcs analysis by Kan, who proved the positive proportion version with:

•
$$\delta_0 = 4/5 = 0.8$$
 in [Kan15]

- $\delta_0 = \frac{4}{3} = 0.8$ m [Kan15], $\delta_0 = \frac{1}{4}(\sqrt{17} 1) \approx 0.781$ in [Kan17], and most recently $\delta_0 = \frac{1}{3}(\sqrt{40} 4) \approx 0.775$, see [Kan21].

Theorem 1.15 (cf. [Kan21]). *Theorem 1.14 holds with* $\delta_0 = 0.775$.

It thus remains to verify that ϑ_{∞} exceeds this value.

Remark 1.16. Combined with Huang's method, it is possible to prove that density one sets of t satisfy Zaremba's conjecture, with these same parameters of δ_0 in Theorem 1.14. While doing this properly with all details would multiply the length of this paper five-fold, we discuss the key steps in §3.3; executing this would make Theorems 1.14, 1.10, 1.5, and 1.2 hold with "positive proportion" replaced by "density one," but would not affect the Main Theorem 1.1.

1.7. Computing Hausdorff dimensions. The numerical estimation of Hausdorff dimensions of fractals corresponding to various families of continued fractions is a major object of study in dynamical systems. Specifically, estimating δ_A has become a benchmark problem in the area.

Famously, Good [Good41] showed that $0.5194 < \delta_2 < 0.5433$, greatly improving upon Jarník's 1928 lower bound $\delta_2 \geq 0.25$. Bumby [Bum85] used Good's approach and computer assistance to calculate $0.5312 < \delta_2 < 0.5314$. Although the exact value is not known, further algorithmic and computational advances led to better estimates, see [OEIS, A279903]. Notably, Hensley [Hen96] improved the bounds to $\delta_2 \approx 0.53128051$, and eventually this value was computed to 25 digits [JP01], 54 digits [Jen04], 100 digits [JP18], and most recently to 200 digits [PV22].

Implementing the Pollicott–Vytnova algorithm [PV22] in our setting, we are able to show the following.

Theorem 1.17 (Main Dimension Theorem). Let ϑ_{∞} denote the Hausdorff dimension of the fractal \mathfrak{C}_{∞} in (1.8). Then $\vartheta_{\infty} > \delta_0$, for $\delta_0 = 0.775$. In particular, for A = 110, we already have that $\vartheta_A > \delta_0 = 0.775$.

This allows us to apply the Main Diophantine Theorem 1.14 and Theorem 1.15 to conclude Theorem 1.10.

Remark 1.18. We have made no effort to optimize the value of A = 110. That said, note that it is possible to prove also the upper bound $\vartheta_{\infty} < 0.799$ (see §4.2), showing that ϑ_{∞} exceeds in the hundredths place the best available value of the parameter $\delta_0 = 0.775$.

1.8. **Paper structure.** As explained above, the Main Theorem 1.1 follows from a combination of the Main Graph Theorem 1.7, the Main Diophantine Theorem 1.14 and the Main Dimension Theorem 1.17. We present proofs of each in separate sections (Sections 2–4). In our presentation, we aim for the middle ground, giving a quick overview of the proof to make it accessible, but filling enough details to make it complete. We conclude with final remarks and open problems in Section 5.

2. Planar graphs and continued fractions

2.1. Recursion identities. Let G = (V, E) be a graph without loops, but possibly with multiple edges. Graph G is called *simple* if it does not have multiple edges.

For an edge $e \in E$, denote by G-e the graph obtained from G by deleting the edge e. Similarly, denote by G/e the graph obtained by contracting vertices incident to e and removing the resulting loops. As in the introduction, denote by $\tau(G)$ the number of spanning trees in G. Recall that $\tau(G) = \tau(G-e) + \tau(G/e)$, see e.g. [Tut84, Thm II.18].

Let G = (V, E) be a graph and let $e \in E$ be an edge in G. A pair (G, e) is called a *marked* graph. Define the spanning tree vector of the marked graph as follows:

$$\boldsymbol{v}(G,e) := (\tau(G-e), \tau(G/e)).$$

We consider two dual operations on marked graphs. Denote by $\Phi^k(G, e)$ the marked graph (G', e') obtained by replacing edge e with a path of length (k + 1), and marking one of the new edges as e' (see Figure 2.1). Note that if G is connected, simple and planar, then so is G'. Denote G' = (E', V'), and observe that |E'| = |E| + k, |V'| = |V| + k.

Lemma 2.1. Suppose that $\tau(G-e) > 0$ and $\tau(G/e) > 0$. Then we have that

$$\boldsymbol{v}(G',e') = \boldsymbol{v}(G,e) \cdot \left(\begin{array}{cc} 1 & k \\ 0 & 1 \end{array}\right).$$

Proof. Note that $\Phi^k = \Phi^1(\Phi^{k-1})$. Thus it suffices to prove the lemma for k = 1, as the general case follows by induction. Observe that $\tau(G' - e') = \tau(G - e)$, since the new edge created by Φ is a bridge to an endpoint in G' - e'. We also have $G'/e' \simeq G$ by the construction of Φ , which gives $\tau(G'/e') = \tau(G) = \tau(G - e) + \tau(G/e)$. This gives the result. \Box

For the second operation, denote by $\Psi^k(G, e)$ the marked graph (G'', e'') obtained by replacing edge e with (k + 1) parallel edges, and marking one of the new edges as e'' (see Figure 2.1). Note that if G is connected and planar, then so is G'', but simplicity is not in general preserved, since we are creating multiple edges. Denote G'' = (E'', V''), and observe that |E''| = |E| + k, |V''| = |V|.

Lemma 2.2. Suppose that $\tau(G-e) > 0$ and $\tau(G/e) > 0$. Then we have:

$$\boldsymbol{v}(G'',e'') = \boldsymbol{v}(G,e) \cdot \left(\begin{array}{cc} 1 & 0 \\ k & 1 \end{array} \right).$$

The proof is similar to the proof of Lemma 2.1 and will be omitted. In fact, Lemma 2.2 follows from Lemma 2.1 via matroid duality, but we will not need this observation. Finally, note that (2.1) $\operatorname{gcd}\left(\tau(G'-e'),\tau(G'/e')\right) = \operatorname{gcd}\left(\tau(G''-e''),\tau(G''/e'')\right) = \operatorname{gcd}\left(\tau(G-e),\tau(G/e)\right),$ for all $k \geq 1$.



FIGURE 2.1. Marked graphs (G, e), $(G', e') = \Phi^2(G, e)$, $(G'', e'') = \Psi^3(G, e)$ and $(H, f) = \Upsilon^2(G, e)$.

2.2. **Proof of Theorem 1.7.** We will prove the theorem as a consequence of the following slightly more general lemma.

Lemma 2.3. Let $t, u \ge 1$ be positive integers with $t \le u$ and gcd(t, u) = 1. Suppose that

(2.2)
$$\frac{t}{u} = [b_1, 1, b_2, 1, \dots, b_m, 1]$$

for some $b_1, \ldots, b_m \ge 1$. Then there exists a simple planar graph G = (V, E) and edge $e \in E$, such that

 $\tau(G-e) = t$, $\tau(G/e) = u$, and $|V| = b_1 + \ldots + b_m + 2$.

Then the graph needed for Theorem 1.7 is G - e from Lemma 2.3. (This explains the one "missing" edge in Figure 1.1.)

Proof of Lemma 2.3. Let G = (V, E) be a connected simple planar graph with a marked edge $e = (x, y) \in E$. For $k \ge 1$, consider a marked graph

$$\Upsilon^k(G,e) := \Phi^k\left(\Psi(G,e)\right)$$

obtained by adding an edge parallel to e, and subdividing e with k vertices. This operation is equivalent to adding a path with k + 1 edges connecting x and y, and marking one of the edges in the new path (see Figure 2.1).

Note that marked graph $(H, f) := \Upsilon^k(G, e)$ is a connected simple planar graph by the description above. Combining Lemma 2.1 and Lemma 2.2, we have:

$$\boldsymbol{v}(H,f) = \boldsymbol{v}(G,e) \cdot \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 1 & k \\ 0 & 1 \end{pmatrix}.$$

Here is a trivial observation (one made already in a completely unrelated setting, namely [KNW17, p. 1042]):

$$\left(\begin{array}{cc}1&0\\1&1\end{array}\right)\left(\begin{array}{cc}1&k\\0&1\end{array}\right)=\left(\begin{array}{cc}0&1\\1&1\end{array}\right)\left(\begin{array}{cc}0&1\\1&k\end{array}\right).$$

Thus

(2.3)
$$\boldsymbol{v}(H,f) = \boldsymbol{v}(G,e) \cdot \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 0 & 1 \\ 1 & k \end{pmatrix}.$$

Now let P_1 be a marked graph consisting of a single marked edge u connecting two vertices. Then $v(P_1, u) = (0, 1)$. Let (G, e) be a marked graph given by

$$(G,e) := \Upsilon^{b_1} \Upsilon^{b_2} \cdots \Upsilon^{b_m} (P_1, u)$$

see Figure 2.2 (and compare to Figure 1.1).



FIGURE 2.2. Marked graph $(G, e) = \Upsilon^3 \Upsilon^1 \Upsilon^2 \Upsilon^4(P_1, u)$.

From above, G = (V, E) is a connected simple planar graph. By induction, we have:

$$|V| = b_1 + \ldots + b_m + 2.$$

Similarly, equation (2.3) gives by induction:

(2.4)
$$\boldsymbol{v}(G,e) = (0,1) \cdot \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 0 & 1 \\ 1 & b_m \end{pmatrix} \cdots \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 0 & 1 \\ 1 & b_1 \end{pmatrix}.$$

There is a well-known general association between the matrix product:

(2.5)
$$\begin{pmatrix} 0 & 1 \\ 1 & a_{\ell} \end{pmatrix} \cdots \begin{pmatrix} 0 & 1 \\ 1 & a_{1} \end{pmatrix} = \begin{pmatrix} * & * \\ t & u \end{pmatrix}$$

and the continued fraction expansion:

$$\frac{t}{u} = [0; a_1, a_2, \dots, a_\ell].$$

It follows that

$$\boldsymbol{v}(G, e) = (\tau(G - e), \tau(G/e)) = (t, u)$$

where gcd(t, u) = 1 and t/u has continued fraction expansion as in (2.2). This completes the proof.

3. BOURGAIN-KONTOROVICH TECHNOLOGY

3.1. Thin Orbits. Given A > 0, let \mathcal{R}_A denote the set of fractions of interest,

(3.1)
$$\mathcal{R}_A := \left\{ \frac{t}{n} = [b_1, 1, b_2, 1, \dots, b_m, 1] : 1 \le b_i \le A \text{ for all } 1 \le i \le m \right\}$$

and let \mathcal{N}_A be the set of its numerators,

$$\mathcal{N}_A := \{t \in \mathbb{N} : \frac{t}{u} \in \mathcal{R}_A \text{ for some } u > t, \ \gcd(t, u) = 1\}.$$

As already suggested by (2.5), a starting point of the analysis is to convert questions about continued fractions into ones about "Thin Orbits," so that we may use techniques from the Orbital Circle Method (see [Kon13, §5]). To this end, it is natural to introduce the matrix semigroup

$$\Gamma_A^{(0)} := \langle \begin{pmatrix} 0 & 1 \\ 1 & a \end{pmatrix} : 1 \le a \le A \rangle^+ \cap \operatorname{SL}(2, \mathbb{Z})$$

which corresponds to all even-length words in matrices of the kind appearing in (2.5), and having all "partial quotients" a bounded by A.

For the application to \mathcal{R}_A in (3.1), we introduce the corresponding semigroup

$$\Gamma_A := \left\langle \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 0 & 1 \\ 1 & b \end{pmatrix} : 1 \le b \le A \right\rangle^+,$$

so that $\binom{*}{t}{u} \in \Gamma_A$ if and only if $t/u \in \mathcal{R}_A$. Then to access the numerator t in \mathcal{N}_A , we simply take the inner product:

(3.2)
$$\mathcal{N}_A = \langle \boldsymbol{v}_1 \cdot \boldsymbol{\Gamma}_A, \boldsymbol{v}_2 \rangle,$$

where $v_1 := (0, 1)$ and $v_2 = (1, 0)$.

Remark 3.1. While the Diophantine Conjecture 1.8 deals with numerators, Zaremba's Conjecture 1.11 is about denominators. In light of (2.5), this merely amounts to replacing $v_2 = (1,0)$ above with $v_2 = (0,1)$.

3.2. Local-Global Principles and Admissibility. The general Bourgain–Kontorovich machinery proves results of the following kind, of which Theorem 1.13 is a special case. Let $\Gamma \subseteq \Gamma_A^{(0)}$ be a finitely generated semigroup, and for fixed $v_1, v_2 \in \mathbb{Z}^2$, consider the set

$$\mathcal{O} := \langle \boldsymbol{v}_1 \boldsymbol{\Gamma}, \boldsymbol{v}_2 \rangle.$$

We say that an integer n is *admissible* if

 $n \in \mathcal{O} \pmod{q}$ for all integers q.

Let $\mathcal{A} = \mathcal{A}(\Gamma, v_1, v_2)$ denote the set of admissible integers, and note that $\mathcal{O} \subseteq \mathcal{A}$. Then the even more precise density-one version of Theorem 1.13 is the following (see [BK14, Theorem 1.8]).

Theorem 3.2. There is a constant $\delta_0 < 1$ such that, if the Hausdorff dimension δ_{Γ} of the limit set of Γ satisfies $\delta_{\Gamma} > \delta_0$, then \mathcal{O} contains a density one set of its admissible integers:

(3.3)
$$\frac{\left|\mathcal{O}\cap\{1,\ldots,N\}\right|}{\left|\mathcal{A}\cap\{1,\ldots,N\}\right|} \to 1 \quad as \quad N \to \infty.$$

This result is interpreted as an asymptotic "local-global" principle for the orbit \mathcal{O} . See §5.6 for further discussion on this result and the rate of convergence in (3.3). To adapt it to our setting, the first question is to determine the set of admissible integers. In practice, this is quite simple, and in this case follows from the following.

Lemma 3.3. Let $A \geq 2$. Then $\Gamma_A \pmod{q} = \operatorname{SL}(2, \mathbb{Z}/q\mathbb{Z})$, for all $q \geq 2$.

Proof. Since reduction mod q turns Γ_A into a finite group (in which every element has finite order), it follows that the reduction mod q of the semigroup Γ_A is equal to the reduction mod q of the group generated by Γ_A . Let $M_b := \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 0 & 1 \\ 1 & b \end{pmatrix}$. Note that $M_1^{-1} \cdot M_2 = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$, and $M_1 M_2^{-1} M_1 = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$. Therefore, the group generated by M_1 and M_2 is all of SL(2, \mathbb{Z}). Hence the mod q reductions of Γ_A and SL(2, \mathbb{Z}) are the same, and equal to SL(2, $\mathbb{Z}/q\mathbb{Z}$).

It follows immediately from (3.2) and Lemma 3.3 that all integers are admissible for \mathcal{N}_A . Then the Main Diophantine Theorem 1.14 follows on applying Theorem 3.2 with $\Gamma = \Gamma_A$. For the reader's convenience, we give a sketch of some of the ingredients going into the proof of Theorem 3.2, referring the reader to, e.g., [Kon13] and the original references for more details.

3.3. Orbital circle method. To show that an integer $n = \Theta(N)$ belongs to \mathcal{O}^3 , it suffices to show that it has a positive number of representatives, where the representation of a number is, roughly speaking, the function:

(3.4)
$$R_N(n) := \sum_{\gamma \in B_N} \mathbf{1}_{\{\langle v_1 \gamma, v_0 \rangle = n\}}.$$

³Note that in number theory and references above, the notation $n \simeq N$ is used here; like $n = \Theta(N)$, it means that $N \ll n \ll N$.

Here B_N is the intersection of Γ with a ball of radius N in $SL(2, \mathbb{R}) \subset \mathbb{R}^4$ with respect to a fixed archimedean norm. The Fourier transform of R_N is the exponential sum

$$S_N(heta) := \sum_{\gamma \in B_N} \exp(heta \langle v_1 \gamma, v_0
angle),$$

so that

(3.5)
$$R_N(n) = \int_{\mathbb{R}/\mathbb{Z}} S_N(\theta) \exp(-n\theta) d\theta.$$

The "circle method" is the idea akin to signal processing that the above sum may sometimes be estimated effectively by decomposing θ into the so-called major and minor arcs, the former corresponding to the signal, and the latter meant to represent noise. The major arcs are values of θ that are "close" to fractions with "small" denominators (for suitable meanings of these words), and minor arcs are simply the complement. This leads to a decomposition (see [BK14, (4.25)])

$$R_N(n) = \mathcal{M}_N(n) + \mathcal{E}_N(n),$$

where the "main" term $\mathcal{M}(n)$ is the contribution to (3.5) coming from θ in the major arcs, and $\mathcal{E}_N(n)$ representing the "error" term.

One shows (see, e.g., [BK14, Theorem 4.28]) that the main term is at least of the expected order:

(3.6)
$$\mathcal{M}_N(n) > C \mathfrak{S}(n) \frac{|B_N|}{N} \quad \text{for some} \quad C > 0,$$

where $\mathfrak{S}(n)$ is the "singular series," which vanishes if n is not admissible, and is roughly of order 1 elsewhere. If one could prove that for individual $n = \Theta(N)$, the error term is significantly smaller,

(3.7)
$$\mathcal{E}_N(n) \stackrel{?}{=} o\left(\frac{|B_N|}{N}\right),$$

then the representation number is asymptotically positive, and hence the full local-global principle holds: every sufficiently large admissible number is represented. Unfortunately, as discussed in §5.6, there are settings which are analytically indistinguishable from this one, for which there are arbitrarily large admissible integers with no representatives; and hence (3.7) is false as the error term *must* be exactly as large as the main term infinitely often. Instead, what one manages to prove in practice is an L^2 -averaged version of (3.7), namely, a statement roughly of the form (see [BK14, Theorem 7.1] and [Hua15, Theorem 9.1]):

(3.8)
$$\sum_{n \le N \le N} |\mathcal{E}_N(n)|^2 = o\left(N \cdot \frac{|B_N|^2}{N^2}\right).$$

An elementary Cauchy–Schwartz type argument concludes Theorem 3.2 from (3.6) and (3.8).

The analysis of both \mathcal{M}_N and \mathcal{E}_N is long and difficult. The former relies on "expander graph" properties for certain families Cayley graphs in congruence towers, and the thermodynamic formalism of Ruelle transfer operators. The analysis of the error term \mathcal{E}_N uses more elementary tools, but involves a rather complicated process coming from the theory of cancellation in exponential sums. It only succeeds to prove (3.8) if the Hausdorff dimension δ_{Γ} of the limit set of Γ exceeds the threshold value δ_0 .

If all one wants is a "positive proportion" result, rather than "density one," then one need not introduce major arcs and can avoid proving (3.6); see §8 in the original **arXiv** version of [BK14].⁴

⁴arxiv.org/abs/1107.3776v1

3.4. Improvements on δ_0 . As mentioned in §1.6, the error term treatment in Frolenkov–Kan [FK14] and series of papers by Kan [Kan15]–[Kan21] improves on the original minor arcs analysis in [BK14], leading to improvements in δ_0 . We give here a few comments on at least some of their ideas, as well as the modifications that would be needed to improve their positive proportion results to density one.

In order to establish the error bound (3.8), it is necessary to modify the representation number R_N defined in (3.4) to create bilinear (in fact, multilinear) forms, as follows. As discussed in [BK14, §3.4], one replaces the full norm ball B_N in Γ by a suitable "ensemble," $\Omega_N \subset B_N$ which, in the original argument, has something like the structure

$$\Omega_N \approx \widetilde{B}_{N^{1/2}} \cdot \widetilde{B}_{N^{1/4}} \cdot \widetilde{B}_{N^{1/8}} \cdots ,$$

where \widetilde{B}_N are some large "modified balls," $\widetilde{B}_N \subset B_N$. We can view this ensemble diagrammatically as follows:



Using this ensemble as is allows one to prove Zaremba's conjecture for a positive proportion of denominators, but the modified balls are not "spectral," that is, not suitable for the major arcs analysis needed for establishing (3.6). Therefore one adds a tiny spectral sliver to the first ball (for the precise statement, which is rather technical, see [BK14, (3.37)]), which can be expressed diagrammatically as:



This makes the modified ensemble amenable to analysis.

The first innovation in [FK14] is to refine the ensemble Ω_N even further, to one of the following rough diagrammatic shape (see [FK14, (3.5)] and [Kan21, §6]):



This ensemble is again missing a spectral component, but it and its further refinements lead to the aforementioned improvements on δ_0 , and eventually to Theorem 1.15. The idea in Huang [Hua15] combines these improvements with the major arcs analysis of [BK14], by reinserting another spectral sliver in the largest component (see [Hua15, (3.51)]), leading to an ensemble of the shape



This again allows one to apply the major arcs analysis (see [Hua15, Theorem 5.5]). In combining the minor arcs analysis in [Kan21] with the major arcs from [BK14, Hua15], one can show the density one result in Theorem 3.2, with the improved values of the parameter δ_0 .

4. Estimating the Hausdorff dimension

4.1. Proof of the Main Dimension Theorem 1.17. A well-known strategy to compute Hausdorff dimensions ϑ_A comes from techniques in thermodynamics, namely to consider certain Ruelletype transfer operators, as follows. For integers $b \ge 1$, let $T_b : [0,1] \to [0,1]$ be the operation corresponding to one step in (2.4), that is, translation by b, inversion, translation by 1 and inversion again,

$$T_b(x) := \frac{1}{1 + \frac{1}{b+x}} = \frac{b+x}{1+b+x}.$$

This has derivative

$$T'_b(x) = \frac{1}{(1+b+x)^2}.$$

For A > 1 (we will take A = 110) and $s \in (0, 1)$ (a parameter used to estimate the Hausdorff dimension), we define the *pressure function*, $P = P_A$, given by

$$P(s) := \lim_{n \to \infty} \frac{1}{n} \log \left(\sum_{1 \le b_1, \dots, b_n \le A} \left| \left(T_{b_1} \circ \dots \circ T_{b_n} \right)'(0) \right|^s \right)$$

It follows from the work (in much greater generality) of Ruelle [Rue82] (see also Bowen [Bow79]), that the pressure function P is a monotone decreasing function of s, and has a unique zero at the Hausdorff dimension $s = \vartheta_A$.

It is difficult in practice to rigorously approximate the dimension from P alone, and one introduces the following *transfer operator* $\mathcal{L}_s = \mathcal{L}_{s,A}$; given a function $f : [0,1] \to \mathbb{R}$, we define:

(4.1)
$$[\mathcal{L}_s f](x) := \sum_{b=1}^A |T'_b(x)|^s (f \circ T_b)(x).$$

A suitable space on which the transfer operator \mathcal{L}_s can act is, e.g., the Banach space $C^1([0,1])$ of Hölder continuous functions with norm $||f|| := ||f||_1 + ||f||_{\infty}$, where

$$||f||_1 := \sup_{0 \le x \ne y \le 1} \frac{|f(x) - f(y)|}{|x - y|}$$

The connection to the pressure function and Hausdorff dimension is given by the Ruelle–Perron– Frobenius theorem, which states that \mathcal{L}_s acting on $C^1([0,1])$ has a unique maximal eigenvalue $\lambda_s = e^{P(s)}$, and the corresponding eigenfunction

(4.2)
$$h_s = h_{s,A} \in C^1([0,1])$$

is strictly positive on [0, 1]. A nice discussion of this theory is given by Pollicott–Vytnova in $[PV22, \S2]$, where the authors develop a relatively simple and easily implementable procedure to give rigorous estimates for Hausdorff dimensions in a wide variety of settings.

Of particular importance to us is the mechanism for producing lower bounds for dimensions, using certain min-max inequalities.

Lemma 4.1. Suppose that there exists some s > 0 and f a polynomial which is positive on [0, 1], so that

$$(4.3) \qquad \qquad [\mathcal{L}_s f](x) > f(x),$$

for all $x \in [0, 1]$. Then $\vartheta_A > s$.

Intuitively, if we want values of s with P near zero, that corresponds to eigenvalues $\lambda_s = e^{P(s)}$ of \mathcal{L}_s near 1. Since λ_s is the largest eigenvalue, (4.3) is evidence that $\lambda_s > 1$.

Proof. Since [0,1] is compact, (4.3) implies that there exists some a > 1 so that $af(x) \leq [\mathcal{L}_s f](x)$ holds for all $x \in [0,1]$. By [PV22, Lemma 3.1(1)], this implies that $a \leq e^{P(s)}$, or $P(s) \geq \log a > 0$. But since P is monotonically decreasing and $P(\vartheta_A) = 0$, we have that $s < \vartheta_A$. This reduces the problem to one of producing, for $s = \delta_0 = 0.775$, say, a sufficiently large value of A and a positive "test function" f for which (4.3) holds. We are grateful to Polina Vytnova for explaining to us the procedure from [PV22] for finding such functions in practice, namely to simply try to approximate the (positive) eigenfunction h_s in (4.2) via Lagrange-Chebyshev interpolation. (In principle, almost any interpolation scheme should work, but this one seems to perform particularly well in this application.)

To this end, fix an order of approximation, N (we will take N = 5), and for $1 \le j \le N$, define the Chebyshev nodes

$$y_j := \frac{1}{2} \left(\cos \left(\frac{2j-1}{2N} \pi \right) + 1 \right) \in [0,1],$$

and corresponding Lagrange interpolation polynomials

$$\ell_j(x) := \prod_{\substack{1 \le k \le N \\ k \ne j}} \frac{x - y_k}{y_j - y_k}.$$

These have the orthonormality property that $\ell_j(y_k) = 1$ if j = k and 0 otherwise. For N large, the nodes become dense in [0, 1], and the Lagrange polynomials span some "significant" subspace of our Banach space $C^1([0, 1])$.

To find a linear combination of the ℓ_j 's which approximates the eigenfunction h_s of the transfer operator \mathcal{L}_s , simply hit them all with \mathcal{L}_s , and evaluate at the nodes. This should excite a harmonic near h_s , giving the approximation. In practice, this is achieved as follows. We evaluate the $N \times N$ matrix, M, say, having entries

$$M_{j,k} := [\mathcal{L}_s \ell_j](y_k),$$

 $(1 \leq j, k \leq N)$, and compute an eigenvector $v_s = (v_s^{(1)}, \ldots, v_s^{(N)})$ corresponding to its largest eigenvalue. By the orthonormality relations, this v_s suggests which linear combination of the ℓ_j should have a high correlation with h_s ; that is, we should try the test function:

$$f_s := \sum_{1 \le k \le N} v_s^{(j)} \ell_j.$$

If it turns out that f_s is positive on [0, 1], and that (4.3) is satisfied, then we learn from Lemma 4.1 that s is a lower bound for ϑ_A . If not, we should increase N (or if that continues to fail, increase A), and hope to get lucky. In [PV22, Prop. 3.10], Pollicott and Vytnova prove that, if in fact $\vartheta_A > s$, then such an f_s will eventually be found by taking N large enough.



FIGURE 4.1. (a) The test function f_s in (4.4), and (b) difference $\mathcal{L}_s f_s - f_s$ on [0, 1].

Setting $s = \delta_0 = 0.775$, N = 5, and A = 110, (chosen for the simple reason that A = 100 fails, even with large values of N), the procedure above gives the eigenvector

 $v_s := (0.3798483, 0.3992862, 0.4366593, 0.4841648, 0.5207676),$

and corresponding test polynomial f_s given by:

$$(4.4) f_s(x) := 0.0121844x^4 - 0.0513245x^3 + 0.116313x^2 - 0.225988x + 0.526229$$

For $x \in [0, 1]$, it can be verified that $f_s(x) > 0.3$, see Figure 4.1(*a*). Moreover, it can be verified that on [0, 1], the difference $\mathcal{L}_s f_s - f_s$ exceeds $7 \times 10^{-5} > 0$, see Figure 4.1(*b*). This confirms (4.3), and hence proves the Main Dimension Theorem 1.17.

4.2. Maximal dimension ϑ_{∞} . We conclude this section by explaining Remark 1.18. The Cantor set \mathfrak{C}_{∞} defined in (1.8) is contained in the fractal

$$\bar{\mathfrak{C}} := \{ [b_1, 1, b_2, 1, \ldots] : b_i \ge 1, \quad \forall i \}.$$

In particular, the former is the intersection of the latter with the set of badly approximable numbers (that is, those with bounded partial quotients). Hence the Hausdorff dimension

$$\bar{\vartheta} := \operatorname{Hdim}(\bar{\mathfrak{C}})$$

is an upper bound for ϑ_{∞} . The method described in §4.1 cannot be applied directly to the problem of estimating $\bar{\vartheta}$, in particular because the corresponding transfer operator (4.1) would now be an infinite sum. Instead, here is an elegant mechanism (again, generously explained to us by Polina Vytnova) for establishing the estimate

(4.5)
$$\vartheta_{\infty} \le \vartheta < 0.799.$$

Let $f(x) = \sum_{n=0}^{N} a_n (x-1)^n$ be a polynomial (expanded about x = 1). Then the transfer operator acting on f can be expressed as

$$\begin{aligned} [\mathcal{L}_s f](x) &:= \sum_{b=1}^{\infty} |T_b'(x)|^s \cdot f(T_b(x)) \ = \ \sum_{b=1}^{\infty} \frac{1}{(1+b+x)^{2s}} \cdot f\left(1 - \frac{1}{1+b+x}\right) \\ &= \ \sum_{b=1}^{\infty} \frac{1}{(1+b+x)^{2s}} \sum_{n=0}^{N} a_n \left(\frac{-1}{1+b+x}\right)^n \ = \ \sum_{n=0}^{N} a_n (-1)^n \zeta(2s+n,2+x) \,, \end{aligned}$$

which is now again a finite sum. Here

$$\zeta(s,x) \; := \; \sum_{b=0}^{\infty} \; \frac{1}{(b+x)^s}$$

is the Hurwitz zeta function. An analogous statement to Lemma 4.1 holds with inequalities reversed: if there is a polynomial f which is positive on [0, 1] such that

$$[\mathcal{L}_s f](x) < f(x)$$

for all $x \in [0, 1]$, then $\overline{\vartheta} < s$.

Now we proceed as before. Set s = 0.799, and N = 5, and using the same Chebyshev–Lagrange polynomials, compute the matrix $M_{j,k} = [\mathcal{L}_s(\ell_j)](y_k)$. (One needs to first extract the polynomials' coefficients.) Its largest eigenvalue has eigenvector

$$v_s = (0.3820795, 0.4007878, 0.4369026, 0.4830608, 0.5187994)$$

and the corresponding test function $f_s = \sum v_s^{(j)} \ell_j$ is computed to be

$$(4.6) f_s(x) = 0.0123381x^4 - 0.0517567x^3 + 0.116202x^2 - 0.221186x + 0.524143x^4 - 0.0517567x^3 + 0.0517567x^3 + 0.0517567x^3 + 0.0517567x^4 - 0.0517567x^3 + 0.0517567x^3 + 0.0517567x^4 - 0.0517567x^4 - 0.0517567x^4 - 0.0517567x^5 - 0.0517567x^4 - 0.0517567x^4 - 0.0517567x^5 - 0.051757x^5 - 0.05175x^5 -$$

This function exceeds 0.3 on [0,1], and $\mathcal{L}_s f_s - f_s$ is less than -0.0002 on [0,1], see Figure 4.2. This confirms (4.5).



FIGURE 4.2. (a) The test function f_s in (4.6), and (b) the difference $\mathcal{L}_s f_s - f_s$.

5. FINAL REMARKS AND OPEN PROBLEMS

5.1. In notation of Main Theorem 1.1, one can ask for the upper bound on the number $|\mathcal{T}(n)|$ of distinct values of $\tau(G)$ over all simple planar graphs with n vertices. The bound $|\mathcal{T}(n)| < 8^n$ is straightforward; indeed, note that for every planar simple graph G = (V, E) on n vertices, we have $\tau(G) < 2^{|E|} < 8^n$, since $|E| \leq 3n - 6$. Writing $\tau(G)$ as a product of eigenvalues of the Laplacian matrix and using the AM–GM inequality shows that $\tau(G) < 6^n$, see [Gri76, Eq. (1)].

Skipping over a series of further improvements, the best known upper bound for the number of spanning trees in a simple planar graph is $\tau(G) = O(5.2852^n)$ given in [BS10]. The optimal base of exponent is probably much smaller, see below. However, since

$$\lim_{n \to \infty} \left(\max \mathcal{T}(n) \right)^{\frac{1}{n}} \ge 5.0295$$

(see [Rib606, §5.1]), there is a limit in the approach above.

In a similar vein, one can ask for an explicit lower bound for the constant c in Theorem 1.1. Our proof of Theorem 1.1 shows one can choose any value of c strictly less than $\phi^{2/111} \approx 1.0087$ for the statement of Theorem 1.1 to hold, where $\phi = \frac{1+\sqrt{5}}{2} \approx 1.618$ is the golden ratio. We made no effort in optimizing this constant. We note, however, that any lower bound achievable by our approach cannot exceed ϕ .

5.2. In view of Theorem 1.1, we make the following natural

Conjecture 5.1. There is a limit $c := \lim_{n \to \infty} (|\mathcal{T}_n|)^{\frac{1}{n}}$.

If Conjecture 5.1 holds, then $c \ge \phi^{2/111} \approx 1.0087$ by Theorem 1.1 and the aforementioned lower bound in §5.1.

5.3. Denote by $\mathcal{T}'(n)$ the set of spanning tree numbers $\tau(G)$ as G ranges over all simple graphs on n vertices, i.e. without the planarity assumptions [Sed69, Sed70]. The sequence $\{|\mathcal{T}'(n)|\}$ starts as

1, 1, 2, 5, 16, 65, 386, 3700, 55784, 1134526, 27053464,

see [OEIS, A182290]. By Theorem 1.1 and Cayley's formula, we have:

$$c^n \leq |\mathcal{T}(n)| \leq |\mathcal{T}'(n)| \leq \tau(K_n) = n^{n-2}.$$

Note that the expected number of spanning trees in a uniform random labeled graph on n vertices is rather large, and equal to $n^{n-2}/2^{n-1}$. Note also that the number of unlabeled graphs on n vertices is even larger, and asymptotically equal to $2^{\binom{n}{2}}/n!$, see e.g. [Noy15, §6.9.2] and [OEIS, A000088]. These suggest the following:⁵

⁵In fact, the first few values of the sequence suggest a stronger bound: $|\mathcal{T}'(n)| = e^{(1-o(1))n \log n}$.

Conjecture 5.2. $|\mathcal{T}'(n)| = e^{\Omega(n \log n)}$.

Following Sedláček [Sed70], let $\alpha'(t)$ denote the minimal number of vertices of a simple (and not necessarily planar) graph with exactly t spanning trees. Clearly, we have $\alpha'(t) \leq \alpha(t)$, but potentially $\alpha'(t)$ is much smaller. Cayley's formula gives a natural lower bound $\alpha'(t) =$ $\Omega(\log t/\log \log t)$ in this case. Sedláček originally conjectured that $\alpha'(t) = o(\log t)$ [Sed70]. Stong's Theorem 1.3 remains the best known upper bound for $\alpha'(t)$.

5.4. By the matrix-tree theorem, the number of spanning trees of a simple graph is a determinant of an integral matrix with off-diagonal entries in $\{0, -1\}$. The set of possible determinant values for various classes of combinatorial matrices is interesting in its own right and closely related to Hadamard's maximal determinant problem and the determinant spectrum problem which remain unresolved. Given the state of art in these two problems, we speculate that Conjecture 1.6 is likely to be more accessible than Conjecture 5.2. We refer to [Shah22] and a discussion in [CP24b, §6.10] for further references.

5.5. In [CP24a, CP24c], the authors used the average cases analysis of the sum of partial quotients of continued fractions, applying sharp bounds of Larcher [Lar86] and Rukavishnikova [Ruk11] to get bounds on $\beta(t)$ discussed in the introduction. While these bounds hold for all t, the resulting extra $O(\log \log t)$ factors are unavoidable with these tools. In the terminology of [YK75] (see also [Knu98, §4.5.3]), the sum of partial quotients is the number of steps of the *subtraction algorithm*, the original (classical) version of the Euclidean algorithm for finding the greatest common divisor that uses only subtractions instead of divisions. Finally, in [CP24c], it was shown that Zaremba's Conjecture 1.11 implies that $\beta(t) = \Theta(\log t)$. The authors also produced a series of weaker conjectures, all of which remain open.

5.6. There is a further refinement in the asymptotic formula in (3.3). Namely, the work of [Bou18] and [MOW19] shows that the asymptotic value 1 in Theorem 3.2 is approached with a power savings: LHS = $1 + O(N^{-\eta})$ for some $\eta > 0$. We further remark that a recent work of Rickards and Stange [RS24], shows that it is possible to exhibit finitely generated semigroups $\Gamma \subset SL_2(\mathbb{Z})$, and base vectors $v_0, v_1 \in \mathbb{Z}^2$, such that none of the elements of the orbit $\mathcal{O} := \langle v_1 \Gamma, v_0 \rangle$ are, e.g., perfect squares. This disproves a conjecture by Bourgain and Kontorovich [BK18, Conj. 1.11], and shows that there can be a Brauer–Manin type reciprocity obstruction, meaning that a power savings asymptotic error is best possible in this generality. The full local-global principle, that, say, \mathcal{O} contains every sufficiently large integer, may well be true for our Γ_A (for A sufficiently large); this would imply the Diophantine Conjecture 1.8. Unfortunately, the general orbital circle method cannot distinguish this setting from the similar setting where the statement is false. Therefore, these techniques will not, without significant further ideas, be able to upgrade Theorem 1.5 to the full Conjecture 1.4.

5.7. In Zaremba's Conjecture 1.11, it is known that A = 4 is not enough (take u = 6 or 54). Hensley (1996) conjectured that one can take A = 2 for d large enough. Zaremba's conjecture is known to hold for integers of the form $2^m 3^n$ and for sufficiently large powers of all primes, where the constant A can depend on the prime, see [Nid86, Shu24] for details. We refer to [BPSZ14, §6.2] for an elegant presentation of the 2^m case.

Acknowledgements. We are thankful to Milan Haiman, Noah Kravitz, Ilya Shkredov and Nikita Shulga for interesting discussions and helpful remarks, and to Oliver Jenkinson for emailing us [Jen04]. We are very grateful to Polina Vytnova for patiently explaining to us the key ideas in [PV22] for rigorously computing Hausdorff dimensions.

SHC was supported by NSF grant DMS-2246845. AK was supported by NSF grant DMS-2302641, BSF grant 2020119 and a Simons Fellowship. IP was supported by NSF CCF-2302173. This paper was written

when AK was visiting Princeton University and IP was a member at the Institute of Advanced Study in Princeton, NJ. We are grateful for the hospitality.

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(Swee Hong Chan) DEPARTMENT OF MATHEMATICS, RUTGERS UNIVERSITY, PISCATAWAY, NJ 08854. Email address: sweehong.chan@rutgers.edu

(Alex Kontorovich) DEPARTMENT OF MATHEMATICS, RUTGERS UNIVERSITY, PISCATAWAY, NJ 08854. *Email address*: alex.kontorovich@rutgers.edu

(Igor Pak) DEPARTMENT OF MATHEMATICS, UCLA, LOS ANGELES, CA 90095. Email address: pak@math.ucla.edu