The Number of Inversions and the Major Index of Permutations are Asymptotically Joint-Independently-Normal (Second Edition)

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Abstract: We use *recurrences* (alias *difference equations*) to prove that the two most important *permutation statistics*, namely the *number of inversions* and the *major index*, are asymptotically joint-independently-normal. We even derive more-precise-than-needed asymptotic formulas for the (normalized) mixed moments.

Preface to the Second Edition

In addition to the considerable interest of the results proved in this article, and the even greater interest of the *methodology*, this article is a landmark case in scholarly publishing. After the first version of this article was *outright* rejected by an *anonymous* referee of the *Proceedings of the American Mathematical Society*, because *too many details were left to the reader* (he or she didn't give us a chance to write a new version with more details), we decided to solicit *nine* non-anonymous reports from world-class experts, assigning them specific parts. The division of labor (quite a few of them did over and above of what we asked them to, and refereed everything), and the full reports, on the first edition, can be gotten from the webpage of this article:

`http://www.math.rutgers.edu/~zeilberg/mamarim/mamarimhtml/invmaj.html` ,

already mentioned in footnote 1. With a few exceptions of stylistic suggestions that we preferred not to adopt, we incorporated *all* of their (excellent!) suggestions. The first edition is also availabe there, for the record.

In order to be faithful to the original version, we have included in square brackets, and smaller font, the extra explanations demanded by the referees. We thank them profusely, and we now

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Exclusively published in the Personal Journal of Shalosh B. Ekhad and Doron Zeilberger `http://www.math.rutgers.edu/~zeilberg/pj.html` and `arxiv.org`. DZ is hereby offering $1000 to the first person to point out a serious flaw in the argument, that would irreparably invalidate the proof. All the meta-mathematical opinions expressed here are those of DZ.

believe that the *formal correctness* and *clarity* far exceeds %99.99 of the articles published in (anonymous!) "peer"-reviewed mathematical journals. Such journals, even electronic ones, will soon become obsolete, together with their *pompous* "editors" and *anonymous* referees. Instead, the present model of author(s)-appointed refereeing and self-publishing, in the authors' personal websites and the *arxiv*, with all the referee reports made public, and the referees acknowledged and given explicit recognition for their trouble, would become the norm, possibly with some tweaking. The present article is *exclusively* published in the Personal Journal of Shalosh B. Ekhad and Doron Zeilberger and `arxiv.org`.

Human Statistics

Human statistics are numerical attributes defined on *humans*, for example, *longevity*, *height*, *weight*, *IQ*, and it is well-known, at least empirically, that these are, each separately, *asymptotically normal*, which means that if you draw a histogram with the statistical data, it would look like a *bell-curve*. It is also true that they are usually joint-asymptotically-normal, but usually not *independently* so. But if you compute empirically the *correlation matrix*, you would get, asymptotically (i.e. for "large" populations) that they are close to being distributed according to a multivariate (generalized) *Gaussian* $exp(-Q(x_1, x_2, ...))$ with $Q(x_1, x_2, \dots)$ a certain *quadratic form* that can be deduced from the correlation matrix.

Permutation Statistics

Let our population be the set of *permutations* of $\{1, 2, \dots, n\}$. They too, can be assigned *numerical attributes*, and the great classical combinatorialist Dominique Foata (who got his *Doctorat de troisième cycle* in statistics!) coined the term *permutation statistics* for them.

The most important permutation statistic is the **number of inversions**, $inv(\pi)$, that counts the number of pairs $1 \leq i < j \leq n$ such that $\pi_i > \pi_j$ (and ranges from 0 to $n(n-1)/2$). For example, $inv(314625) = 5$, corresponding to the set of pairs $\{[1,2],[1,5],[3,5],[4,5],[4,6]\}$. It features in the definition of the determinant, and Netto proved that the *probability generating function* (the polynomial in q such that its coefficient of q^i is the probability that a uniformly-at-random n-permutation has i inversions) is given by

$$\frac{(1)(1+q)(1+q+q^2)\cdots(1+q+q^2+\ldots+q^{n-1})}{n!} = \frac{\prod_{i=1}^{n}(1-q^i)}{n!(1-q)^n} \quad .$$

The *second most* important permutation statistic is the *major index*, $maj(\pi)$, that is the sum of the places i, where $\pi_i > \pi_{i+1}$. For example, $maj(314625) = 1 + 4 = 5$, because at $i = 1$ and $i = 4$ we have descents. Major Percy Alexander MacMahon [M] famously proved that the probability generating function for the major index is also given by that very same formula. In other words the permutation statistics *inv* and *maj* are *equidistributed*. Dominique Foata [Fo] gave a lovely seminal *bijective* proof that proved the stronger statement that *inv* and *maj* are equi-distributed also when restricted to permutations ending at a given integer.

William Feller ([Fe], 3rd ed., p.257) proved that the number of inversions (and hence also the major

index) is *asymptotically normal* in the following sense. Feller easily computed the *expectation*,

$$E[inv] = m_n = n(n-1)/4 \quad ,$$

and the *variance*,

$$\sigma_n^2 = \frac{2n^3 + 3n^2 - 5n}{72} \quad .$$

If we denote by X_n the *centralized* and *normalized* random variable

$$X_n = \frac{inv - m_n}{\sigma_n} \quad ,$$

then $X_n \to \mathcal{N}$, as $n \to \infty$, in distribution, where $\mathcal{N}$ is the Gaussian distribution whose *probability density function* is $e^{-x^2/2}/\sqrt{2\pi}$.

A *computer-generated* proof, that gives much more detail about the rate of convergence to $\mathcal{N}$, can be obtained using Zeilberger's Maple package
`http://www.math.rutgers.edu/~zeilberg/tokhniot/AsymptoticMoments` that accompanies the article [Z].

So both *inv* and *maj* are *individually* asymptotically normal, but what about their *interaction*? In this article, we prove that they are asymptotically *joint-independently-normal*. In other words, defining,

$$X_n(\pi) := \frac{inv(\pi) - m_n}{\sigma_n} \quad , \quad Y_n(\pi) := \frac{maj(\pi) - m_n}{\sigma_n} \quad ,$$

we have that

$$Pr(X_n \leq s \ , \ Y_n \leq t) \to \frac{1}{2\pi} \int_{-\infty}^{s} \int_{-\infty}^{t} e^{-x^2/2 - y^2/2} \, dy \, dx \quad as \quad n \to \infty.$$

A Brief History

It all started when the great Swedish probabilist Svante Janson (member of the Swedish Academy of Science, that awards the Nobel prizes) asked Donald Knuth (one of the greatest computer scientists of all time, winner of the Turing and Kyoto prizes, among many other honors) about the *asymptotic covariance* of *inv* and *maj*. Neither of these luminaries knew the answer, so Don Knuth asked one of us (DZ). DZ didn't know the answer either, so he asked his beloved servant, Shalosh B. Ekhad, who *immediately* ([E]) produced, not just the asymptotics, but the *exact answer*! It turned out to be $n(n-1)/8$. In particular, the *correlation coefficient*, $Cov(inv, maj)/\sigma_n^2 = \frac{\frac{n(n-1)}{8}}{\frac{2n^3+3n^2-5n}{72}} = \frac{9}{2n} + O(1/n^2)$ tends to zero as n goes to infinity. It followed that in the long-run, *inv* and *maj* are practically *uncorrelated*.

But there are lots of pairs of random variables that are uncorrelated yet not independent. A convenient way to prove that X_n and Y_n are asymptotically independent (we already know that they are both normal) is to use the method of moments, and to prove that the *mixed moments*

$$M_{r,s}(n) := E[X_n^r Y_n^s] \quad ,$$

tend to the mixed moments of $\mathcal{N} \times \mathcal{N}$, as $n \to \infty$. In other words, for $r, s \geq 1$:

$$\lim_{n\to\infty} M_{2r,2s}(n) = \frac{(2r)!}{2^r r!}\frac{(2s)!}{2^s s!} \quad , \tag{EE}$$

$$\lim_{n\to\infty} M_{2r-1,2s}(n) \; = \; 0 \quad , \tag{OE}$$

$$\lim_{n\to\infty} M_{2r,2s-1}(n) \; = \; 0 \quad , \tag{EO}$$

$$\lim_{n\to\infty} M_{2r-1,2s-1}(n) \; = \; 0 \quad . \tag{OO}$$

Ekhad's brilliant approach to the Janson-Knuth question merely settled the case $r = 1, s = 1$ of (OO). Of course, because of symmetry (OE) and (EO) are trivially true (before taking the limits!, i.e. $M_{2r,2s-1}(n) \equiv 0$ and $M_{2r-1,2s}(n) \equiv 0$).

[Indeed if $com([\pi_1, \ldots, \pi_n]) := [n+1-\pi_1, \ldots, n+1-\pi_n]$ then trivially $maj(\pi)+maj(com(\pi)) = n(n-1)/2$ and $inv(\pi) + inv(com(\pi)) = n(n-1)/2$. So $Pr(\{X_n = i, Y_n = j\}) = Pr(\{X_n = -i, Y_n = -j\})$ for all i, j, and it follows that

$$M_{2r,2s-1}(n) = \sum_{i,j} i^{2r} j^{2s-1} Pr(\{X_n = i, Y_n = j\}) = \sum_{i,j} i^{2r} j^{2s-1} Pr(\{X_n = -i, Y_n = -j\}) =$$

$$\sum_{i,j} (-i)^{2r} (-j)^{2s-1} Pr(\{X_n = i, Y_n = j\}) = -\sum_{i,j} (i)^{2r} (j)^{2s-1} Pr(\{X_n = i, Y_n = j\}) = -M_{2r,2s-1}(n) \quad ,$$

so $M_{2r,2s-1}(n)$ equals to its negative, and so must vanish. The proof that $M_{2r-1,2s}(n) = 0$ is similar.]

One natural approach would be to extend Ekhad's brilliant derivation of $M_{1,1}(n)$ to the general case, and try to derive closed-form expressions for $M_{r,s}(n)$ for larger r and s. Since Ekhad's proof [E] is so brief, we can cite it here in full.

"Svante Janson asked Don Knuth, who asked me, about the covariance of inv and maj. The answer is $\binom{n}{2}/4$. To prove it, I asked Shalosh to compute the average of the quantity $(inv(\pi) - E(inv))(maj(\pi) - E(maj))$ over all permutations of a given length n, and it gave me, for $n = 1, 2, 3, 4, 5$, the values $0, 1/4, 3/4, 3/2, 5/2$, respectively. Since we know a priori[2] that this is a polynomial of degree ≤ 4, this must be it! $\square$".

Obviously this brute-brute-force approach would be hopeless for deriving polynomial expressions for the moments $M_{r,s}(n)$ for larger r and s. As we will soon see, the degree of the polynomial $M_{2r,2s}(n)$ is $3(r + s)$, so for example, in order to (rigorously) guess $M_{10,10}(n)$, we would need 31

[2] This is the old trick to compute moments of combinatorial 'statistics', described nicely in [GKP], section 8.2, by changing the order of summation. It applies equally well to covariance. Rather than actually carrying out the gory details, we observe that this is always a polynomial whose degree is trivial to bound. [Added in 2nd edition: the referees didn't find this obvious and asked for an explanation. See the bottom of page 8 and the top of page 9 in the present article.].

data points, and we would have to ask our computers to examine more than $31! > 0.822 \cdot 10^{34}$ permutations.

However, an *inspired*, still "empirical" (yet **fully rigorous**) "brute-force" approach *does* work. The first step would be to have a more efficient way to compute the moments $M_{r,s}(n)$, for *specific* n and *specific* r and s. We will do it by first designing an efficient way to generate the *probability generating function*, let's call it $G(n)(p,q)$, for the pair of statistics (inv, maj). There are beautiful closed-form expressions for $G(n)(p,1)$ and $G(n)(1,q)$ (the same one, actually, due to Netto and MacMahon, given in page 2), but no such closed-form expression seems to exist for the bi-variate generating function, so the best that we can hope for is to find a *recurrence scheme*.

A Combinatorial Interlude

Let us forget about probability for a few moments, and focus on a fast algorithm for computing

$$H(n)(p,q) := \sum_{\pi \in S_n} p^{inv(\pi)} q^{maj(\pi)} \quad ,$$

for n up to, say, $n = 50$.

[Referee Dan Romik believes that we should mention, at this point, the "explicit" formula of Roselle[R] (mentioned by Knuth[K]) in terms of a certain infinite double product for the q-exponential generating function of $H(n)(p,q)$. Romik believes that this may lead to an alternative proof, that would even imply a stronger result (a local limit law). We strongly doubt this, and DZ is hereby offering $300 for the first person to supply such a proof, whose length should not exceed the length of this article.

This was proved by Joshua Swanson in his nice article http://arxiv.org/abs/1902.06724, and he won the prize.

Referee Christian Krattenthaler believes that we should also mention the beautiful extension of Roselle's result, by Adriano Garsia and Ira Gessel [GG], handling more permutation statistics.]

Define the *weight* of a permutation π to be $p^{inv(\pi)} q^{maj(\pi)}$. Suppose that $\pi \in S_n$ ends with i, so we can write $\pi = \pi' i$, where π' is a permutation of $\{1, \ldots, i-1, i+1, \ldots n\}$.

When you chop off i from π to form π' you always lose $n - i$ inversions (that is, π' has $n - i$ fewer inversions than π). The major index, however, decreases by $n - 1$ if the last letter of π', let's call it j, is larger than i. If $j < i$ the major index does not change at all. So writing $\pi = \pi'' ji$, we have

$$inv(\pi'' ji) = inv(\pi'' j) + n - i \quad ,$$

$$maj(\pi'' ji) = \begin{cases} maj(\pi'' j), & \text{if } j < i \ ; \\ maj(\pi'' j) + n - 1, & \text{if } j > i. \end{cases}$$

Combining, we have

$$weight(\pi'' ji) = \begin{cases} p^{n-i} weight(\pi'' j), & \text{if } j < i \ ; \\ p^{n-i} q^{n-1} weight(\pi'' j), & \text{if } j > i. \end{cases}$$

So in order to compute $H(n)(p,q)$, we need to introduce the more general weight-enumerators of those permutations in S_n that end with an i. Let's call these $F(n,i)(p,q)$. In symbols:

$$F(n,i)(p,q) := \sum_{\substack{\pi\in S_n \\ \pi_n=i}} p^{inv(\pi)}q^{maj(\pi)} \quad .$$

It follows that (let's omit the arguments (p,q) from now on):

$$F(n,i) = p^{n-i}\sum_{j=1}^{i-1} F(n-1,j) + p^{n-i}q^{n-1}\sum_{j=i}^{n-1} F(n-1,j) \quad . \qquad (Fni)$$

Note that, when we chop off the last entry, i, from $\pi = \pi''ji$, $\pi''j$ is a permutation of $\{1,\ldots,i-1,i+1,\ldots,n\}$. We then "reduce" $\pi''j$ to a permutation of $\{1,\ldots,n-1\}$ by diminishing all entries larger than i by 1. Hence the summation ranges from $j=i$ to $j=n-1$ rather than from $j=i+1$ to $j=n$.

Replacing i by $i+1$ in the above equation, we have:

$$F(n,i+1) = p^{n-i-1}\sum_{j=1}^{i} F(n-1,j) + p^{n-i-1}q^{n-1}\sum_{j=i+1}^{n-1} F(n-1,j) \quad .$$

Subtracting the former equation from p times the latter we get

$$F(n,i) - pF(n,i+1) =$$

$$p^{n-i}\sum_{j=1}^{i-1} F(n-1,j) + p^{n-i}q^{n-1}\sum_{j=i}^{n-1} F(n-1,j)$$

$$-p^{n-i}\sum_{j=1}^{i} F(n-1,j) - p^{n-i}q^{n-1}\sum_{j=i+1}^{n-1} F(n-1,j)$$

$$= p^{n-i}\left(\sum_{j=1}^{i-1} F(n-1,j) - \sum_{j=1}^{i} F(n-1,j)\right) + p^{n-i}q^{n-1}\left(\sum_{j=i}^{n-1} F(n-1,j) - \sum_{j=i+1}^{n-1} F(n-1,j)\right)$$

$$= -p^{n-i}F(n-1,i) + p^{n-i}q^{n-1}F(n-1,i) = p^{n-i}(q^{n-1}-1)F(n-1,i) \quad .$$

Rearranging, we get:

$$F(n,i) = pF(n,i+1) + p^{n-i}(q^{n-1}-1)F(n-1,i) \quad for \quad 1 \le i < n \quad . \qquad (RecF)$$

We still need to specify $F(n,n)$, and for this we do need the $\sum$ symbol, namely we use Eq. (Fni) with $i=n$:

$$F(n,n) = \sum_{j=1}^{n-1} F(n-1,j) \quad . \qquad (Fnn)$$

The recurrence $(RecF)$ together with the final condition (Fnn), and the trivial *initial condition* $F(1,i) = \delta_{i,1}$, enables us to efficiently compute $F(n,i)$ for numeric (n,i), for $\{(n,i)|1 \le i \le n \le N\}$ for any finite N (not too large, but not too small either: e.g., $N=100$ is still feasible). In particular, we can compile a table of $H(n)(p,q) = F(n+1,n+1)(p,q)$, (the generating function for *all* n-permutations) for $n \le N-1$.

A crash course in multivariable enumerative probability

Suppose that you have a finite set of objects S and several *statistics* $f_1(s), \ldots, f_r(s)$. The *multivariable generating function* (weight-enumerator under the weight $x_1^{f_1(s)} \cdots x_r^{f_r(s)}$) is defined to be:

$$\sum_{s \in S} x_1^{f_1(s)} \cdots x_r^{f_r(s)} \quad .$$

Suppose that you pick an element $s \in S$ *uniformly at random* and you want the multivariable generating function such that the coefficient of $x_1^{a_1} \cdots x_r^{a_r}$ would give you the probability that $f_1(s) = a_1, \ldots, f_r(s) = a_r$. It is given by:

$$P(x_1, \ldots, x_r) = \frac{1}{|S|} \sum_{s \in S} x_1^{f_1(s)} \cdots x_r^{f_r(s)} \quad .$$

The *expectations*, $\bar{f}_1, \ldots, \bar{f}_r$ are simply

$$\bar{f}_i = \left(\frac{\partial}{\partial x_i} P \right) (1, \ldots, 1) \quad .$$

The *centralized* probability generating function is

$$\tilde{P}(x_1, \ldots, x_r) = \frac{P(x_1, \ldots, x_r)}{x_1^{\bar{f}_1} \cdots x_r^{\bar{f}_r}} \quad .$$

The mixed moments (about the mean) of the statistics $f_1(s), \ldots, f_r(s)$ are defined by

$$Mom[a_1, \ldots, a_r] = \frac{1}{|S|} \sum_{s \in S} (f_1(s) - \bar{f}_1)^{a_1} \cdots (f_r(s) - \bar{f}_r)^{a_r} \quad .$$

Often it is more convenient to consider the mixed *factorial moments*, using the combinatorial "powers" $z^{(r)} := z(z-1)\cdots(z-r+1)$, better known as the *falling-factorials*. The mixed factorial moments are defined analogously by

$$FM(a_1, \ldots, a_r) = \frac{1}{|S|} \sum_{s \in S} (f_1(s) - \bar{f}_1)^{(a_1)} \cdots (f_r(s) - \bar{f}_r)^{(a_r)} \quad .$$

Once you know the FM's for all $a_1, \ldots, a_r \leq M$, you can easily figure out the Mom's (for $a_1, \ldots, a_r \leq M$), using Stirling numbers of the second kind (see [Z]). It is well-known and easy to see that one can just as well use the method of factorial moments in order to prove asymptotic independence. In other words, it would suffice to prove the analogs of $(EE), (OE), (EO), (OO)$ with the moments $E[X_n^r Y_n^s]$ replaced by the factorial moments $E[X_n^{(r)} Y_n^{(s)}]$.

The best way to compute $FM(a_1, \ldots, a_r)$ for all $0 \leq a_1, \ldots, a_r \leq M$, for some fixed positive integer M, is via the *Taylor expansion* of $\tilde{P}(x_1, \ldots, x_r)$ around $(x_1, \ldots, x_r) = (1, \ldots, 1)$, or equivalently, the *Maclaurin expansion* of $\tilde{P}(1 + x_1, \ldots, 1 + x_r)$

$$\tilde{P}(1 + x_1, \ldots, 1 + x_r) = \sum_{\alpha_1, \ldots, \alpha_r \geq 0} \frac{FM(\alpha_1, \ldots, \alpha_r)}{\alpha_1! \cdots \alpha_r!} x_1^{\alpha_1} \cdots x_r^{\alpha_r} \quad .$$

Back to inv-maj

In the present approach, we need to put up with the more general discrete function $F(n,i)(p,q)$ even though ultimately we are only interested in $H(n)(p,q) = F(n+1,n+1)(p,q)$. We will prove the stronger statement that even if you restrict attention to those $(n-1)!$ permutations that end with a specific i, the pair (inv, maj) is still asymptotically-joint-independently normal.

Since the averages of both inv and maj over the permutations that end in i is $n-i+(n-1)(n-2)/4$ [for inv it is obvious, the last entry i contributes $n-i$ to the number of inversions, and removing the last entry yields an $n-1$-permutation, and for maj this follows from Foata's bijection mentioned above that maps inv to maj preserving the last entry], the *centralized probability generating function* corresponding to $F(n,i)(p,q)$ is:

$$G(n,i)(p,q) := \frac{F(n,i)(p,q)}{(n-1)!(pq)^{n-i+(n-1)(n-2)/4}} \quad .$$

The recurrence $(RecF)$ becomes

$$G(n,i) = \frac{1}{q}G(n,i+1) + \frac{p^{n-i}(q^{n-1}-1)}{(pq)^{n/2}(n-1)}G(n-1,i) \quad , \qquad (RecG)$$

and the *final condition* becomes

$$G(n,n) = \frac{1}{n-1}\sum_{j=1}^{n-1}(pq)^{n/2-j}G(n-1,j) \quad . \qquad (Gnn)$$

We also need the obvious *initial condition* $G(1,i) = \delta_{i,1}$.

Guessing Polynomial Expressions for the Factorial Moments

Equipped with these efficient recurrences, our computer computes $G(n,i)(p,q)$ for many values of n and i. Then for each of these it computes the (r,s) (mixed) factorial moments $FM(r,s)(n,i)$, for many small numeric (r,s) by computing the (initial terms of the) Maclaurin series for $G(n,i)(1+p,1+q)$. We then fix numeric values of (r,s) and use *polynomial interpolation* to *guess* explicit polynomial expressions for $FM(r,s)(n,i)$ as polynomials in (n,i) for that pair (r,s). The process is repeated for all pairs (r,s) for which $0 \le r,s \le M$, for some pre-determined specific positive integer M. We note that it is obvious, both from the combinatorics and from the recurrences, that the $FM(r,s)(n,i)$ are always polynomials in (n,i), for any fixed numeric r and s.

[As we have already mentioned in footnote 2, most of the referees didn't find this obvious. The proof via the recurrences $(RecG')$ and (Gnn') to be derived in page 10 is by induction on (r,s) and the fact that the indefinite sum of a polynomial (in this case with respect to i) is yet another polynomial, and the "operator" on the right of (Gnn') is polynomial-preserving.

Let's sketch the combinatorial proof that the mixed moments $M_{r,s}(n)$ are polynomials in n. The combinatorial proof that $FM(r,s)(n,i)$ are polynomials in both n and i is similar. Write $inv(\pi)$ and $maj(\pi)$ as a sum of "atomic" events, e.g. for inv the sum of $\chi(\pi_j > \pi_i)$ over all pairs of integers (i,j) satisfying $1 \le i < j \le n$. Here

$\chi(S) = 1$ if S is true and $\chi(S) = 0$ if it is false. maj can be similarly expressed as a sum of $\chi(i \leq j \ AND \ \pi_j > \pi_{j+1})$. The sum of $inv(\pi)^r maj(\pi)^s$ over all permutations $\pi \in S_n$ can be expressed as a multi-sum with the outer sum ranging over S_n and the inner sums with $2(r+s)$ sigma signs. Now do discrete Fubini! Bring the formerly outer-sigma, over S_n, *all the way inside* past all the other $2(r+s)$ sigma signs. Each individual sigma sign involves one index of summation, and collectively these $2(r+s)$ sigma signs involve $\leq 2(r+s)$ such indices, corresponding to locations in a generic n-permutation π, some of whom may coincide. Let's call them $i_1 < ... < i_k$ (where $k \leq 2r+2s$). These can be placed in lots of possible intertwining ways, and so can the values of π in those places. There are still finitely many scenarios. (Formally, one gets a Cartesian product of two partially ordered sets, one for the domain and one for the range, each of which has finitely many linear extensions. This is reminiscent of Richard Stanley's theory of P-partitions). For each particular such scenario (linear extension of the domain-poset and the range-poset), there are $\binom{n}{k}$ ways to choose the participant indices, and $\binom{n}{k}$ ways to choose their occupants (i.e. the values of π there) and the remaining $n-k$ entries can, of course, be arranged in $(n-k)!$ ways, yielding $\binom{n}{k}^2 (n-k)!$ ways. Dividing by $n!$ gives $\binom{n}{k}^2 (n-k)!/n! = \binom{n}{k}/k!$, a polynomial in n of degree k. Now the whole thing is a sum of finitely many such $\binom{n}{k}/k!$ for $0 \leq k \leq 2r+2s$, and since a finite sum of polynomials is still a polynomial, we are done! Now isn't that *obvious*?!]

It would have been nice if we could guess *closed-form* expressions for $FM(r,s)(n,i)$ for *symbolic* (r,s), but no such closed-form exists as far as we know, and besides it is too much to ask for and more than we need. To prove *asymptotic normality* we only need the *leading terms*. Viewing the leading terms, our beloved computer easily conjectures the following expressions. For integers $r \geq 0, s \geq 0$ we have:

$$FM(2r,2s)(n,i) = \frac{(2r)!}{2^r r!}\frac{(2s)!}{2^s s!}\left(\frac{1}{36}\right)^{r+s} n^{3r+3s} + (lower-total-degree-terms-in-(n,i)) \quad .$$

[Note that the coefficients of $n^{3r+3s-1}i, n^{3r+3s-2}i^2$ etc. are all zero, hence they don't show up!].

For integers $r \geq 0, s \geq 1$ we have:

$$FM(2r,2s-1)(n,i) = \frac{(2r)!}{2^r r!}\frac{(2s)!}{2^s s!}\left(\frac{1}{36}\right)^{r+s-1} n^{3r+3s-6}[-(s-1)n^3 - 6rn^2 i + 18rni^2 - 12ri^3]$$

$$+ (lower-total-degree-terms-in-(n,i)) \quad .$$

For integers $r \geq 1, s \geq 0$ we have:

$$FM(2r-1,2s)(n,i) = -\frac{(2r)!}{2^r r!}\frac{(2s)!}{2^s s!}\left(\frac{1}{36}\right)^{r+s-1} (r-1)n^{3r+3s-3} + (lower-total-degree-terms-in-(n,i)) \quad .$$

For integers $r \geq 1, s \geq 1$ we have:

$$FM(2r-1,2s-1)(n,i) = \frac{(2r)!}{2^r r!}\frac{(2s)!}{2^s s!}\left(\frac{1}{36}\right)^{r+s-1} \frac{9}{2} n^{3r+3s-6}(n-2i)^2 + (lower-total-degree-terms-in-(n,i)) \quad .$$

Nice conjectures but what about proofs?

While we prefer the empirical approach of guessing, an alternative approach to finding many $FM(r,s)(n,i)$'s, that is also necessary in order to rigorously prove our conjectures, is to first use $(RecG)$ and (Gnn). Write $G(n,i)(1+p,1+q)$ as an *infinite* **generic** Taylor series around $(0,0)$, and write down the implied *infinite-order* recurrences expressing $FM(r,s)$ in terms of $FM(r',s')$ with $r'+s'<r+s$.

[The infinite-order recurrence for the $FM(r,s)(n,i)$, obtained from $(RecG)$ is gotten by expanding $1/(1+q)$ and

$$\frac{(1+p)^{n-i}(1+q)^{n-i}-1}{((1+p)(1+q))^{n/2}}$$

as Maclaurin series in (p,q), using the binomial theorem and manipulations on formal power series (that Maple does automatically to any desired order), and combining terms. Similarly, the implication of (Gnn) is obtained by expanding $((1+p)(1+q))^{n/2-j}$ using the binomial theorem. Both of these tasks are accomplished by procedure `MOP` in the Maple package `InvMaj`.]

Note that in order to compute $FM(r,s)(n,i)$, for any specific, numeric r and s, we only need finitely many terms (actually $rs-1$ of them) of the infinite-order recurrence, since eventually all the contributions will be zero. Of course, as we have already commented, there is no hope for finding a general expression for $FM(2r,2s)(n,i)$, $FM(2r,2s-1)(n,i)$, $FM(2r-1,2s)(n,i)$ and $FM(2r-1,2s-1)(n,i)$, depending *explicitly* on r and s (i.e. **symbolically** in terms of r and s) as well as on n and i, but to prove, by induction on r,s, that the above *leading terms* are valid, all we need to do is to verify that the leading terms of the implied recurrences for the $FM(r,s)$'s (that we have just talked about) are consistent with the above explicit expressions.

The implication of $(RecG)$ is

$$FM(r,s)(n,i)-FM(r,s)(n,i+1)=$$

$$-sFM(r,s-1)(n,i+1)+sFM(r,s-1)(n-1,i)+\frac{rs(n-2i)}{2}FM(r-1,s-1)(n-1,i)+$$

$$(lower-order-terms) \quad , \qquad (RecG')$$

while the implication of (Gnn) is:

$$FM(r,s)(n,n)=\frac{1}{n-1}\sum_{j=1}^{n-1}FM(r,s)(n-1,j)-\frac{s}{2(n-1)}\sum_{j=1}^{n-1}(2j-n)FM(r,s-1)(n-1,j)-$$

$$\frac{r}{2(n-1)}\sum_{j=1}^{n-1}(2j-n)FM(r-1,s)(n-1,j)+\frac{rs}{4(n-1)}\sum_{j=1}^{n-1}(2j-n)^2FM(r-1,s-1)(n-1,j)$$

$$+(lower-order-terms) \quad . \qquad (Gnn')$$

[Referee Guoniu Han correctly commented that one must say something about the degree of the polynomials in (n,i) that reside in the "lower-order terms" that feature as coefficients in these infinite-order recurrences. It turns out that

the coefficient of $FM(r-r',s-s')(n-1,i)$ in $(RecG')$ and the coefficient of $FM(r-r',s-s')(n,i+1)$ (and similarly for (Gnn')) are polynomials in (n,i) of total degree $\leq r'+s'$, thanks to the binomial theorem, and this would imply, by induction on (r,s) that not only are $FM(2r,2s)(n,i)$ etc. polynomials, but their leading terms look as claimed above.]

The next step is to spell out these two recurrences each into the four cases according to whether (r,s) is $(even,even)$, $(even,odd)$, $(odd,even)$, and (odd,odd).

Once you have these eight recurrences, for each and every one of them, you plug in the above conjectured expressions for the leading terms, and verify that up to the leading terms, things agree. At the end of the day, after dividing by

$$\frac{(2r-2)!}{(r-1)!2^{r-1}}\frac{(2s-2)!}{(s-1)!2^{s-1}} \quad ,$$

this boils down to proving equalities among certain low-degree polynomials in (r,s) (namely the leading coefficients, in (n,i)), that in turn, reduces (since $A=B$ iff $A-B=0$) to proving that certain low-degree polynomials in (r,s) are identically zero.

So in order to check that these low-degree polynomials in (r,s) are all identically zero, it is enough to check each of them for finitely (and not-too-many) *numeric* r,s. Typing `Check1(FM8m(n,i),n,i);` and `Check2(FM8m(n,i),n,i);` in the Maple package `InvMaj` does exactly that, by checking that if you plug in the conjectured leading terms of $FM(2r,2s)(n,i)$, $FM(2r-1,2s)(n,i)$, $FM(2r,2s-1)(n,i)$ and $FM(2r-1,2s-1)(n,i)$ and subtract the right sides from the left sides (for each of the eight cases) you get lower-order polynomials, in (n,i), for all $1\leq r,s\leq 8$. This proves all these claims (rigorously), with a vengeance! The $(8/2)^2=16$ special cases are much more than is needed, since the relevant polynomials in (r,s) are easily seen to have degree ≤ 2 so $(2+1)^2=9$ agreements would have sufficed.

The Maple package InvMaj

All the nitty-gritty calculations described above, that constitute **a fully rigorous** proof, may be found in the Maple package `InvMaj` accompanying this article. This package is available from the webpage of the present article:

`http://www.math.rutgers.edu/~zeilberg/mamarim/mamarimhtml/invmaj.html` ,

where the reader can also find some sample input and output. The direct url of the package is:
`http://www.math.rutgers.edu/~zeilberg/tokhniot/InvMaj` .

La Grande Finale

The special cases $r=1,s=0$ and $r=0,s=1$ of the now proved formula for $FM(2r,2s)$ (displayed at the middle of page 9) give

$$FM(2,0)(n,i)=\frac{1}{36}n^3+O(n^2) \quad ,$$

$$FM(0,2)(n,i) = \frac{1}{36}n^3 + O(n^2) \quad .$$

So (recall that we are interested in the *normalized* mixed factorial moments)

$$\frac{FM(2r,2s)(n,i)}{FM(2,0)(n,i)^r FM(0,2)(n,i)^s} = \frac{(2r)!}{2^r r!}\frac{(2s)!}{2^s s!} + O(1/n) \quad ,$$

$$\frac{FM(2r,2s-1)(n,i)}{FM(2,0)(n,i)^r FM(0,2)(n,i)^{s-1/2}} = o(1/n) \quad ,$$

$$\frac{FM(2r-1,2s)(n,i)}{FM(2,0)(n,i)^{r-1/2} FM(0,2)(n,i)^s} = o(1/n) \quad ,$$

$$\frac{FM(2r-1,2s-1)(n,i)}{FM(2,0)(n,i)^{r-1/2} FM(0,2)(n,i)^{s-1/2}} = O(1/n) \quad .$$

And we see that as $n \to \infty$ these indeed converge to the mixed moments of the famous mixed moments of the bivariate independent normal distribution $e^{-a^2/2-b^2/2}/(2\pi)$ □.

Encore: A more refined asymptotics for the (Normalized) Mixed Moments

With more effort, we (or rather, our computer) can guess-and-prove the following asymptotics for the case of interest $(n+1, n+1)$, i.e. the asymptotic expressions for the centralized-and-normalized (genuine, not factorial) mixed-moments, let's call them $\alpha(r,s)(n)$, for the pair of random variables (inv, maj) acting on the set of permutations of length n. Indeed, according to S. B. Ekhad, we have:

$$\alpha(2r,2s)(n) = \frac{(2r)!}{2^r r!}\frac{(2s)!}{2^s s!}\left(1 - \frac{9(r^2+s^2-r-s)}{25}\cdot\frac{1}{n} + O(\frac{1}{n^2})\right) \quad ,$$

$$\alpha(2r-1,2s-1)(n) = \frac{(2r)!}{2^r r!}\frac{(2s)!}{2^s s!}\left(\frac{9}{2n} + \left(-\frac{81}{50}(r^2+s^2) + \frac{243}{50}(r+s) - \frac{1773}{100}\right)\frac{1}{n^2} + O(\frac{1}{n^3})\right) \quad .$$

Of course, by *symmetry* (see page 4) $\alpha(2r,2s-1)(n)$ and $\alpha(2r-1,2s)(n)$ are *identically* (not just asymptotically!) zero.

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