# Efficient algorithms for tensor scaling, quantum marginals, and moment polytopes



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• Marginals: marginals of joint probability distributions with constrained supports



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- History, special cases, and applications
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- Algorithm
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# Quantum marginals, moment polytopes, and tensor scaling

If Alice, Bob, and Carol's qubits are jointly in a pure quantum state, the one-body marginals are 2 × 2 PSD matrices  $\rho^{(A)}, \rho^{(B)}, \rho^{(C)}$ 



**One body quantum marginal problem,** d = 3**Input:** PSD matrices  $P_A$ ,  $P_B$ ,  $P_C$ **Output:** Whether there is a pure state with marginals  $P_A$ ,  $P_B$ ,  $P_C$ 

**Fact:** the answer depends only on spec( $P_A$ ), spec( $P_B$ ), spec( $P_C$ ).

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*d*-tensors: denoted  $(\mathbb{C}^n)^{\otimes d}$ 

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Three different ways of viewing a  $n \times n \times n$  tensor as an  $n \times n^2$  matrix.



Flattening

 $i^{th}$  flattening  $M^{(i)}$  of X is an  $n \times n^{d-1}$  matrix rows = slices of X orthogonal to the  $i^{th}$  direction

Marginal

 $i^{th}$  marginal of X is the  $n_i imes n_i$  matrix

$$\rho^{(i)} = M^{(i)} \left( M^{(i)} \right)^{\dagger}$$

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 $\Delta(\mathbb{X}) = \{ \text{tuples of spectra of marginals of elements of } \mathbb{X} \text{ with trace } 1 \}$ 

e.g. ((.75, .25), (.5, .5), (.9, .1)) for 3 qubits

Amazing fact [Mumford, Ness '84] For many natural X, including those in this talk,  $\Delta(X)$  is a

called the moment polytope; can be defined in much greater generality

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## Moment polytope membership for $\mathbb{X} = (\mathbb{C}^n)^{\otimes d}!$

Quantum marginal problem Input: p list of spectra Output: Whether p is in the moment polytope  $\Delta(X)$  for  $X = (\mathbb{C}^n)^{\otimes d}$ . Moment polytope membership for  $\mathbb{X} = (\mathbb{C}^n)^{\otimes d}!$ 

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Given X, can we locally change basis to obtain specific marginals?

- Changing basis on the  $i^{th}$  vector space by  $g_i$  changes flattening by  $M^{(i)} \gets g_i M^{(i)}$
- Scaling: Simultaneous local basis change  $(g_1, \ldots, g_d) \cdot X := g \cdot X$

**Question:** Tensor scaling

**Input:** *p* list of spectra, tensor *X* 

**Output:** whether there exists scalings of X with spectra of marginals approaching **p** 

• say  $G \cdot X$ : set of all scalings of X,  $\overline{G \cdot X}$  its *closure*  **Tensor scaling** is moment polytope membership for  $\mathbb{X} = \overline{G \cdot X}$ ! **Fact: Quantum marginal** problem for  $p \iff$  **tensor scaling** problem for p and *random* X!

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# History, special cases, and applications

## Applications of tensor scaling

#### Matrix scaling: reweighting rows and columns

- deterministically approximating permanent
- **Operator scaling: tensor scaling with two marginals**
- noncommutative rational identity testing
- Forster's radial isotropic position
- computing the Brascamp-Lieb constant in analysis
- Horn's problem on eigenvalues of sums of matrices
- Uniform tensor scaling: target marginals  $I_n/n$
- Nullcone problem in invariant theory: do all invariants vanish on X?
- equivalence under SLOCC to locally maximally mixed state
- One body quantum marginal problem: tensor scaling for random XThe *Kronecker polytope* in representation theory
- Nonuniform tensor scaling:
- entanglement polytopes: comparing different types of entanglement
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## Nonuniform tensor scaling:

- entanglement polytopes: comparing different types of entanglement
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Produce scaling of input X with marginals  $\varepsilon$ -close to input p

#### Approximate quantum marginals:

Produce tensor with marginals  $\varepsilon$ -close to input **p** 

#### Approximate moment polytope membership:

Correctly output one of

## $\boldsymbol{p} \in \Delta(\mathbb{X}) + B(\varepsilon) \text{ or } \boldsymbol{p} \in \Delta(\mathbb{X})^{c} + B(\varepsilon).$

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### One body quantum marginal problem (nonuniform)

• [Higuchi,Sudbery'02] qubits; [Klaychko'04] polytope, [WDGC13] algebraic algorithm, [BCMW'17]: membership in **NP** ∩ **coNP** 

#### Theorem (BFGOWW '18, Tensor scaling)

There is a randomized poly( $n^d$ ,  $\langle X \rangle + \langle p \rangle$ ,  $1/\varepsilon$ )-time algorithm for approximate tensor scaling on input  $X, p, \varepsilon$  with success probability 1/2.

The algorithm requires  $O\left(dn^{2d} \frac{\langle X \rangle + \langle p \rangle + d \log dn}{\varepsilon^2}\right)$  iterations, each dominated by computing a Cholesky decomposition of some  $n \times n$  matrix. **Corollary (BFGOWW '18, Quantum marginals)** There is a randomized poly $(\langle p \rangle, 1/\varepsilon)$ -time algorithm for approximate quantum marginals on input  $p, \varepsilon$  with success probability 1/2.

Corollary (BFGOWW '18, Approximate moment polytope membership)

There is a randomized algorithm for approximate moment polytope membership running in time poly( $n^d$ ,  $\langle parameterization of X \rangle$ , p,  $1/\varepsilon$ )

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**Input:** X, p with integer coordinates,  $\varepsilon$ .

**Output:** scaling of X s.t. spectra of marginals  $\varepsilon$ -close to p, or OUTSIDE POLYTOPE

- Repeat *T* times:
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(ignoring damage done to other marginals!)

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For 
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# Analysis

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#### Our lower bounds are objects arising from representation theory

Definition (Highest weights of type p)

Homogeneous polynomials on  $(\mathbb{C}^n)^{\otimes d}$  that are eigenfunctions for the action of the lower triangular matrices;

 $P(g \cdot Y) = scalar(g, p)P(Y)$ 

scalar(g, p) is just the *eigenvalue*.

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**Open problems** 

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# Thank you!











