## 8 The Ordering on $\omega$

**Definition 8.1.**  $\in_{\omega}$  is the binary relation on  $\omega$  defined by

$$\in_{\omega} = \{ \langle m, n \rangle \in \omega \times \omega \mid m \in n \}.$$

In this section, we shall prove:

**Theorem 8.2.**  $\in_{\omega}$  is a linear order on  $\omega$ .

**Definition 8.3.** A set A is *transitive* iff whenever  $x \in a \in A$ , then  $x \in A$ .

Example 8.4.

- 1.  $\{\{\emptyset\}\}\$  is *not* transitive, since  $\emptyset \in \{\emptyset\} \in \{\{\emptyset\}\}\$  but  $\emptyset \notin \{\{\emptyset\}\}$ .
- 2.  $\{\emptyset, \{\emptyset\}\}$  is transitive.

**Lemma 8.5.** If  $n \in \omega$ , then n is transitive.

*Proof.* It is enough to prove that the set

$$T = \{ n \in \omega \mid n \text{ is transitive} \}$$

is inductive. First  $\emptyset$  is trivially transitive and so  $\emptyset \in T$ . Next suppose that  $n \in T$  and that

$$x \in a \in n^+ = n \cup \{n\}.$$

There are two cases to consider.

**Case 1** Suppose that  $a \in n$ . Since  $n \in T$  and  $x \in a \in n$ , is follows that  $x \in n$  and so  $x \in n^+$ .

**Case 2** Suppose that a = n. Then  $x \in n$  and so  $x \in n^+$ . Thus in both cases,  $x \in n^+$ . Hence  $n^+ \in T$ .

**Remark 8.6.** In other words, if  $a, b, c \in \omega$ , then

 $a \in b$  and  $b \in c$  implies  $a \in c$ .

Thus  $\in_{\omega}$  is a transitive relation on  $\omega$ .

## Lemma 8.7.

(a) For any  $n, m \in \omega$ ,

 $m \in n$  iff  $m^+ \in n^+$ .

(b) For all  $n \in \omega$ ,  $n \notin n$ .

*Proof.* (a) First suppose that  $m^+ \in n^+$ . Then

$$m \in m^+ \in n^+ = n \cup \{n\}.$$

There are two cases to consider.

**Case 1** Suppose that  $m^+ \in n$ . Then  $m \in m^+ \in n$  and so  $m \in n$ .

**Case 2** Suppose that  $m^+ = n$ . Then  $m \in n$ .

Thus in either case,  $m \in n$ .

To prove the converse, we use induction. In other words, we prove that

$$T = \{ n \in \omega \mid (\forall m \in n) \ m^+ \in n^+ \}$$

is inductive. First  $\emptyset \in T$  vacuously. Next suppose that  $n \in T$ . We must prove that

(\*) if  $m \in n^+$ , then  $m^+ \in n^{++}$ .

So suppose that  $m \in n^+ = n \cup \{n\}$ . If m = n then

$$m^{+} = n^{+} \in n^{++} = n^{+} \cup \{n^{+}\}.$$

Otherwise,  $m \in n$  and so since  $n \in T$ m

$$m^+ \in n^+ \subset n^{++}$$

and so  $m^+ \in n^{++}$ . Hence  $n^+ \in T$ . (b) It is enough to show that

$$S = \{ n \in \omega \mid n \notin n \}$$

is inductive. Clearly  $\emptyset \in S$ . Next suppose that  $n \in S$ . FOr the sake of contradiction, assume that  $n^+ \in n^+$ . By (a),  $n \in n$ , which contradicts the fact that  $n \in S$ . Thus  $n^+ \notin n^+$  and so  $n^+ \in S$ .

**Lemma 8.8.** For any  $n, m \in \omega$  at most one of the following holds:

$$m \in n, \quad m = n, \quad n \in m$$

*Proof.* By Lemma 8.7 (b), if two hold, then we must have that  $m \in n$  and  $n \in m$ . By Lemma 8.5,  $m \in m$ , which contradicts Lemma 8.7 (b).

**Lemma 8.9.** For any  $n, m \in \omega$  at least one of the following holds:

$$m \in n, \quad m = n, \quad n \in m.$$

2006/10/11

*Proof.* It is enough to show that

$$T = \{ n \in \omega \mid (\forall m \in \omega) \ (m \in n \text{ or } m = n \text{ or } n \in m) \}$$

is inductive.

**Exercise 8.10.** Prove that for all  $m \in \omega$ ,  $m = \emptyset$  or  $\emptyset \in m$ .

Hint Argue by induction on m.

Thus  $\emptyset \in T$ . Next suppose that  $n \in T$ . Let  $m \in \omega$  be arbitrary. Since  $n \in T$ , we have that

$$m \in n$$
 or  $m = n$  or  $n \in m$ .

If  $m \in n$  or m = n, then  $m \in n^+ = n \cup \{n\}$ . If  $n \in m$ , then Lemma 8.7 (a) implies

$$n^+ \in m^+ = m \cup \{m\}$$

and so  $n^+ \in m$  or  $n^+ = m$ . In either case, we have that

$$m \in n^+$$
 or  $m = n^+$  or  $n^+ \in m$ .

Thus  $n^+ \in T$ .

This completes the proof that  $\in_{\omega}$  is a linear order on  $\omega$ .

**Notation** (Different from Enderton) From now on, if  $m, n \in \omega$ , then we use the following notation interchangably:

$$\begin{array}{ll} m \in n & \text{iff} \quad m < n \\ m \stackrel{\epsilon}{-} n & \text{iff} \quad m \leq n \end{array}$$

**Exercise 8.11.** Let < be a linear order on a A. If  $a, b \in A$  satisfy  $a \leq b$  and  $b \leq a$  then a = b.

**Theorem 8.12 (Well-ordering of**  $\omega$ ). If  $\emptyset \neq A \subseteq \omega$ , then there exists  $m \in A$  such that  $m \leq a$  for all  $a \in A$ ; ie  $m \in a$  or m = a for all  $a \in A$ .

*Proof.* Assume that no such element exists. Define

$$B = \{ n \in \omega \mid (\forall k \in n) \ k \in \omega \backslash A \}$$

We shall prove that B is inductive. Clearly  $\emptyset \in B$  vacuously. Next suppose that  $n \in B$ . Thus (i) If  $k \in n$ , then  $k \notin A$ . Suppose that  $n^+ \notin B$ . Then there exists  $k \in n^+ = n \cup \{n\}$  such that  $k \in A$ . By (i) we must have that (ii)  $n \in A$  Now let  $a \in A$  be arbitrary. By Trichotomy, either

$$a \in n$$
 or  $a = n$  or  $n \in a$ .

By (i),  $a \notin n$ . Thus for all  $a \in A$ , a = n or  $n \in a$ , contradicting our assumption. Hence  $n^+ \in B$ . By Induction,  $B = \omega$ . But this means that  $A = \emptyset$ , which is a contradiction.  $\Box$ 

2006/10/11

**Theorem 8.13 (Strong Induction Principle for**  $\omega$ ). Let  $A \subseteq \omega$  and suppose that for every  $n \in \omega$ ,

(\*) if 
$$m \in A$$
 for all  $m < n$ , then  $n \in A$ .

Then  $A = \omega$ .

*Proof.* Suppose that  $A \neq \omega$ . Then  $\omega \setminus A \neq \emptyset$  and so there exists a least element  $k \in \omega \setminus A$ . Since k is the least such element, it follows that  $m \in A$  for all m < k. Then (\*) implies that  $k \in A$ , which is a contradiction.

**Definition 8.14.** Suppose that  $<_A, <_B$  are linear orders on A, B respectively. Then a function  $f: A \to B$  is order-preserving iff for all  $a_1, a_2 \in A$ ,

(\*) if  $a_1 <_A a_2$ , then  $f(a_1) <_B f(a_2)$ .

**Exercise 8.15.** Suppose that  $f: A \to B$  is order-preserving. Then the following statements are true.

- If  $a_1, a_2 \in A$ , then  $a_1 <_A a_2$  iff  $f(a_1) <_B f(a_2)$ .
- f is an injection.
- If f is a bijection, then  $f^{-1}: B \to A$  is also order-preserving.

**Theorem 8.16.** If  $f: \omega \to \omega$  is order-preserving, then  $f(n) \ge n$  for all  $n \in \omega$ .

*Proof.* If not, then

$$C = \{ n \in \omega \mid f(n) < n \} \neq \emptyset.$$

Let  $k \in C$  be the least element. Then f(k) < k. Since f is order-preserving, this implies that f(f(k)) < f(k). Hence  $f(k) \in C$ , which contradicts the minimality of k.

**Corollary 8.17.** If  $f: \omega \to \omega$  is an order-preserving bijection, then f(n) = n for all  $n \in \omega$ .

Proof. Since f is order-preserving,  $f(n) \ge n$  for all  $n \in \omega$ . Since  $f^{-1}$  is order-preserving,  $f^{-1}(n) \ge n$  for all  $n \in \omega$ . This implies that  $f(f^{-1}(n)) \ge f(n)$  and so  $n \ge f(n)$  for all  $n \in \omega$ .  $\square$  Hence f(n) = n for all  $n \in \omega$ .

**Remark 8.18.** The above remark fails for  $\mathbb{Z}$ . For example, the function  $f: \mathbb{Z} \to \mathbb{Z}$  defined by f(z) = z + 1 is an order-preserving bijection.

**Exercise 8.19.** Suppose that A is a transitive set. Then

- $\mathcal{P}(A)$  is also a transitive set.
- $A \subseteq \mathcal{P}(A)$ .

Informal discussion of  $V = \bigcup_{\alpha \in On} V_{\alpha}$ , where

$$V_0 = \emptyset$$
  

$$V_{\alpha+1} = \mathcal{P}(V_{\alpha})$$
  

$$V_{\lambda} = \bigcup_{\alpha < \lambda} V_{\alpha} \text{ for lim } \lambda$$

Axioms so far	
Extensionality	By the transitivity of $V_{\alpha}$
Empty Set	$\emptyset \in V_1$
Subset Axiom	If $A \in V_{\alpha}$ and $B \subseteq A$ , then $B \in V_{\alpha+1}$
Union Axiom	If $A \in V_{\alpha}$ , then $\bigcup A \in V_{\alpha+1}$
Pairing Axiom	If $A, B \in V_{\alpha}$ , then $\{A, B\} \in V_{\alpha+1}$
Powerset Axiom	If $A \in V_{\alpha}$ , then $\mathcal{P}(A) \in V_{\alpha+2}$
Infinity	$\omega \in V_{\omega+1}$
The following results will be crucial in our construction of $\mathbb{Z}$	

**Theorem 8.20.** For any  $m, n, p \in \omega$ , we have that

$$m < n$$
 iff  $m + p < n + p$ .

Proof. Reading Exercise, Enderton p. 85-86.

Corollary 8.21 (Cancellation Law). For any  $m, n, p \in \omega$ , if m + p = n + p, then m = n.

*Proof.* Suppose that m + p = n + p. By Trichotomy, if  $m \neq n$ , then either m < n or n < m. By Theorem 8.20, if m < n then m + p < n + p, which contradict Trichotomy. Similarly, if n < m, then n + p < m + p, which also contradicts Trichotomy. Hence m = n.

The following resuls will be crucial in our construction of  $\mathbb{Q}$ .

**Theorem 8.22.** If  $m, n, p \in \omega$  and  $p \neq 0$ , then

$$m < n$$
 iff  $m \cdot p < n \cdot p$ .

Proof. Reading Exercise, Enderton p. 85-86.

**Corollary 8.23 (Cancellation Law).** If  $m, n, p \in \omega$  and  $p \neq 0$ , then  $m \cdot p = n \cdot p$  implies m = n.

5