

## 4 Relations

Functions  $\equiv$  Operations

Relations  $\equiv$  Properties

**Remark 4.1.** In this course, we will only consider *binary* relations.

**Example 4.2.** Consider the order relation on  $A = \{0, 1, 2\}$  of *less than*. Then  $0 < 1$ ,  $0 < 2$ , and  $1 < 2$ . But  $2 \not< 1$ , *etc.* We shall define the order relation to be

$$< = \{\langle 0, 1 \rangle, \langle 0, 2 \rangle, \langle 1, 2 \rangle\}.$$

**Example 4.3.** Consider the relation  $E$  on  $A = \{0, 1, 2\}$  defined by

$$xEy \text{ iff } x - y \text{ is even.}$$

Then we shall define

$$E = \{\langle 0, 2 \rangle, \langle 2, 0 \rangle, \langle 0, 0 \rangle, \langle 1, 1 \rangle, \langle 2, 2 \rangle\}.$$

**Definition 4.4.** Let  $A$  be a set. Then  $R$  is a *relation* on  $A$  iff  $R \subseteq A \times A$ .

**Remark 4.5.** In this case, we would usually write  $xRy$  instead of  $\langle x, y \rangle \in R$ . *eg* “ $0 < 2$ ” is more natural than “ $\langle 0, 2 \rangle \in <$ ”.

**Definition 4.6 (More General Definition).** A *relation* is a set of ordered pairs.

*eg*  $R = \{\langle 0, 3 \rangle, \langle 1, 4 \rangle, \langle 2, 5 \rangle, \langle 3, 6 \rangle\}$  is a relation.

**Definition 4.7.** Let  $R$  be a relation. Then we define  $\text{dom } R$ ,  $\text{ran } R$ , and  $\text{fld } R$  by

$$x \in \text{dom } R \text{ iff there exists } y \text{ such that } \langle x, y \rangle \in R$$

$$x \in \text{ran } R \text{ iff there exists } y \text{ such that } \langle y, x \rangle \in R$$

$$\text{fld } R = \text{dom } R \cup \text{ran } R$$

**Remark 4.8.** Thus  $R$  is a relation on  $\text{fld } R$ .

**Example 4.9.** Let  $R = \{\langle l, n \rangle \mid l \text{ is the } n^{\text{th}} \text{ letter of the alphabet}\}$ . Then  $\text{dom } R = \{l \mid l \text{ is a letter of the alphabet}\}$  and  $\text{ran } R = \{1, 2, \dots, 26\}$ .

## 5 Equivalence Relations

**Definition 5.1.** Let  $R$  be a binary relation on  $A$ .

- $R$  is *reflexive* iff  $xRx$  for all  $x \in A$ .
- $R$  is *symmetric* iff for all  $x, y \in A$ , if  $xRy$  then  $yRx$ .
- $R$  is *transitive* iff for all  $x, y, z \in A$ , if  $xRy$  and  $yRz$  then  $xRz$ .
- $R$  is an *equivalence relation* iff  $R$  is reflexive, symmetric, and transitive.

**Example 5.2.** Let  $E$  be the relation on  $\mathbb{Z}$  defined by

$$xEy \quad \text{iff} \quad 3|x - y$$

Then  $E$  is an equivalence relation.

*Proof.* We check that  $E$  is reflexive, symmetric, and transitive.

- If  $x \in \mathbb{Z}$ , then  $3|x - x = 0$  since  $0 = 3 \cdot 0$ . Thus  $xEx$ .
- Suppose that  $xEy$ . Then  $3|x - y$  and so there exists  $z \in \mathbb{Z}$  such that  $x - y = 3 \cdot z$ . Then  $y - x = 3 \cdot (-z)$  and so  $3|y - x$ . Thus  $yEx$ .
- Suppose that  $xEy$  and  $yEz$ . Thus  $3|x - y$  and  $3|y - z$ . Hence there exist  $a, b \in \mathbb{Z}$  such that  $x - y = 3a$  and  $y - z = 3b$ . It follows that

$$\begin{aligned} x - z &= (x - y) + (y - z) \\ &= 3a + 3b \\ &= 3(a + b) \end{aligned}$$

and so  $3|x - z$ . Thus  $xEz$ . □

**Exercise 5.3.** Define a relation  $R$  on  $\mathbb{N} \times \mathbb{N}$  by

$$\langle a, b \rangle R \langle c, d \rangle \quad \text{iff} \quad a + d = c + b$$

Prove that  $R$  is an equivalence relation on  $\mathbb{N} \times \mathbb{N}$ .

**Definition 5.4.** Suppose that  $R$  is an equivalence relation on  $A$ . Then, for each  $x \in A$ , the  $R$ -equivalence class of  $x$  is

$$[x]_R = \{y \in A \mid xRy\}.$$

**Theorem 5.5.** Suppose that  $R$  is an equivalence relation on  $A$ .

- $x \in [x]_R$  for each  $x \in A$ .

- If  $x, y \in A$  and  $[x]_R \cap [y]_R \neq \emptyset$ , then  $[x]_R = [y]_R$ .

*Proof.*

- If  $x \in A$ , then  $xRx$  and so  $x \in [x]_R$ .
- Suppose that  $z \in [x]_R \cap [y]_R$ ; ie  $xRz$  and  $yRz$ . Since  $yRz$ , it follows that  $zRy$ . Since  $xRz$  and  $zRy$ , it follows that  $xRy$ . We claim that  $[y]_R \subseteq [x]_R$ . To see this, suppose that  $t \in [y]_R$ ; ie  $yRt$ . Since  $xRy$  and  $yRt$ , it follows that  $xRt$  and so  $t \in [x]_R$ . Similarly,  $[x]_R \subseteq [y]_R$  and so  $[x]_R = [y]_R$ .  $\square$

**Definition 5.6.** Let  $A$  be a set. Then  $\Pi$  is a *partition* of  $A$  iff:

- $\Pi$  is a set of nonempty subsets of  $A$ .
- If  $B, C \in \Pi$  are distinct, then  $B \cap C = \emptyset$ .
- $\bigcup \Pi = A$ .

**Example 5.7.**  $\Pi = \{0, 2\}, \{1, 3\}, \{4\}$  is a partition of  $A = \{0, 1, 2, 3, 4\}$ .

**Theorem 5.8.** Suppose that  $R$  is an equivalence relation on  $A$ . Then

$$\{[x]_R \mid x \in A\}$$

is a partition of  $A$ .  $\square$

**Example 5.9.** Let  $E$  be the equivalence relation on  $\mathbb{Z}$  defined by

$$xEy \text{ iff } 3 \mid x - y$$

Then

$$[0]_E = \{\dots, -6, -3, 0, 3, 6, \dots\}$$

$$[1]_E = \{\dots, -5, -2, 1, 4, 7, \dots\}$$

$$[2]_E = \{\dots, -4, -1, 2, 5, 8, \dots\}$$

Thus  $\{[0]_E, [1]_E, [2]_E\}$  is the corresponding partition of  $\mathbb{Z}$ .

**Theorem 5.10.** Suppose that  $\Pi$  is a partition of the set  $A$ . Define the relation  $R_\Pi$  by:

$$xR_\Pi y \text{ iff there exists } B \in \Pi \text{ such that } x, y \in B.$$

Then:

- $R_\Pi$  is an equivalence relation on  $A$ .
- $\Pi$  is the set of  $R_\Pi$ -equivalence classes.

*Proof.* Exercise! □

**Definition 5.11.** Suppose that  $R$  is an equivalence relation on  $A$ . Then the *quotient set* is

$$A/R = \{[x]_R \mid x \in A\}$$

and the *natural map / canonical map*  $\varphi: A \rightarrow A/R$  is defined by

$$\varphi(x) = [x]_R.$$